



# The Gathering Problem for Two Oblivious Robots with Unreliable Compasses

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Joint work with

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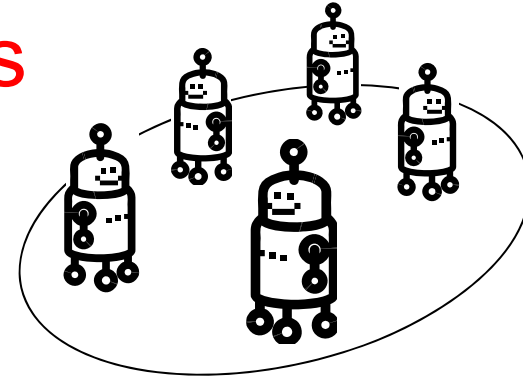
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Masafumi Yamashita\*(Kyusyu University, Japan)

\* : attendees of this meeting

# Coordination of Autonomous Mobile Robots

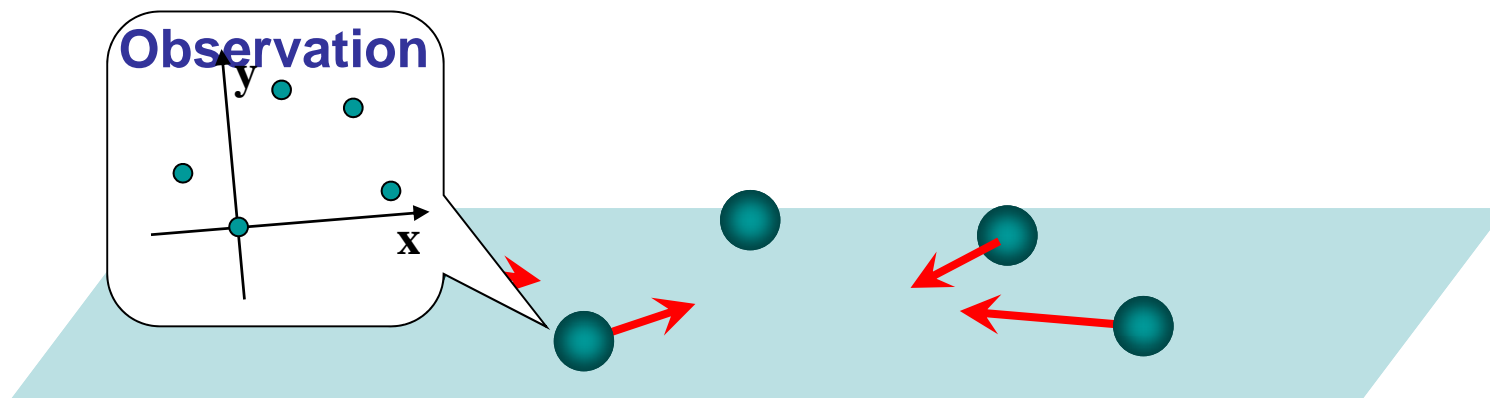
- **Autonomous Mobile Robots**  
Multiple, Fully decentralized



- Coordination task of Mobile Robots
  - Gathering, Convergence, Formation ...
- Challenges from the theoretical aspect
  - Clarifying the "weakest capability" to solve a given task

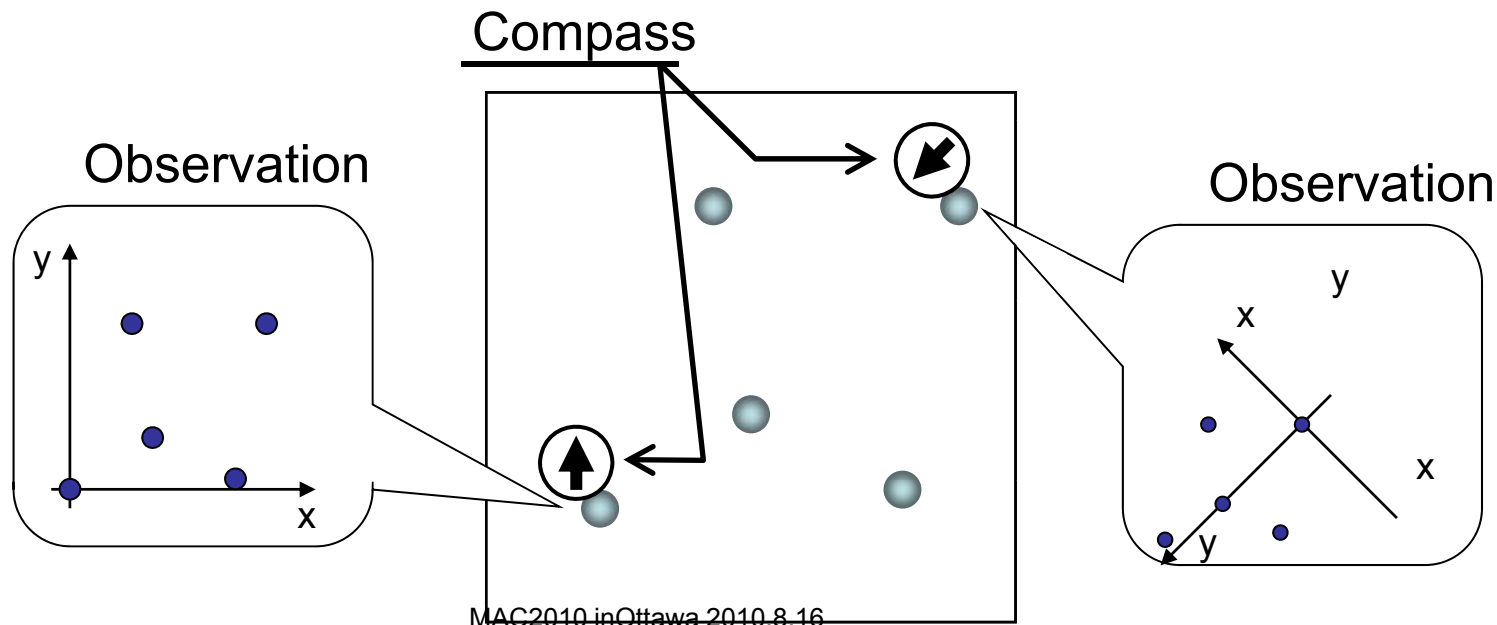
# Autonomous Mobile Robots

- Robot: Point on an infinite 2D-space
  - Anonymous (No distinguished ID)
  - Oblivious (No memory)
  - Deterministic
  - No communication (Observe the environment and Move)



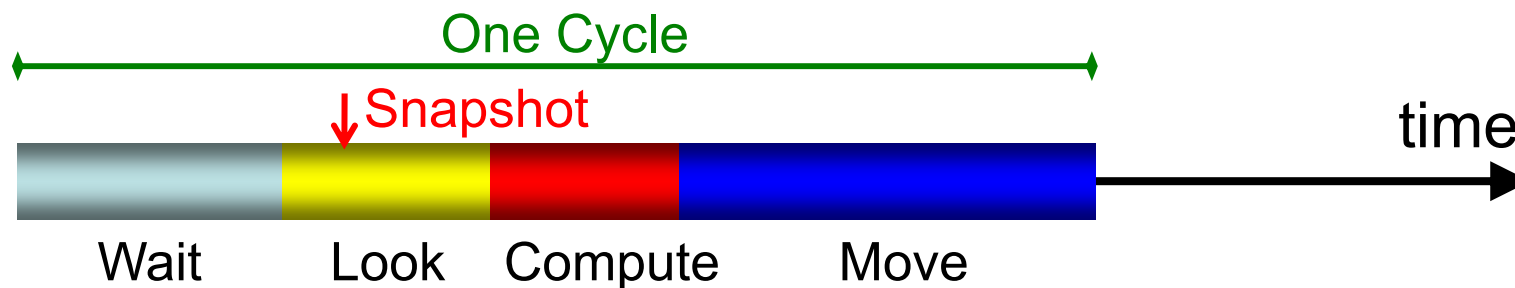
# Observation

- Each robot has a local x-y coordinate system(LCS)
  - The current position is the origin
  - The +direction of y-axis follows the **local compass**
- Agreement level of LCSs depends on the model (**compass model**)



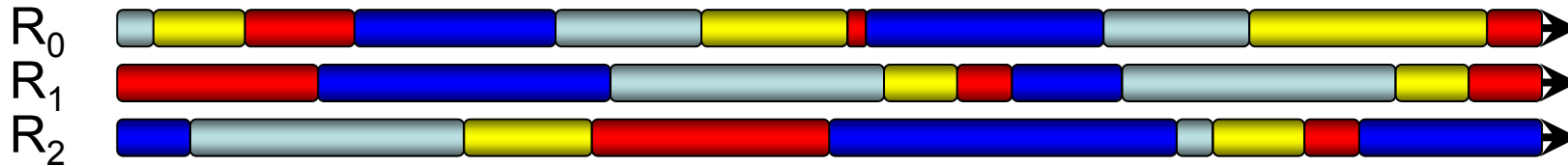
# Execution of Robots (Behavior of Each Robot)

- Wait-look-compute-move cycle
  - Wait: Idle state
  - Look: Take a snapshot of all robots' current locations (in terms of LCS)
  - Compute: Deciding the next position
  - Move: Move to the next position(unpredictable move)



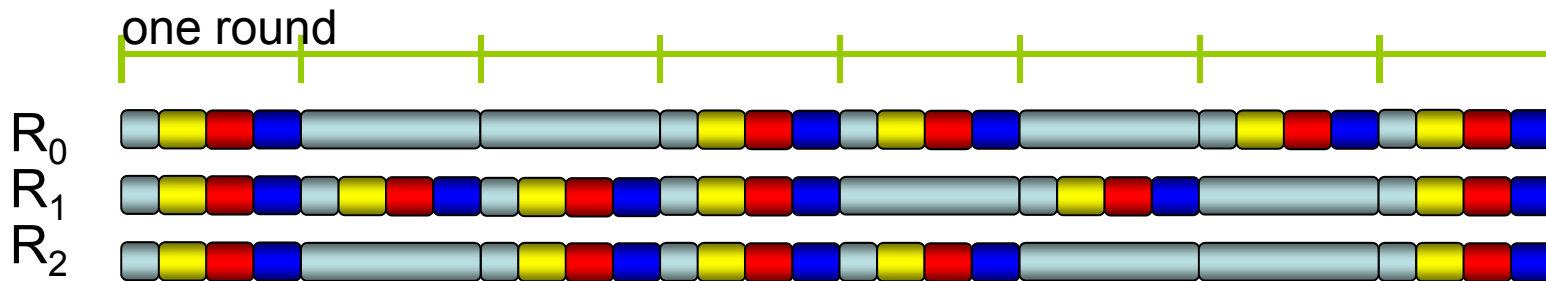
# Timing Model(How Cycles are Synchronized)

- Asynchronous(**CORDA**): No bound for length of each step

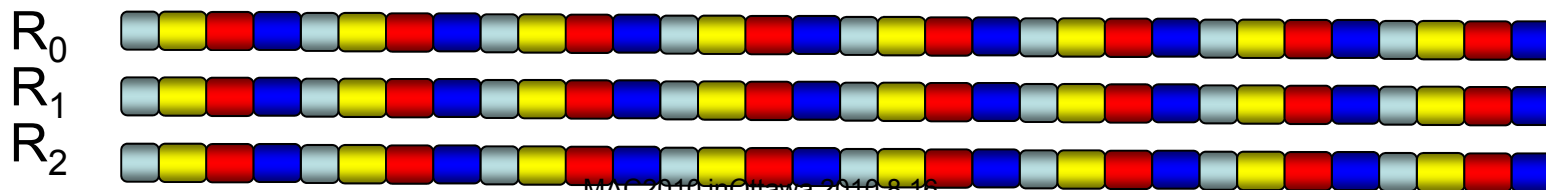


- Semi-synchronous(**SYm**, ATOM): Synchronized Round (one cycle=one round)

- Only a subset of all robots becomes **active** in each round



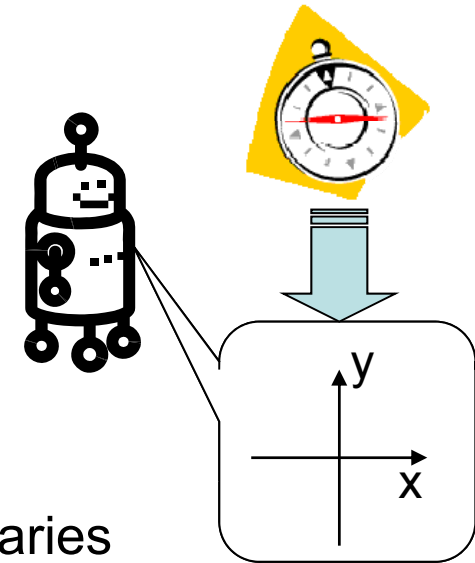
- Synchronous: All robots are completely synchronized



# Compass Models

## Inaccurate Compasses

- Every robot has its own local coordinate system
  - Compass gives y-axis' positive direction of the local coordinate system.
    - a compass varies, a local coordinate system varies



## ■ Inaccuracy of Compass

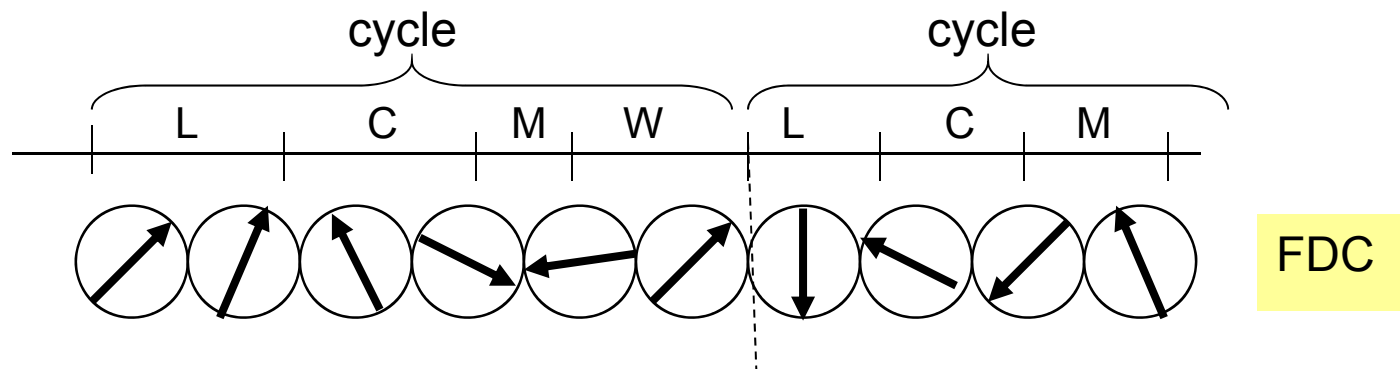
- **Variance of Compasses**
  - the variance of indicated directions of compasses
- **Deviation from the absolute direction**
  - the difference of indicated direction between compasses

# Compass Models – Variance –1

- Fully-Dynamic Compass(FDC)
- Semi-Dynamic Compass(SDC)
- FiXed Compass(FXC)

## Fully-dynamic Compass (FDC)

A compass whose indicated direction may **vary at any time during execution.**



Gathering is impossible on FDC.



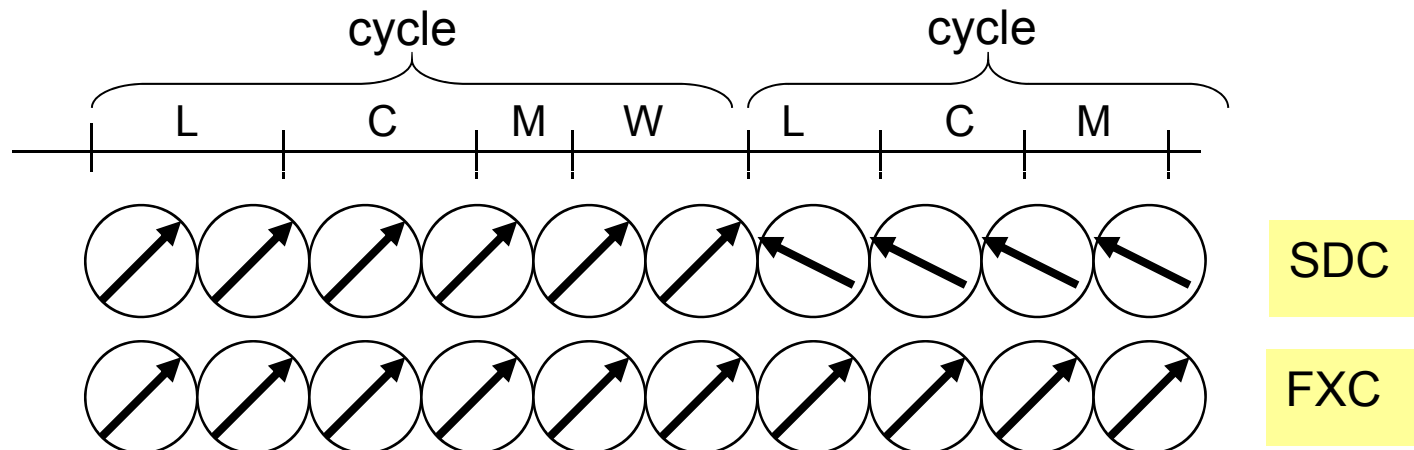
# Compass Models – Variance –2

## Semi-dynamic Compass (SDC)    Dynamic Compass

A compass whose indicated direction may **vary at the time between any two cycles** (never change during one cycle).

## Fixed Compass (FXC)    Static Compass

A compass whose indicated direction **never varies.**



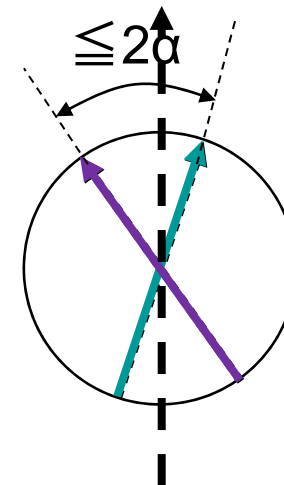
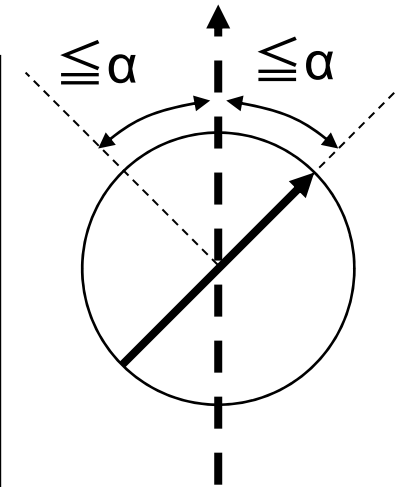
# Compass Models –Deviation–

## $\alpha$ -error Compass

A direction of “**the absolute north**” is assumed. The each angle which is formed by **the indicated direction of robots’ compass and the absolute north** is at most  $\alpha$ .

Note that the angle between two robots’ compasses is at most  $2\alpha$  on  $\alpha$ -error compass model

The absolute north



# Gathering Problem

- All robots meet at one point on a plane

- Not convergence

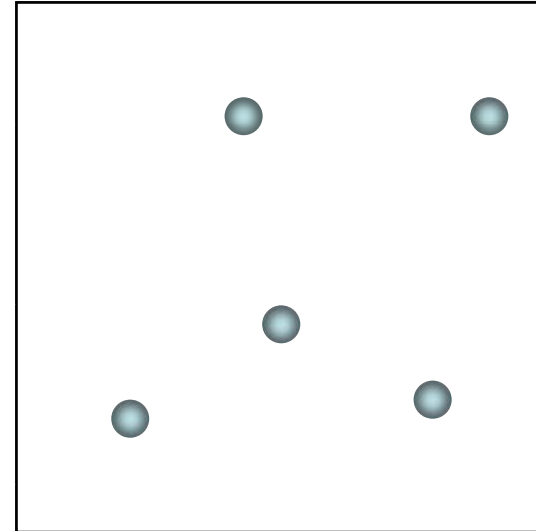
- Known Results

- Agreed Compass : Solvable

- CORDA / Arbitrary #robots

- Disagreed Compass : Unsolvable

- SYm / #robots = 2

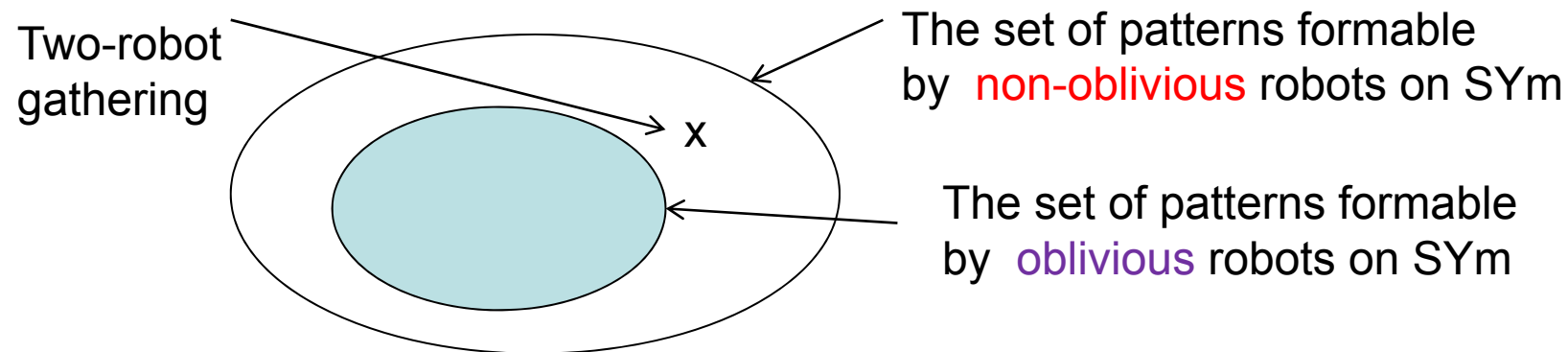


# Our results (summary)

## Two-robot Gathering problem on $\alpha$ -error compass

	SYm	CORDA
Semi-DC	impossible( $\alpha=\pi/4$ ) possible( $\alpha<\pi/4$ )	open possible( $\alpha<\pi/6$ )
FiXedC	impossible( $\alpha=\pi/2$ )[1] →	← possible( $\alpha<\pi/2$ )

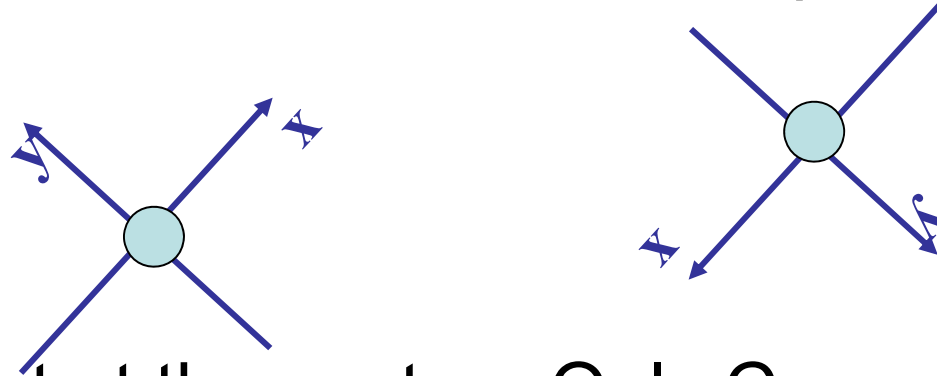
[1] I. Suzuki, M. Yamashita, SIAM J. Computing, 28, 4, 1347-1363, 1999.



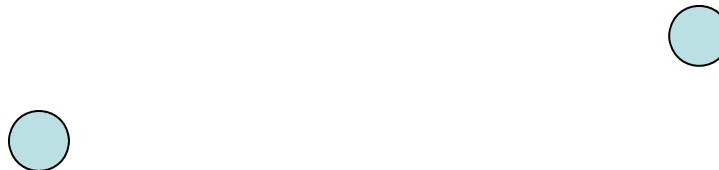
# Impossibility( $\pi/2$ -error compass , FXC and SYm)

- Opposite directions of two compasses

- Approach to another : Swap occurs



- Meet at the center : Only Convergence

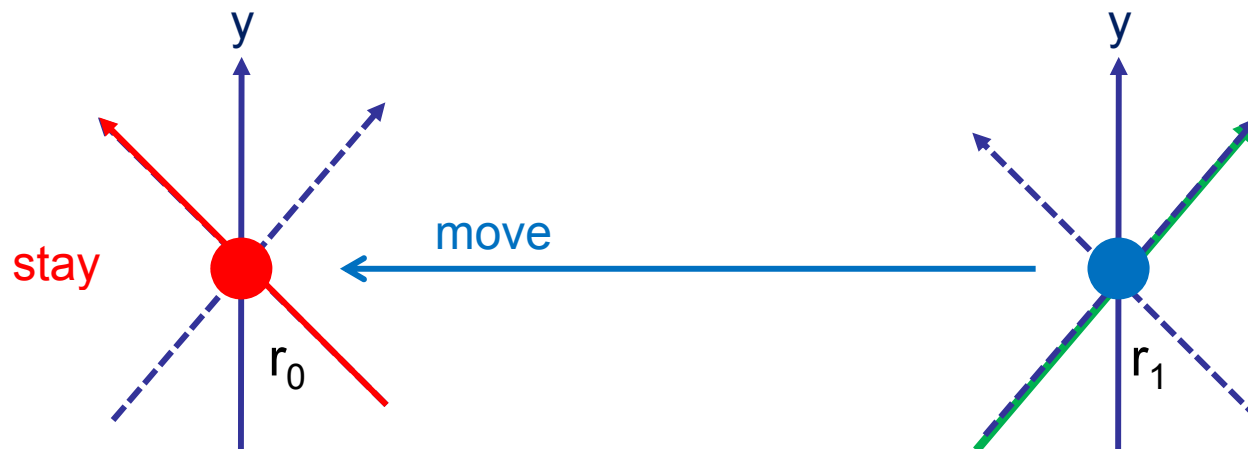


# Impossibility( $\pi/4$ -error compass , SDC and SYm)

A necessary condition for any gathering algorithm :

## stable configuration

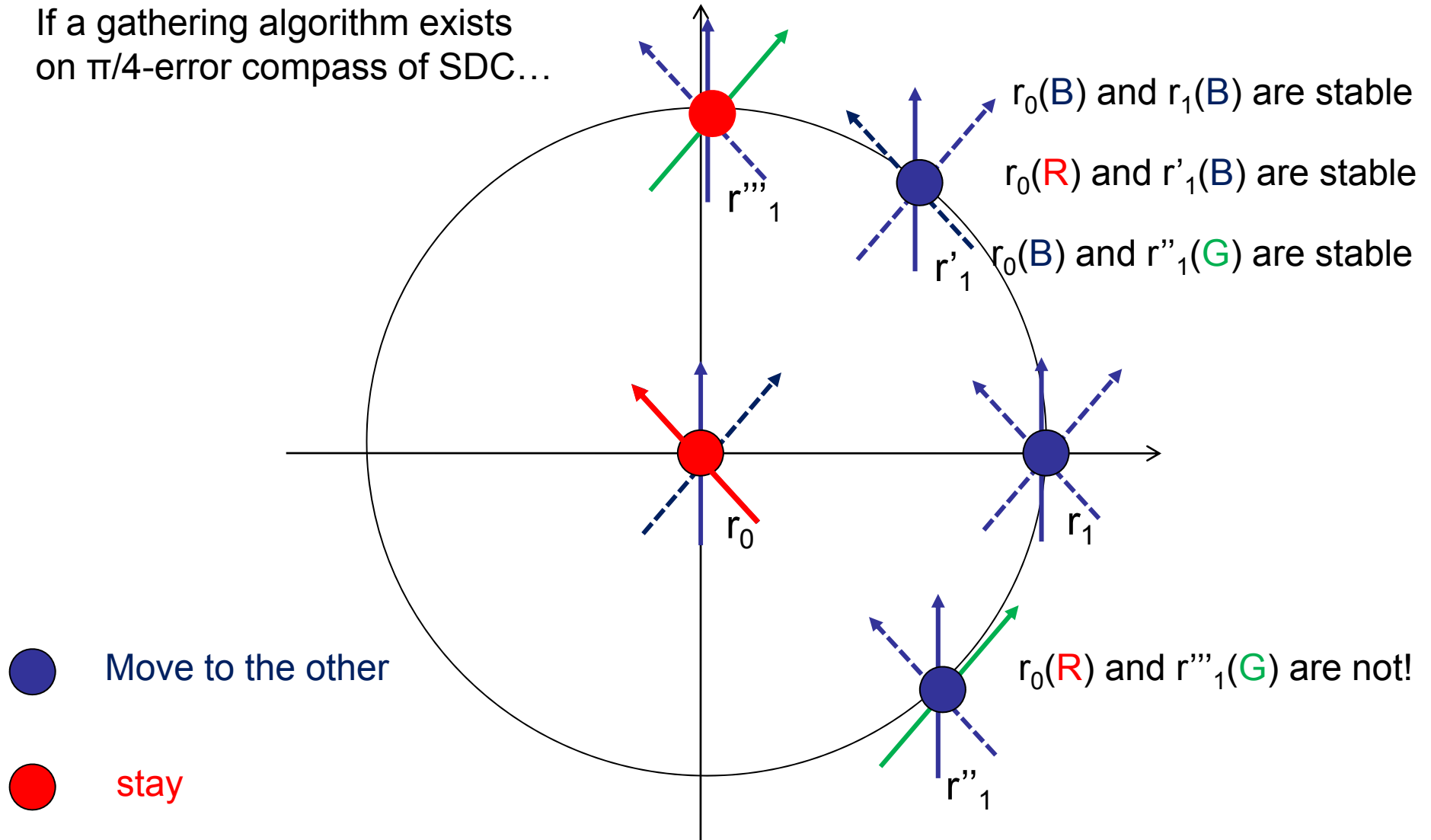
- a) There exists a configuration such that
  - 1) One robot  $r_0$  stays at own position
  - 2) Another robot  $r_1$  moves to the robot  $r_0$
- b) This configuration is regardless of the current local coordinate systems of both robots



Stable configuration

# Impossibility( $\pi/4$ -error compass , SDC and Sym)

If a gathering algorithm exists on  $\pi/4$ -error compass of SDC...



# Possibility results

## Two-robot Gathering problem on $\alpha$ -error compass

	SYm	CORDA
Semi-DC	impossible( $\alpha=\pi/4$ ) possible( $\alpha<\pi/4$ )	open possible( $\alpha<\pi/6$ )
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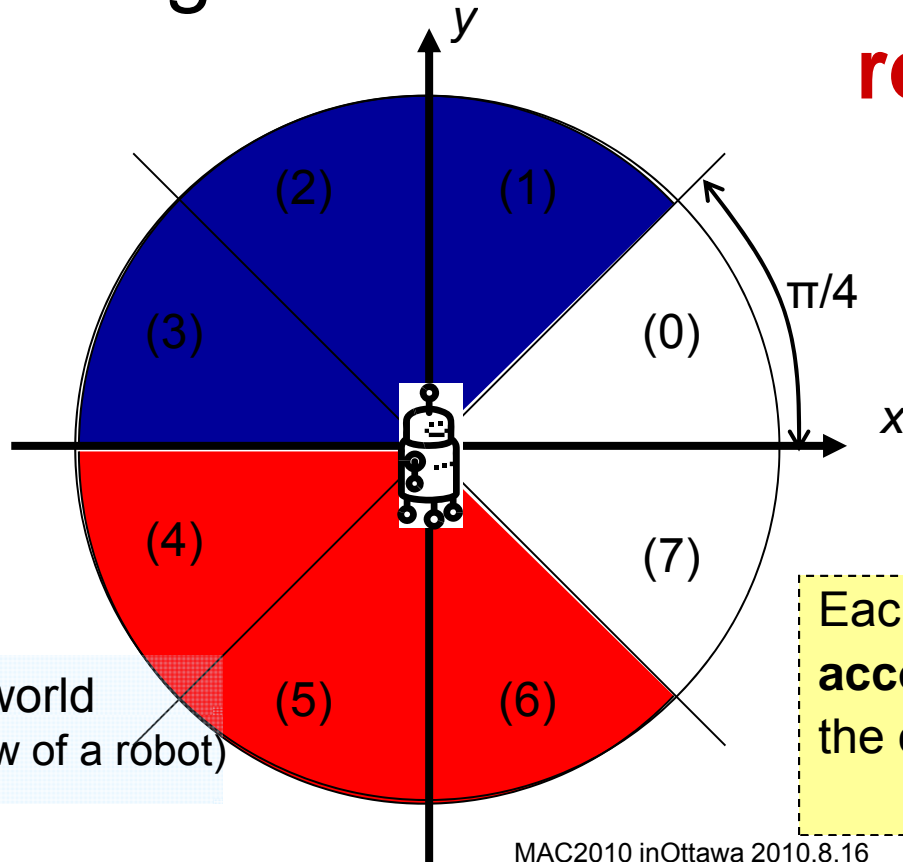


# $\pi/8$ -error SDC Algorithm on SYm

**Point:** How to decide the robots' behavior ?

Dividing the world (a view of a robot) into 8 sectors.  
Coloring the divided world with three colors:

**red, blue and white.**



The world  
(a view of a robot)

Each robot **decides its behavior according to the sector** in which the other robot is observed.

# $\pi/8$ -error SDC Algorithm on SYm

## Algorithm

Result of observing the other robot

case: no robot except me

gathering is achieved

case: in blue sectors (1), (2) or (3)

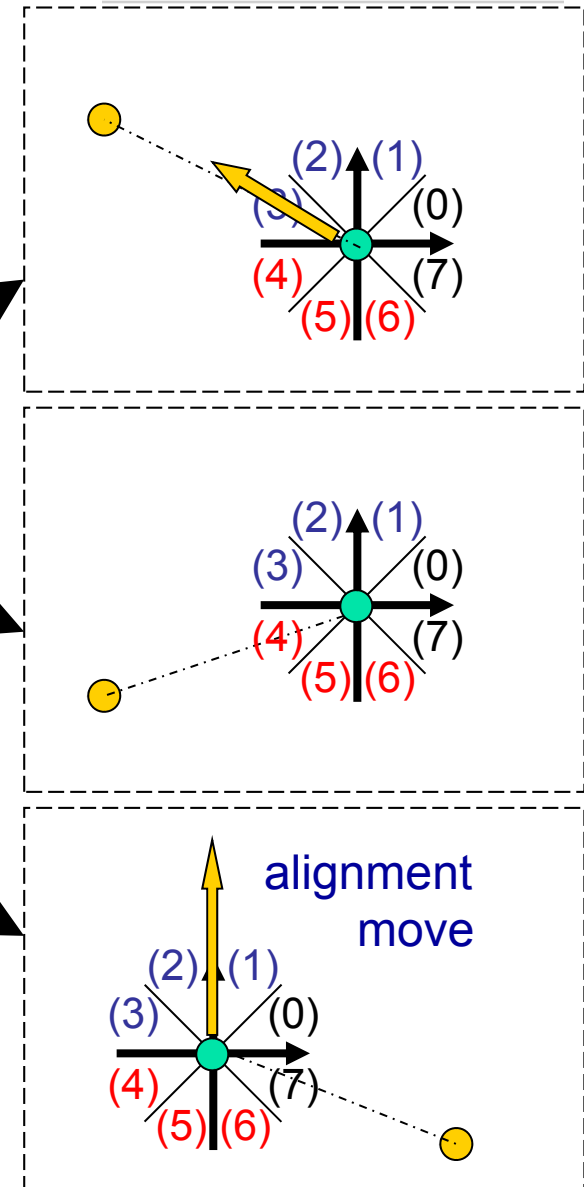
move toward the other

case: in red sectors (4), (5) or (6)

no move

case: in white sectors (7) or (0)

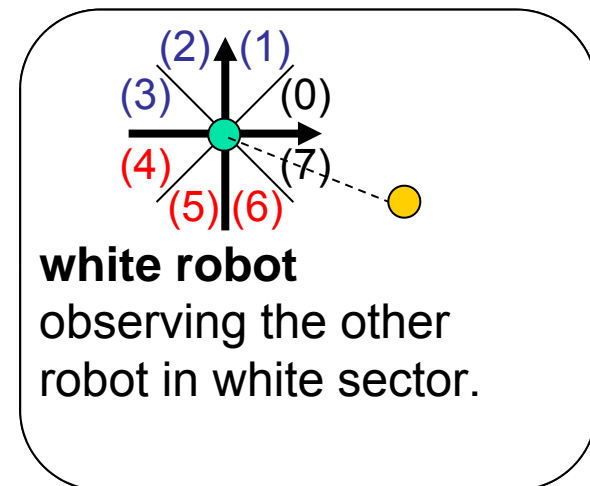
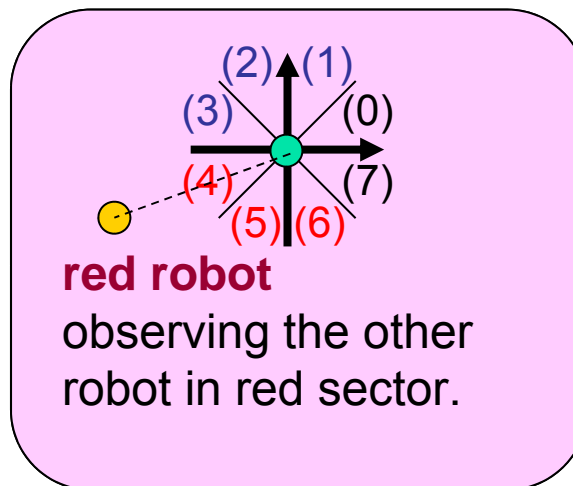
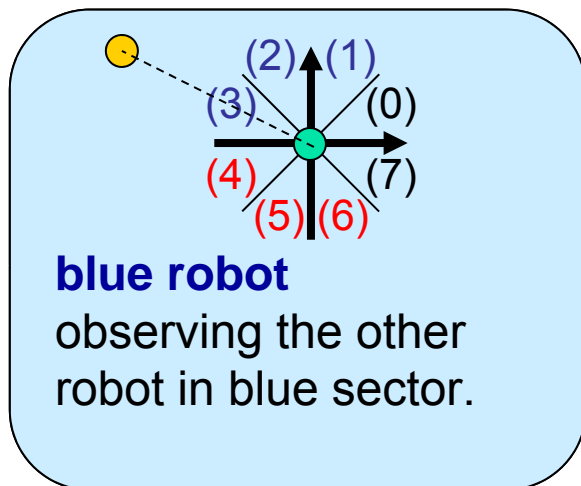
move toward a right above point where I will be able to observe the other robot in the sector (6)



# $\pi/8$ -error SDC Algorithm on SYm

## Why the robots can gather ?

To show the correctness, three names of robots are introduced:



### Dangerous Configurations

red - red : deadlock!!

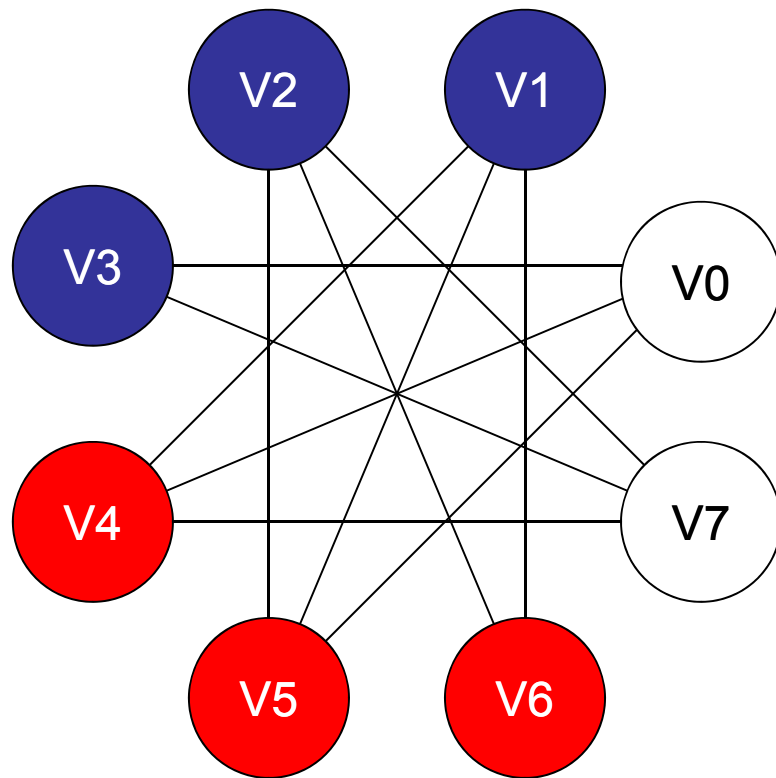
blue - blue : swapping!! (loop)

### We must show:

- dangerous configurations never occur
- blue-red configuration is eventually reached

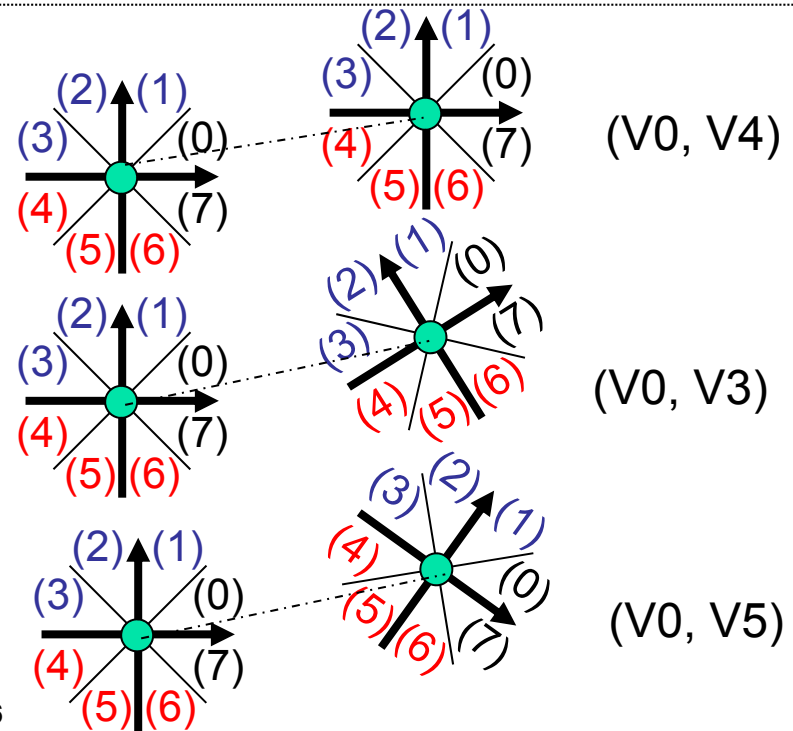
# $\pi/8$ -error SDC Algorithm on SYm

The Observation-Relation Graph



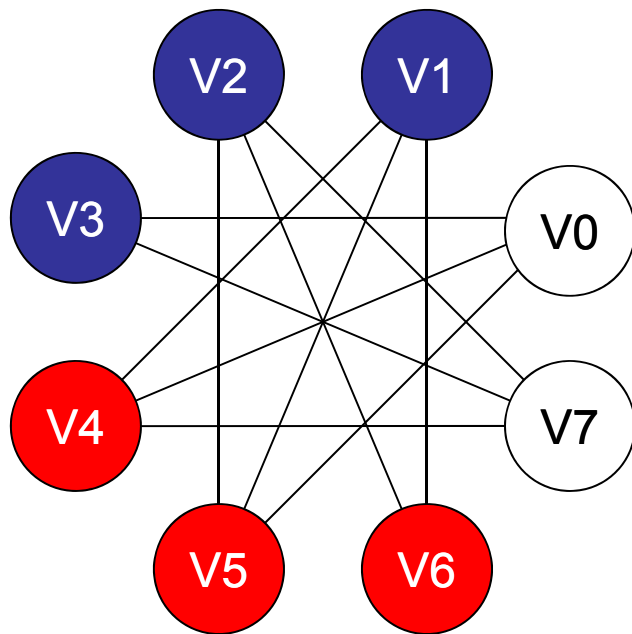
All nodes have three edges because of deviation of compass.

- $V_i$  represents a robot who observes the other in **sector (i)**.
- An edge  $(V_i, V_j)$  represents that a configuration can exist such that robots observe each other in **sector (i)** and **(j)**, respectively.



# $\pi/8$ -error SDC Algorithm on SYm

- Dangerous Configurations never occur
  - From the observation-relation graph with our sectoring and coloring, we know “red-red / blue-blue configurations never occur through executions.”



Only  
**blue-red, blue-white, red-white**  
configurations can exist.

# $\pi/8$ -error SDC Algorithm on SYm

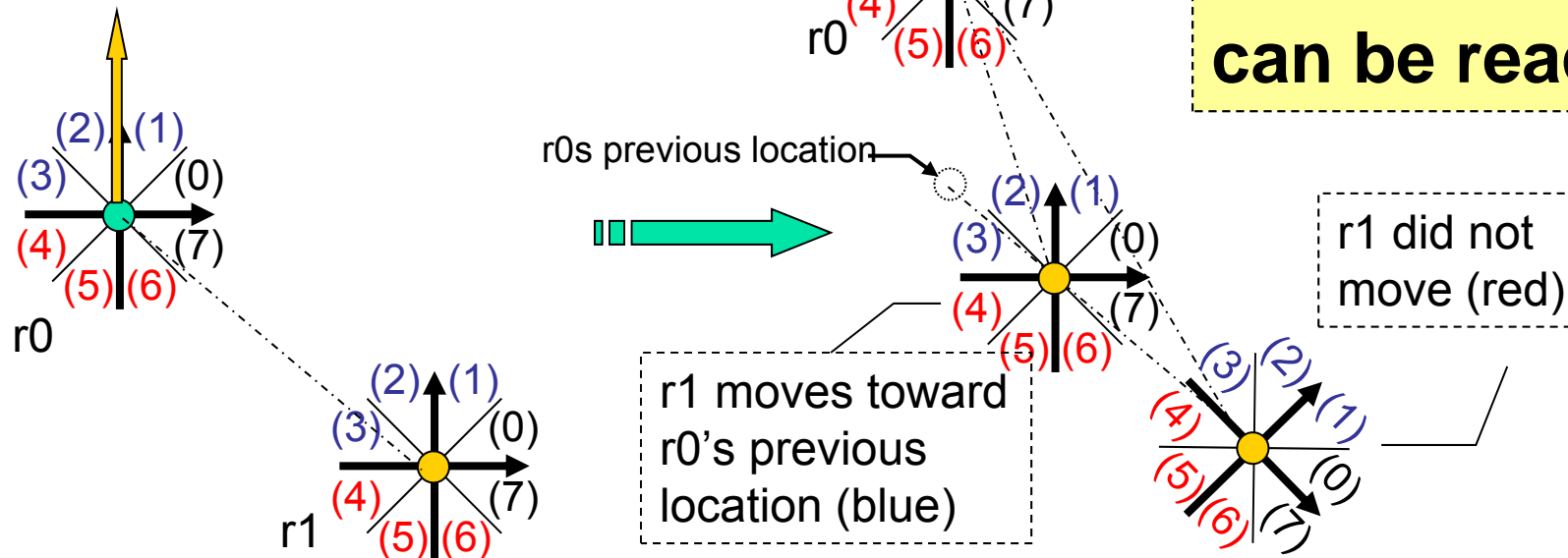
- Blue and Red configuration will be eventually reached.

- We need to show

“From blue-white/red-white configuration, if  $r_0$  moves right above,  $r_0$  can always observe  $r_1$  in

We need to consider two cases :  $r_1$  is “red” or “blue”

In both case, **blue-red conf.** can be reached.



# $\pi/8$ -error SDC Algorithm on SYm



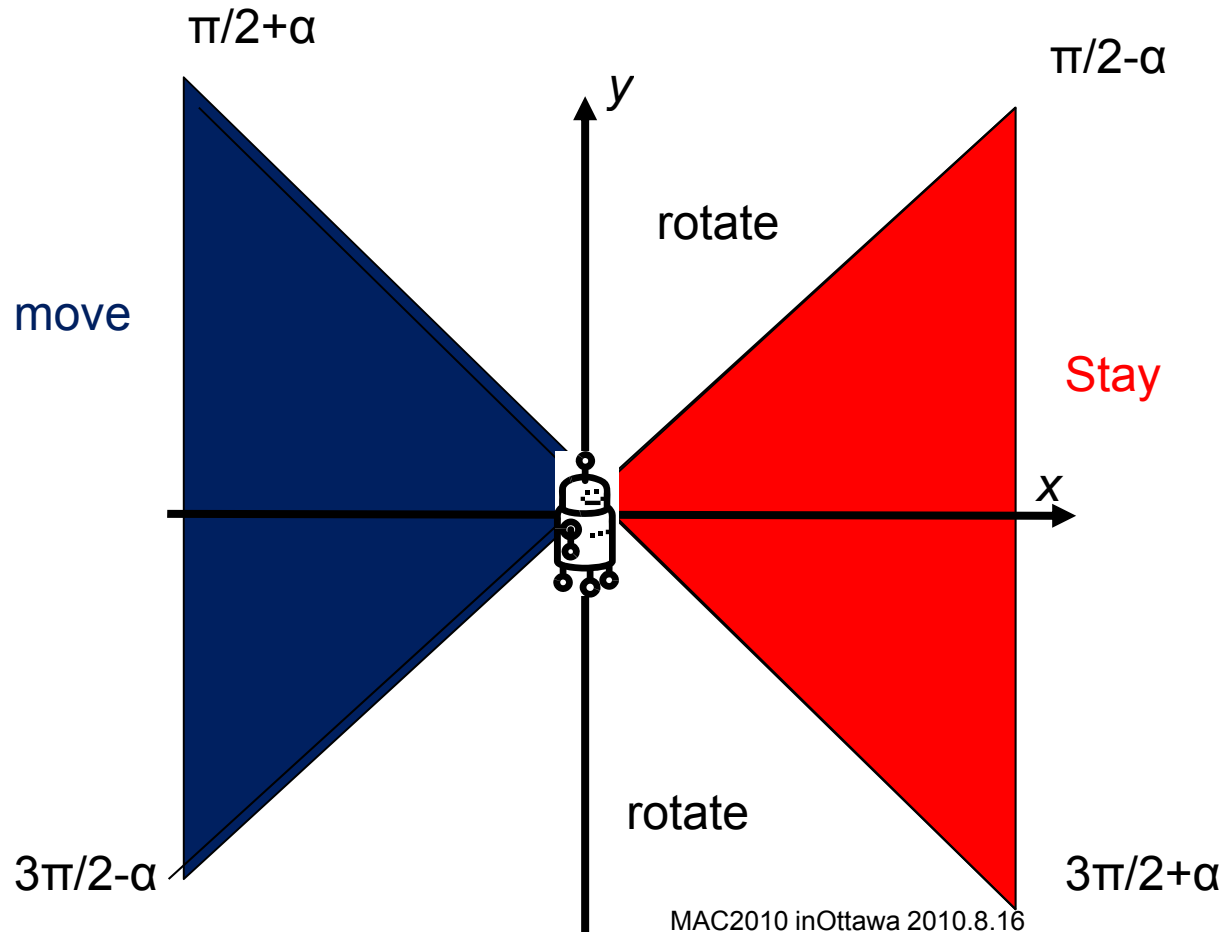
- This algorithm can behave on CORDA
- The difficulty of proof on CORDA
  - Some robot  $r_0$  observes  $r_1$ ,  $r_1$  may be moving
    - The relation when  $r_1$  stops is different from the relation when  $r_0$  observed.
- (In SYm, such situation can not occur.)
- Fast robot and very slow robot
  - Most problems do not occur for 2 robots



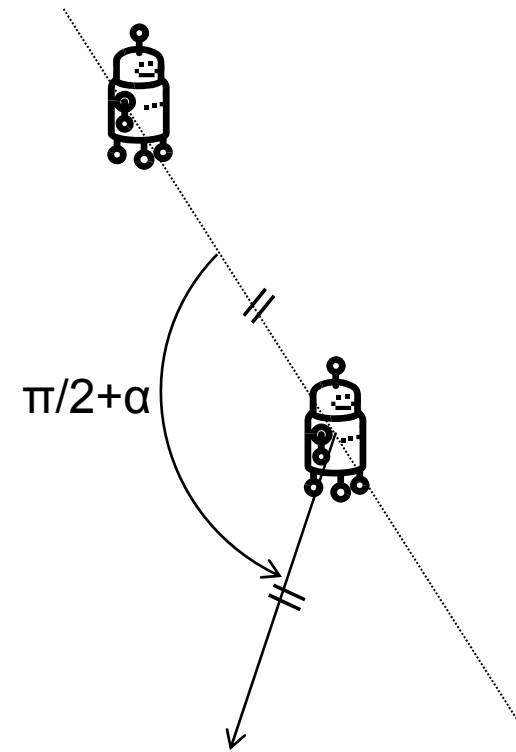


# $\alpha$ -error SDC Algorithm on SYm

$\alpha < \pi/4$

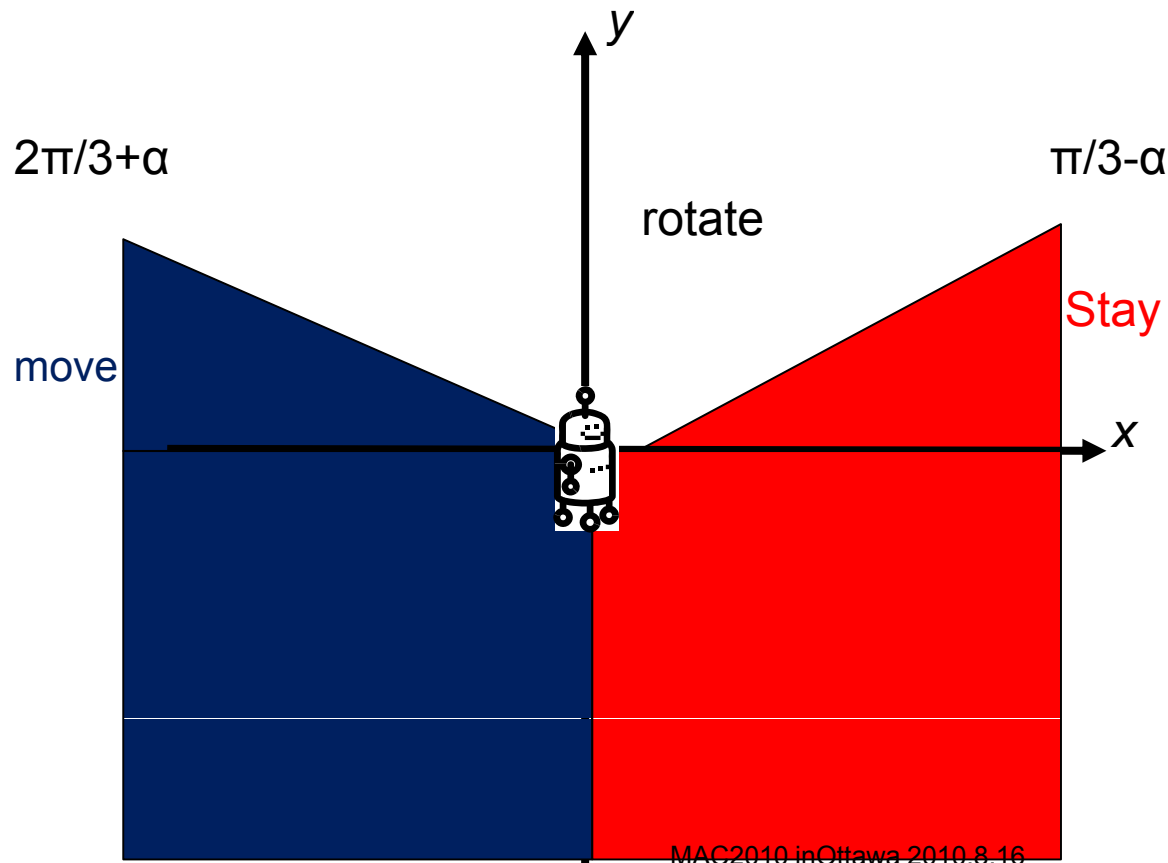


rotate

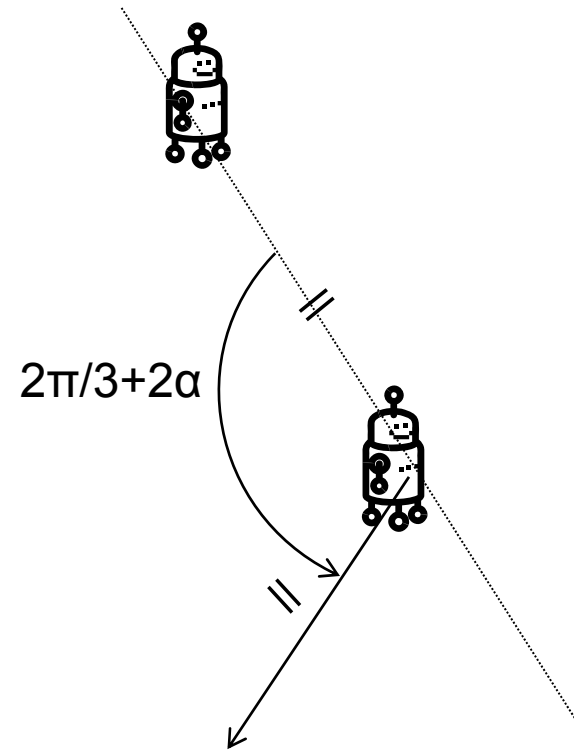


# $\alpha$ -error SDC Algorithm on CORDA

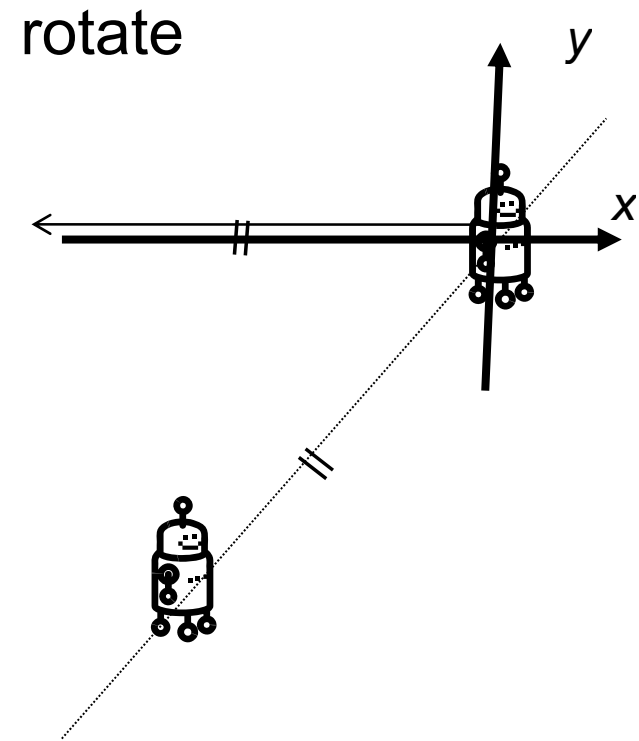
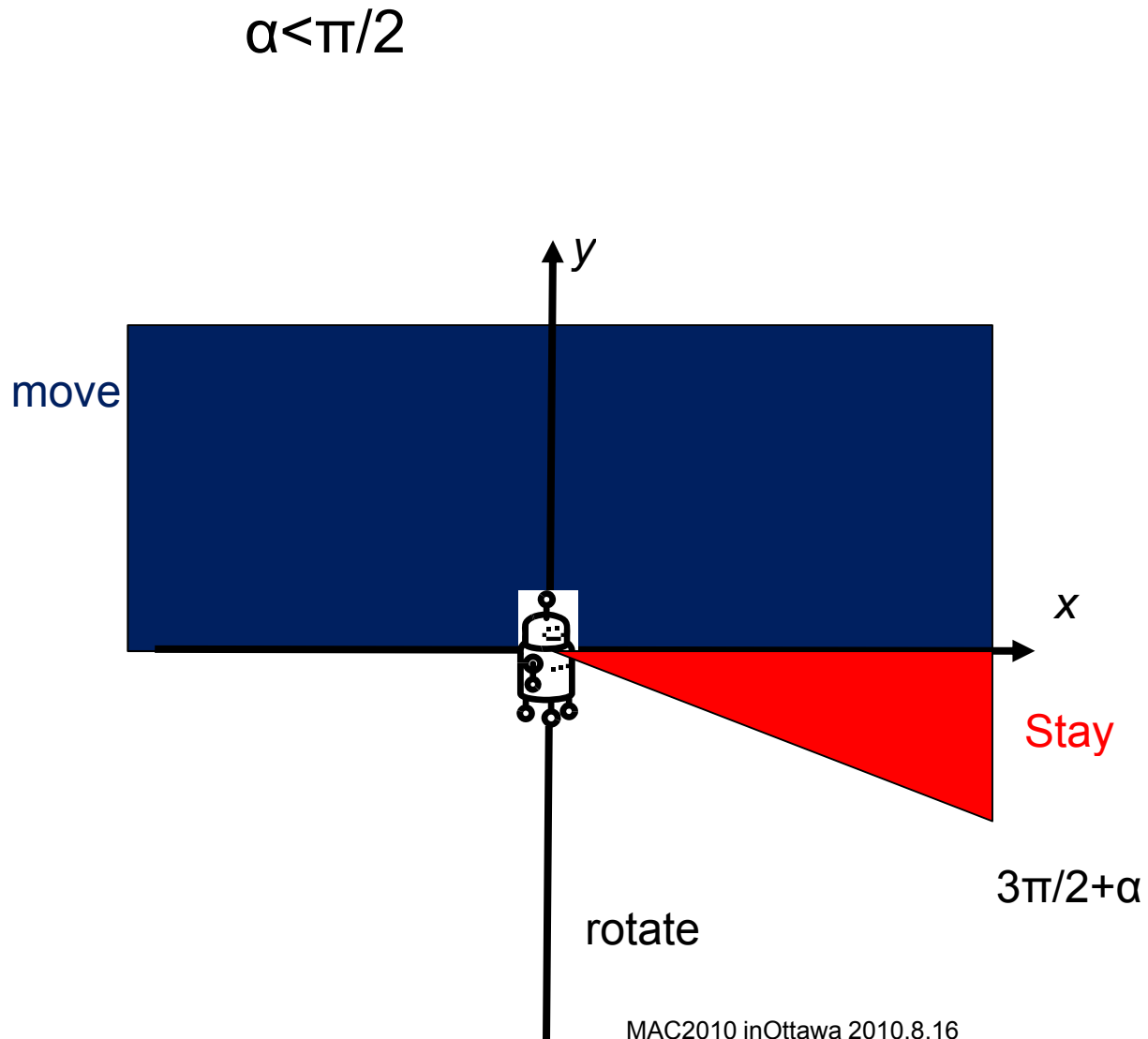
$\alpha < \pi/3$



rotate



# $\alpha$ -error FXC Algorithm on CORDA(SYm)



# Conclusions

## Two-robot Gathering problem on $\alpha$ -error compass

	SYm	CORDA
Semi-DC	impossible( $\alpha=\pi/4$ ) possible( $\alpha<\pi/4$ )	open possible( $\alpha<\pi/6$ )
FiXedC	impossible( $\alpha=\pi/2$ )[1] →	← possible( $\alpha<\pi/2$ )

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# Open Problems

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- Angle gap of SDC on CORDA
  - Impossible for  $\alpha < \pi/6$  on CORDA
  - Possible for  $\pi/4 > \alpha > \pi/6$  on CORDA
- Extension to n-robot system
  - SDC( $\alpha < \pi/4$ ) on SYm is possible [DISC2007]

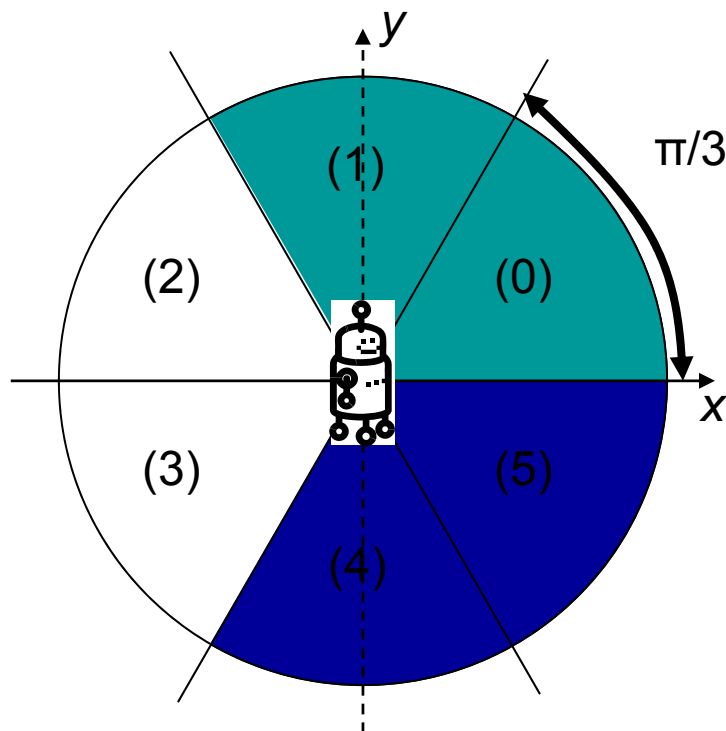


Thank you!

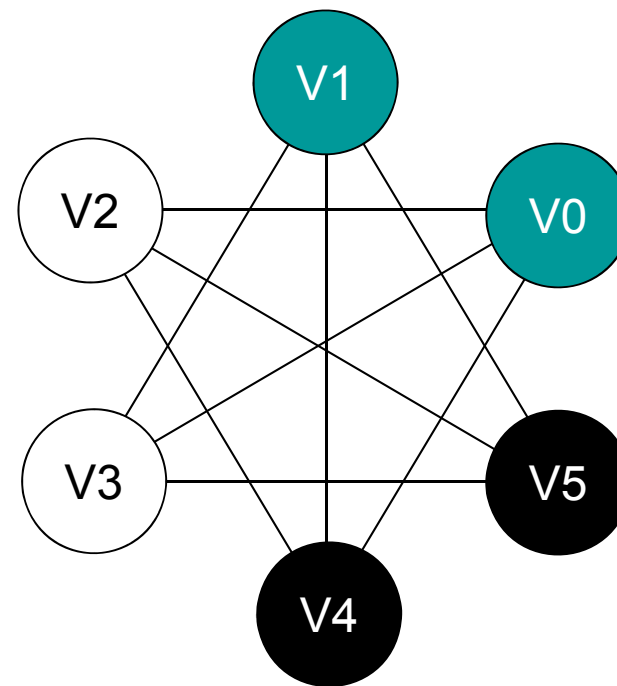
# $\pi/3$ -relative error FXC Algorithm

Basic idea is same with  $\pi/4$ -absolute error SDC algorithm

Dividing the world into 6 sectors

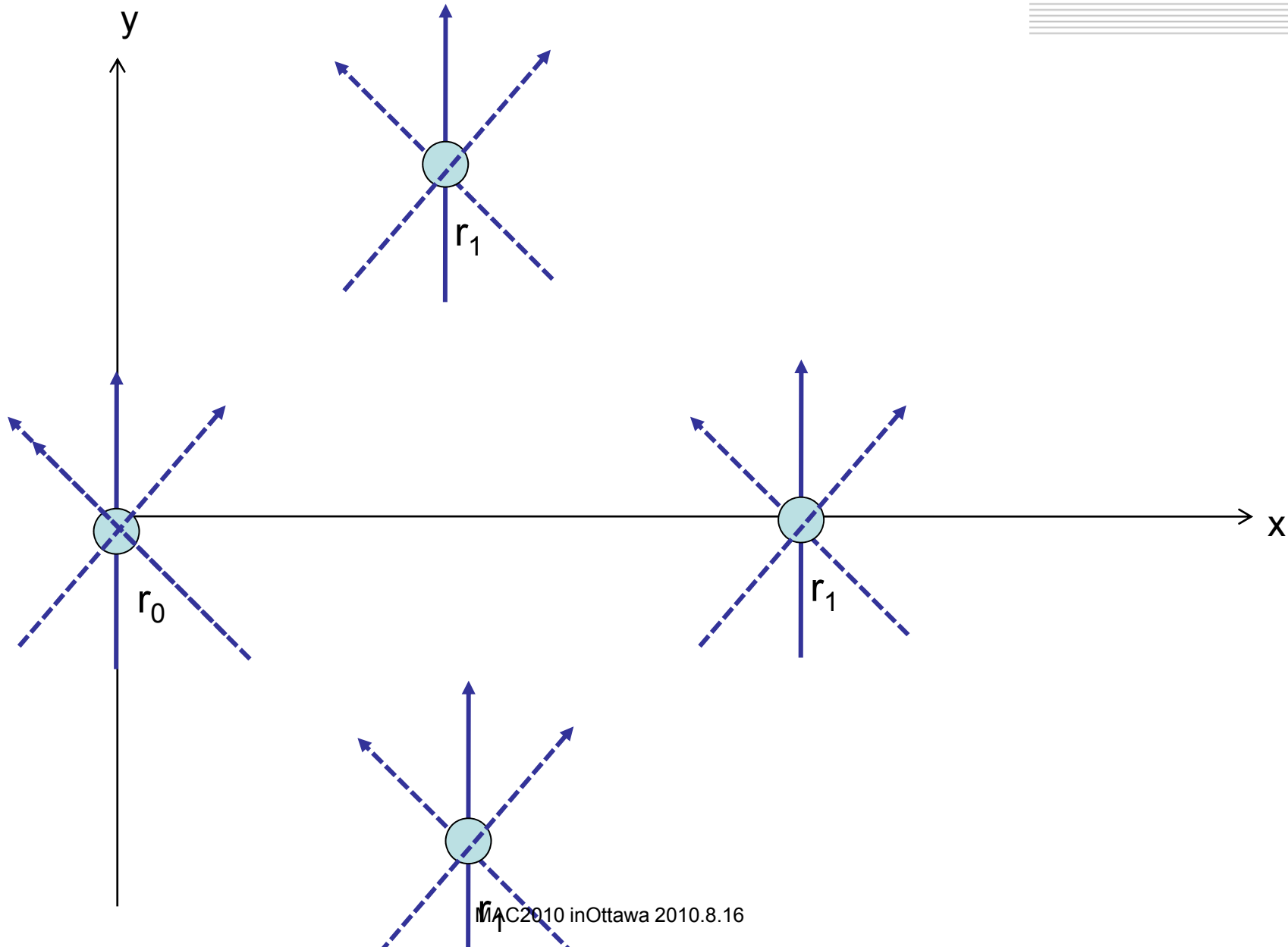


The observation-relation graph



no red-red, blue-blue

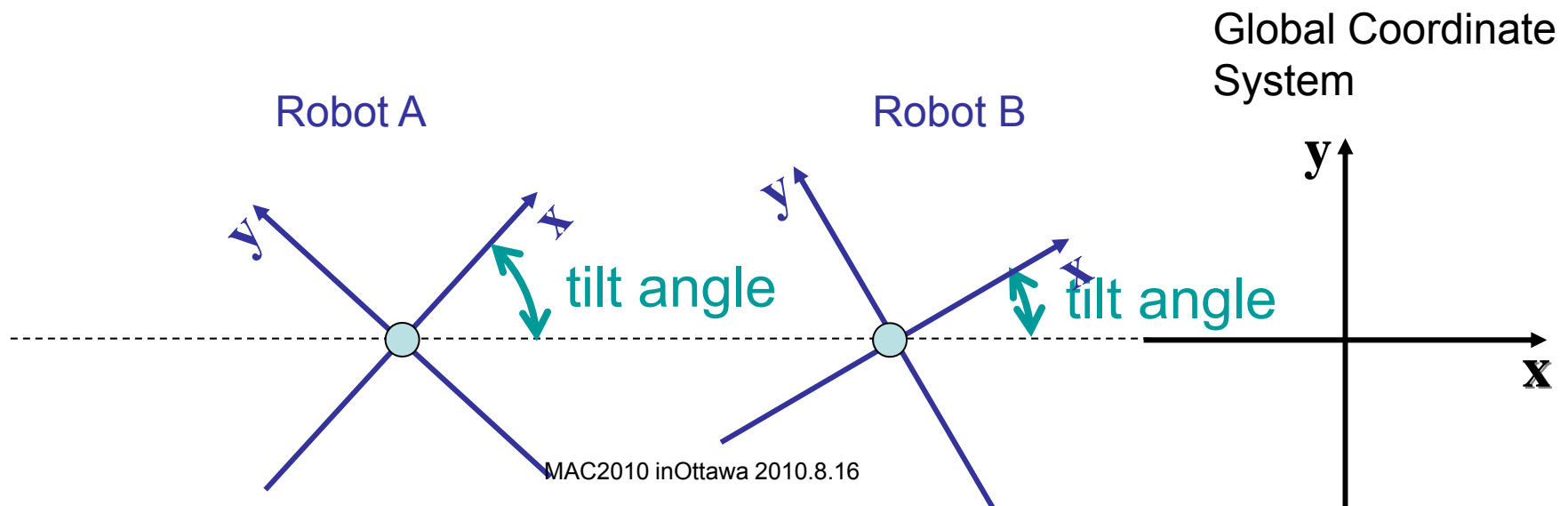
# Impossibility( $\pi/4$ -error compass , SDC and Sym)





# Weakly-agreed compass

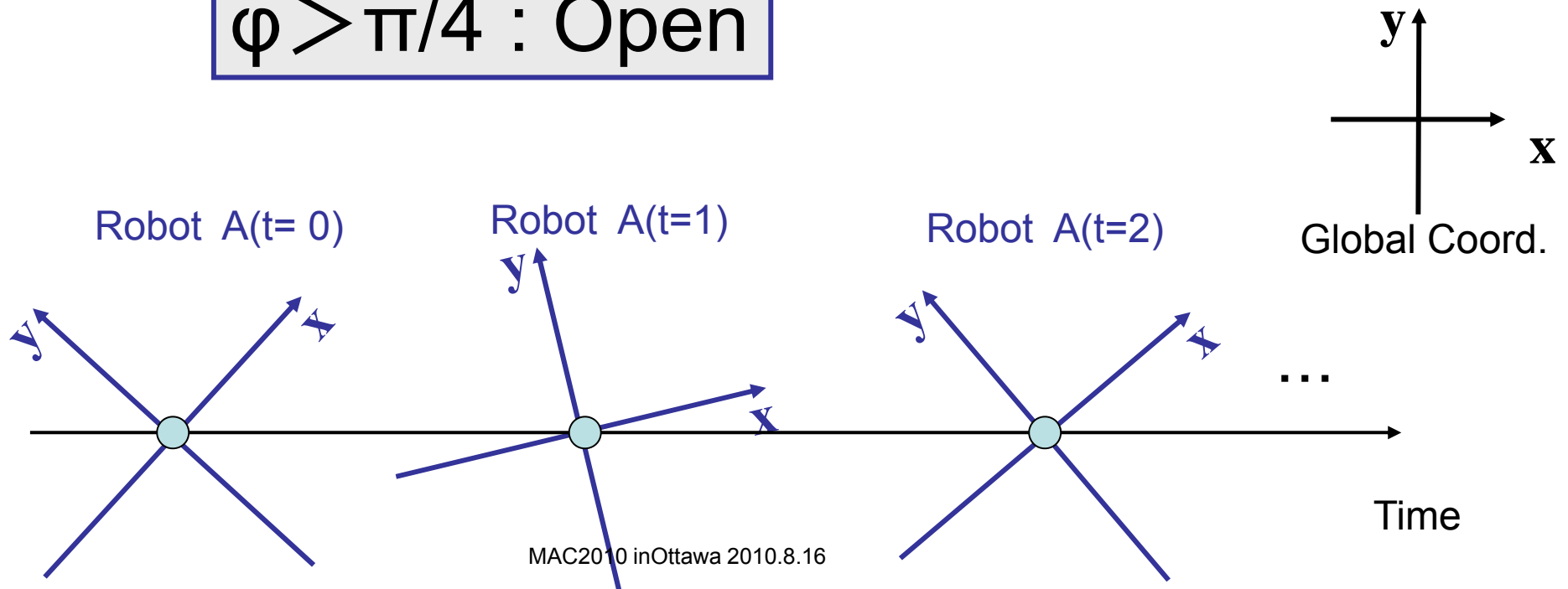
- Measuring Compass Agreement Level by tilt angle [Imazu et al., 05][Souissi et al., 06]
  - Tilt angle =  $\angle$  formed by the global and local axis
- Tilt angle of every robot  $< \pi$  : Solvable [Yamashita et al., 07]
  - Asynchronous / #robots = 2



# Dynamic Compass

- Tilt Angle varies with time [Katayama et al., 07]
  - At the beginning of each cycle
  - Bounded by  $\varphi$
- $\varphi \leq \pi/4$  : Solvable [Katayama et al., 07]
  - Asynchronous / #robots = 2

$\varphi > \pi/4$  : Open



# Our Contribution

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- Dynamic compass with  $\varphi \leq \pi/2 - \varepsilon$ : Solvable

- $\varepsilon (>0)$ : Arbitrary small constant

- Semi-synchronous / #robots  $n$  is arbitrary

(The first result considering any #robots with disagreed compasses)

- $\varphi \geq \pi/2$ : Unsolvable

- Semi-synchronous / #robots = 2

➡ Our Result is optimal in terms of maximum tilt angle

# Algorithm Design

- Algorithm for 2 robots with  $\varphi = \pi/2 - \varepsilon$



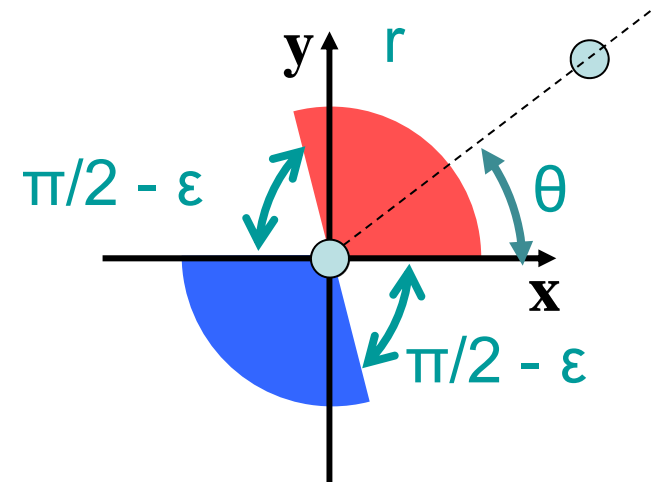
- "Conditional" Algorithm for n robots
  - Working correctly if the initial configuration has a unique **Longest Distance Segment(LDS)**
- LDS election algorithm
  - Starting any configuration, terminate a configuration with unique LDS

Composition

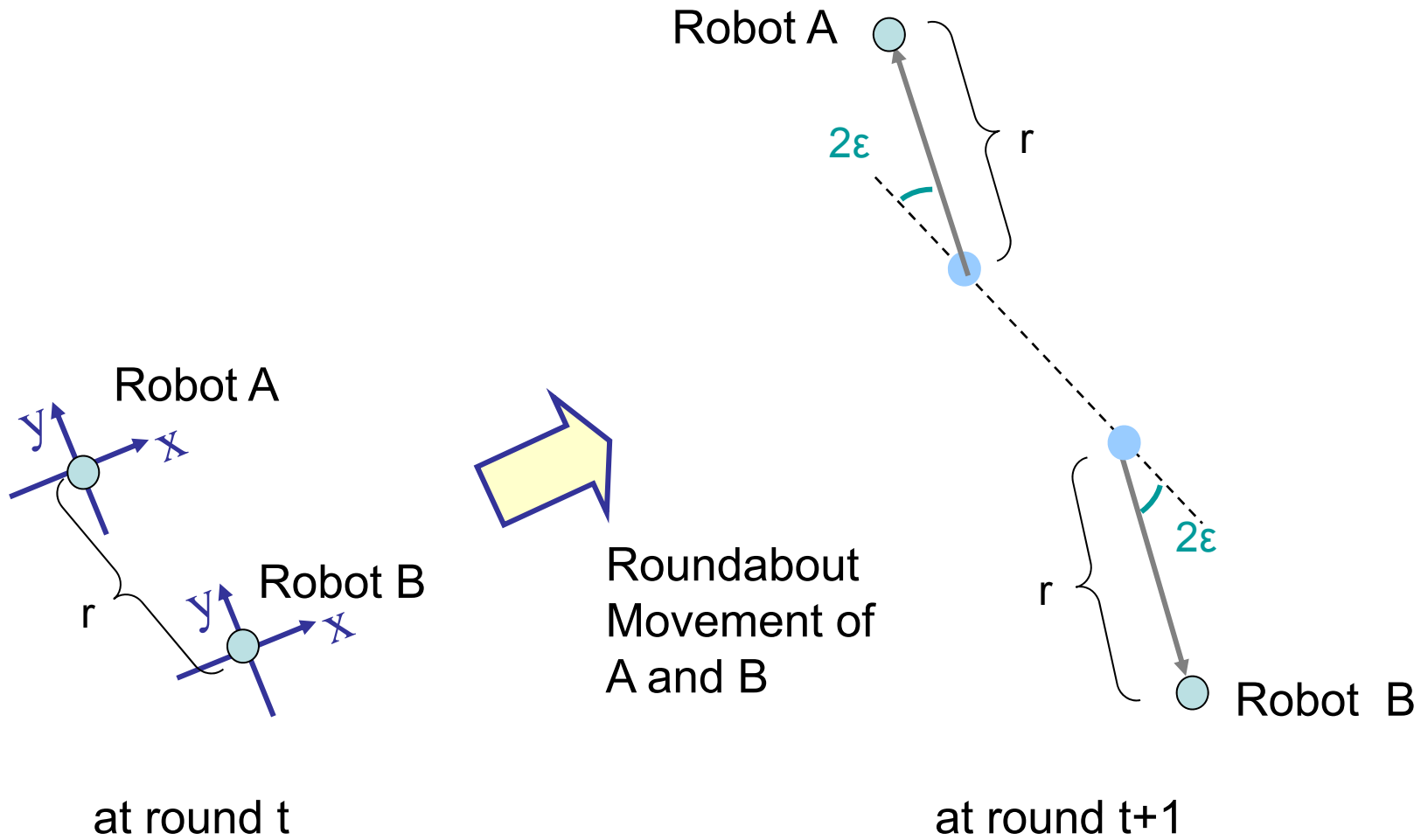
**n-robot gathering algorithm**

# 2-robot Algorithm

- Consists of three types of movement
  - $\theta$  : the angle at which a robot sees its partner
  - $r$  : distance between two robots  
(in terms of observer's local coordinate sys.)
  - $0 \leq \theta < \pi/2 + \epsilon$  : Wait
    - No movement
  - $\pi \leq \theta < 3\pi/2 + \epsilon$  : Approach
    - Move to the partner's location
  - Otherwise : Roundabout
    - Move toward the angle  $\theta + \pi - 2\epsilon$  with distance  $r$



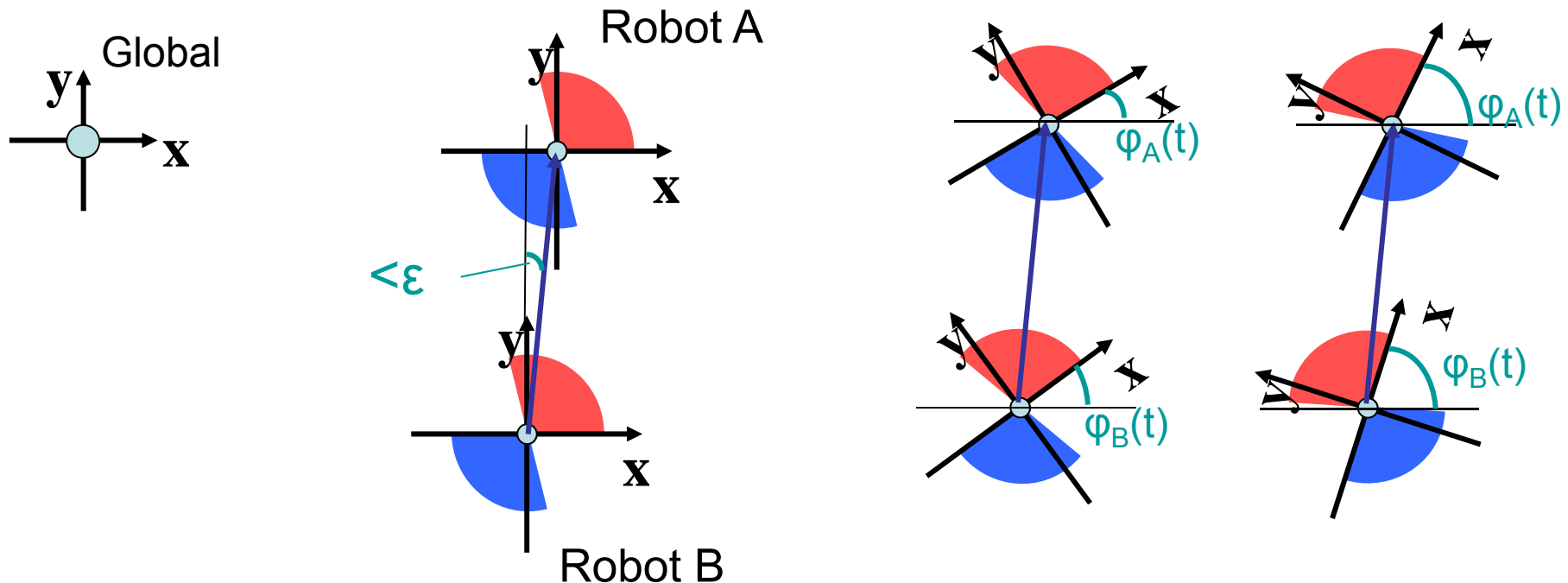
# Roundabout Movement



# Correctness (1/5)

## ■ Lemma 1

- $\angle ABy = \angle$  formed by AB and the global y-axis  $< \epsilon$   
→ **Wait-Approach** Relation is guaranteed  
(regardless of tilt angles of robot A and B)



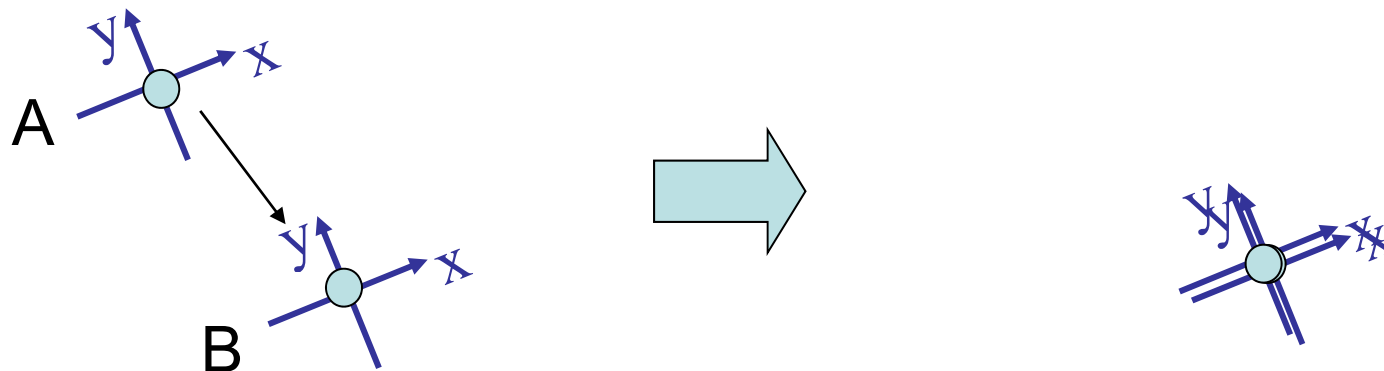
# Correctness(2/5)

## ■ Lemma 2

- At any round,  $\angle ABy$  decreases by  $\epsilon \sim 2\epsilon$  unless gathering is achieved

■ A: Approach move

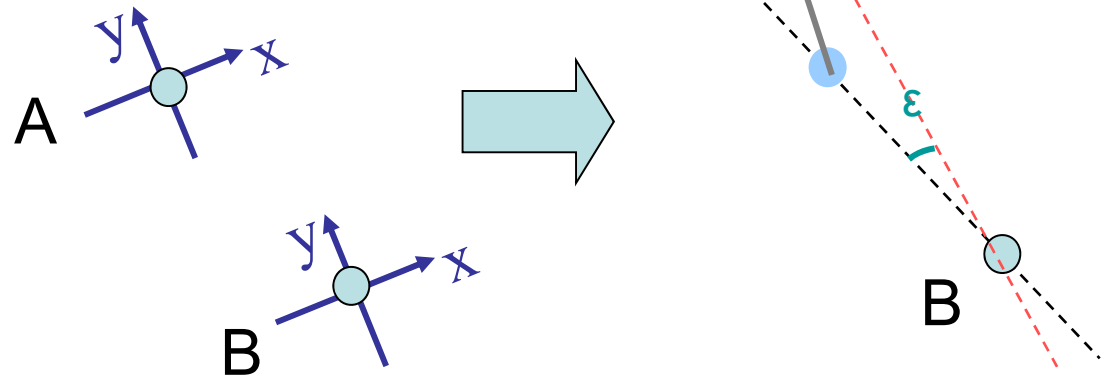
■ B: No movement (Wait or inactive)



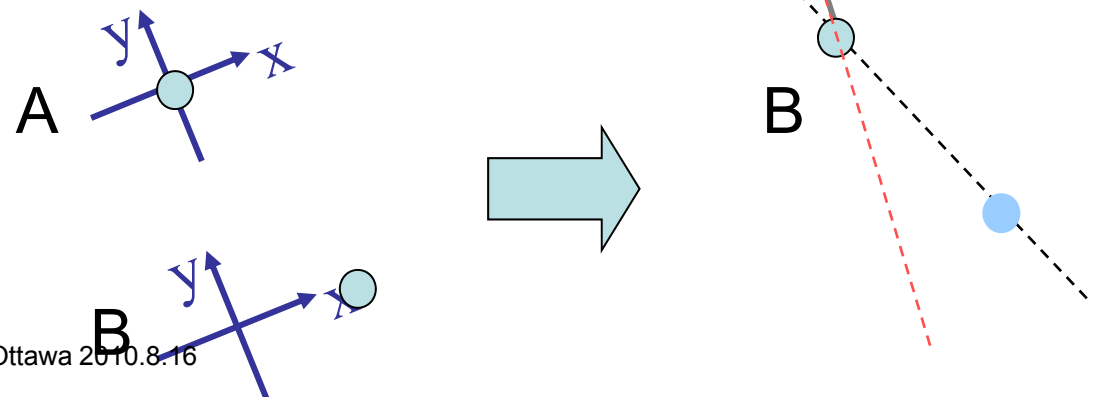


# Correctness(3/5)

- A: Roundabout Move  
B: No movement  
→ decrease by  $\epsilon$

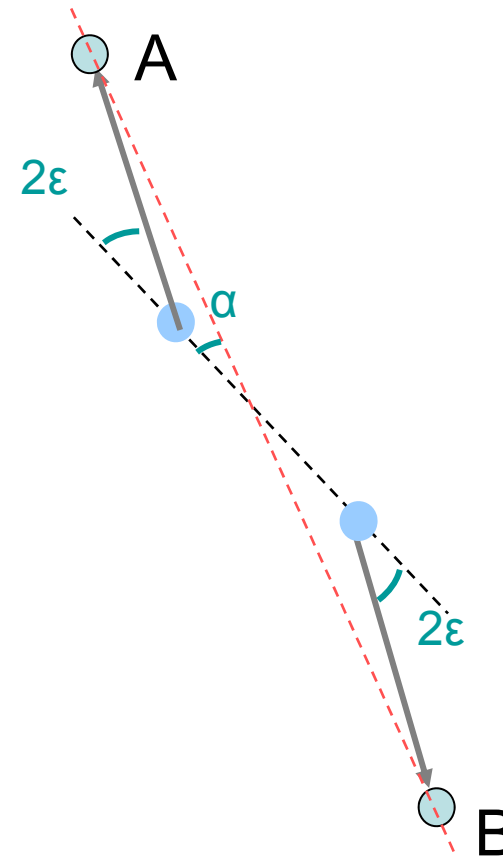
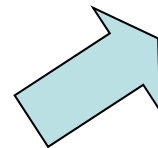
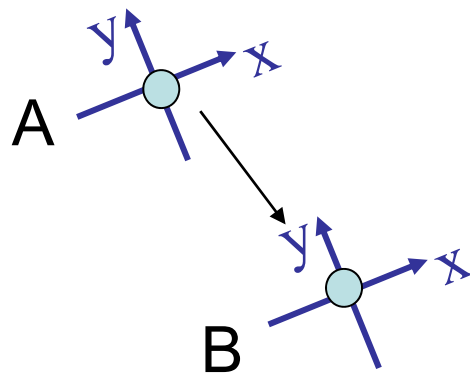


- A: Roundabout Move  
B: Approach Move  
→ decrease by  $2\epsilon$



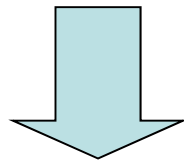
# Correctness (4/5)

- A: Roundabout move
- B: Roundabout move
- decrease by  $\varepsilon \sim 2\varepsilon$



## Correctness (5/5)

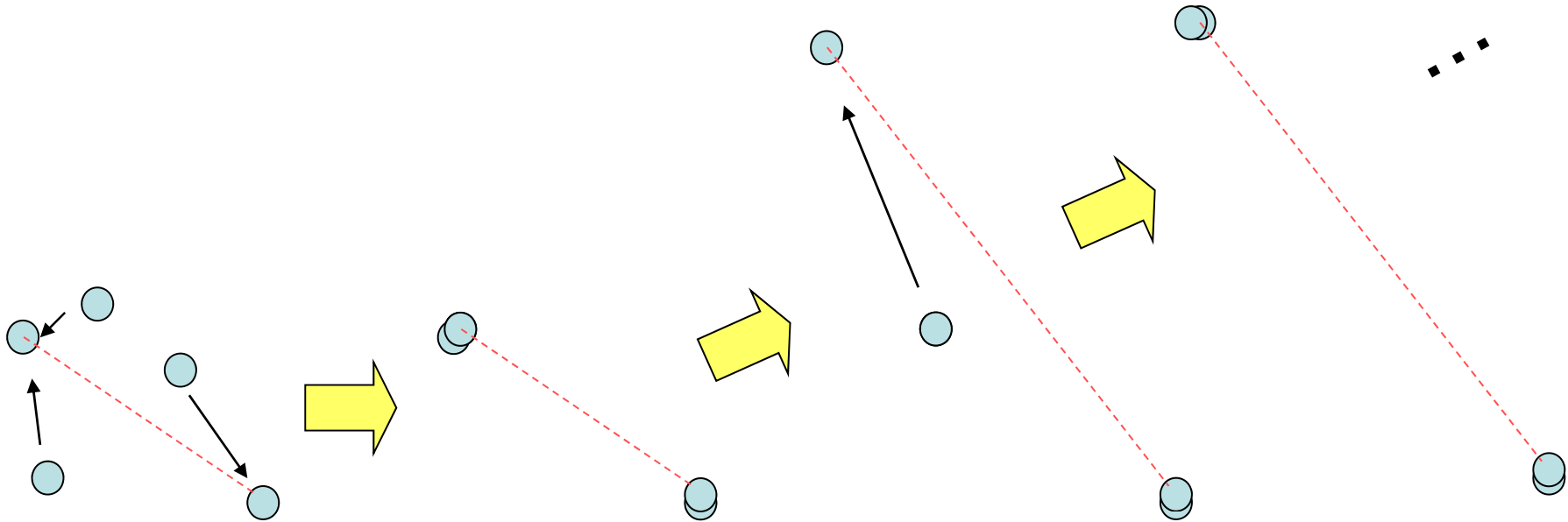
- From Lemma 2,  
$$-\varepsilon \leq \angle ABy < +\varepsilon \text{ eventually holds}$$
- From Lemma 1,  
If  $-\varepsilon \leq \angle ABy < +\varepsilon$  holds, one robot approaches and the other waits.



**Gathered !**

# n-robot algorithm under unique LDS(1/2)

- Robots are located at two points  
→ All robots execute the two-robot algorithm
- Robots are located at more than two points  
→ All robots move to one of two endpoints of LDS



# Correctness of Conditional n-robot Alg.



## ■ Lemma 3

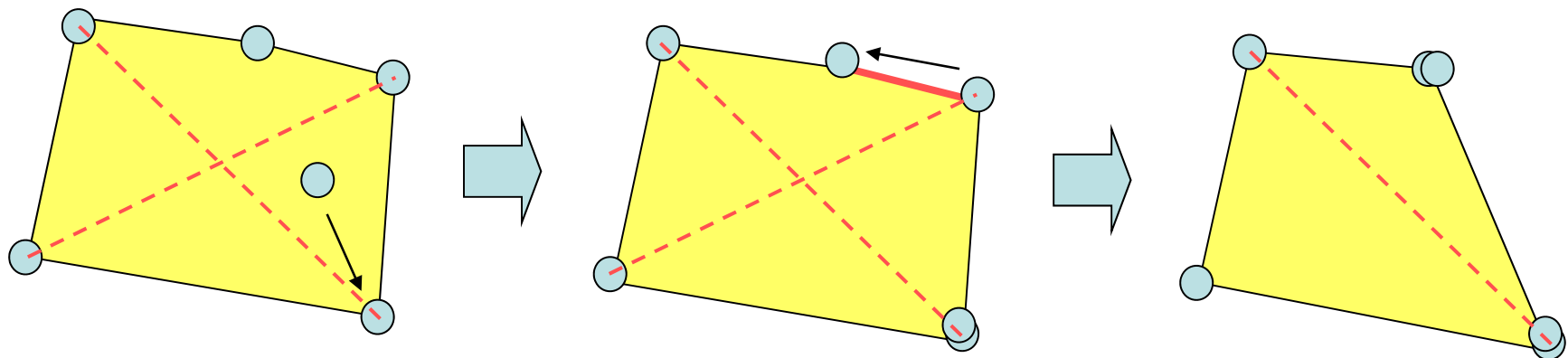
- $\angle \text{LDSy} = \angle$  formed by LDS and the global y-axis  $< \varepsilon$   
→ **Wait-Approach** Relation is guaranteed  
(regardless of the title angle of each robots)

## ■ Lemma 4

- At any round,  $\angle \text{LDSy}$  decreases by  $\varepsilon \sim 2\varepsilon$   
unless gathering is achieved

# Unique LDS Election (1/2)

- If two or more LDSs exist, each robot calculates the convex hull(CH)
  - Robots on the boundary : Wait
  - Inner robots : Moves to one of vertices
- Contracting the shortest edge of the CH

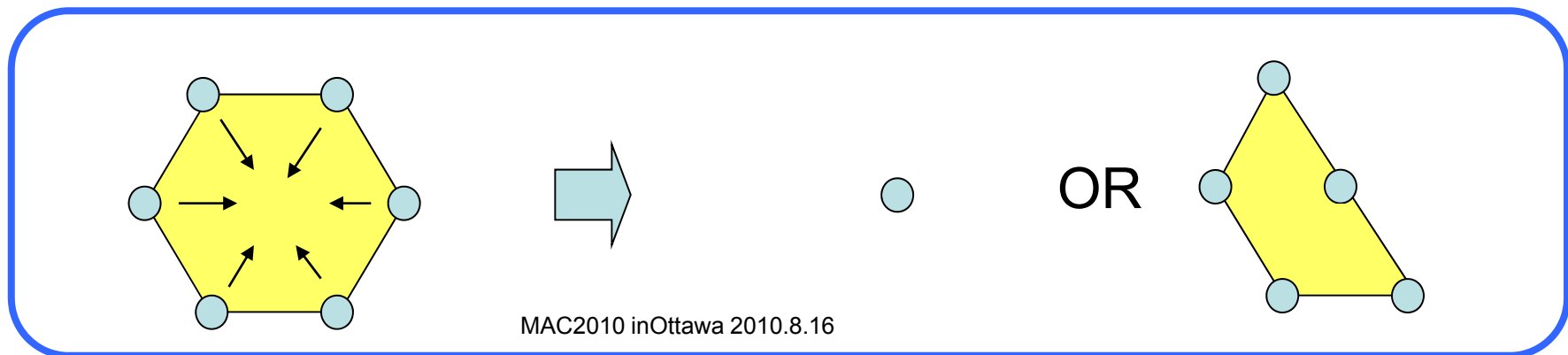
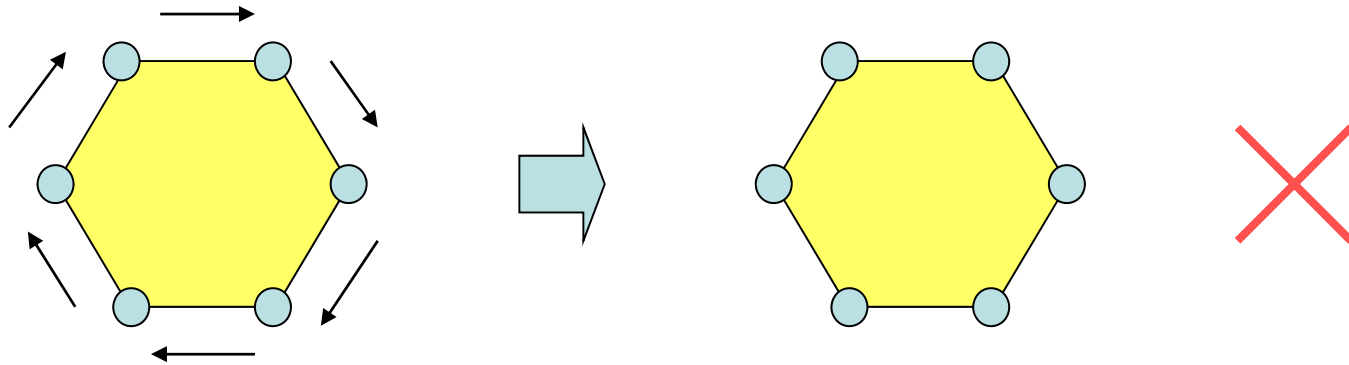


#edges of the CH decreases

→ Eventually unique LDS is elected (or gathered)

# Unique LDS Election(2/2)

- If all edges have a same length
  - Robots moves to the center-of-gravity of the CH
    - All robots simultaneously move → gathered
    - A part of robots move → Symmetry is broken



# Conclusion

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- Gathering mobile robots with dynamic compasses
  - Tilt angle  $\leq \pi/2 - \varepsilon$  (**Optimal**)
  - Semi-synchronous model
  - Arbitrary #robots
- Open problem
  - Asynchronous model
    - $\pi/2 < \text{Maximum Tilt angle} < \pi/4$
    - Recently, two robots are solved for  $< \pi/3$
    - #robots = 2, dynamic compass