

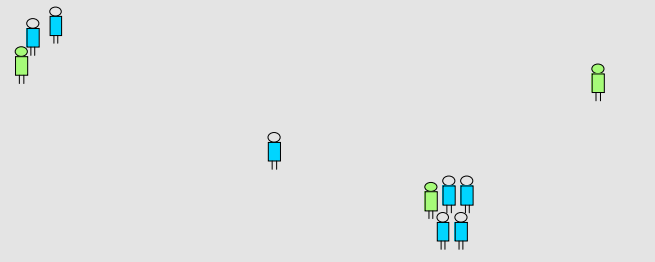
DISTRIBUTED COMPUTATIONS

BY

MOBILE ENTITIES

Nicola Santoro
Carleton University

COMPUTATIONAL MOBILE ENTITIES

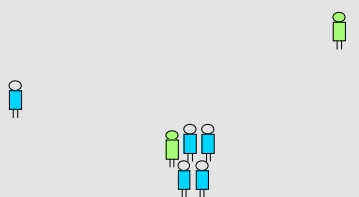


A diagram showing several small human-like figures in various colors (green, blue, cyan) scattered across a light gray background. Some are grouped together, while others are isolated.

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CONTINUOUS : 2D/3D SPACE - TERRAIN -

ROBOTS

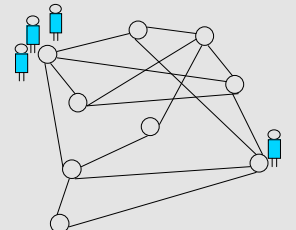


A diagram showing several small human-like figures in various colors (green, blue, cyan) scattered across a light gray background. Some are grouped together, while others are isolated.

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DISCRETE : NETSCAPE - GRAPH WORLD -

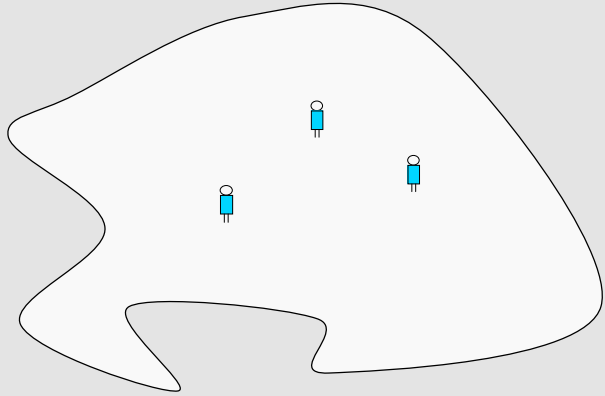
AGENTS



A diagram showing a network of nodes connected by lines, representing a graph world. Several small human-like figures in various colors (green, blue, cyan) are positioned at some of the nodes.

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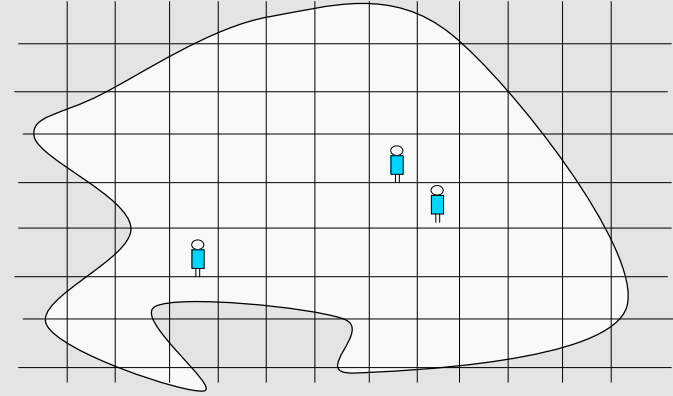
DISCRETIZED CONTINUOUS SPACE



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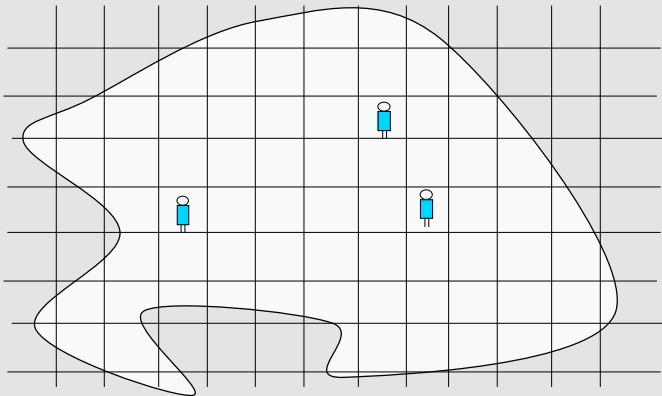
DISCRETIZED CONTINUOUS SPACE



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6

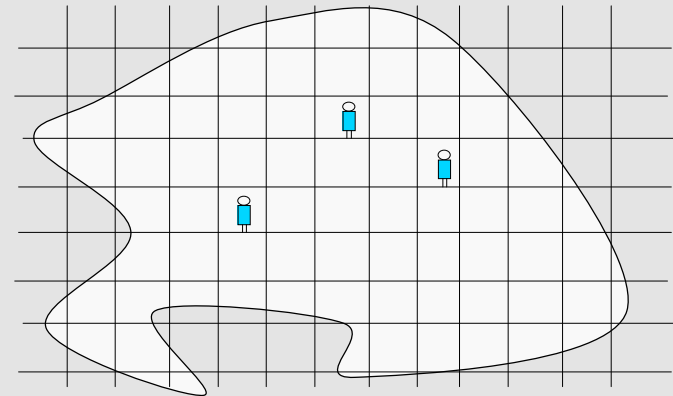
DISCRETIZED CONTINUOUS SPACE



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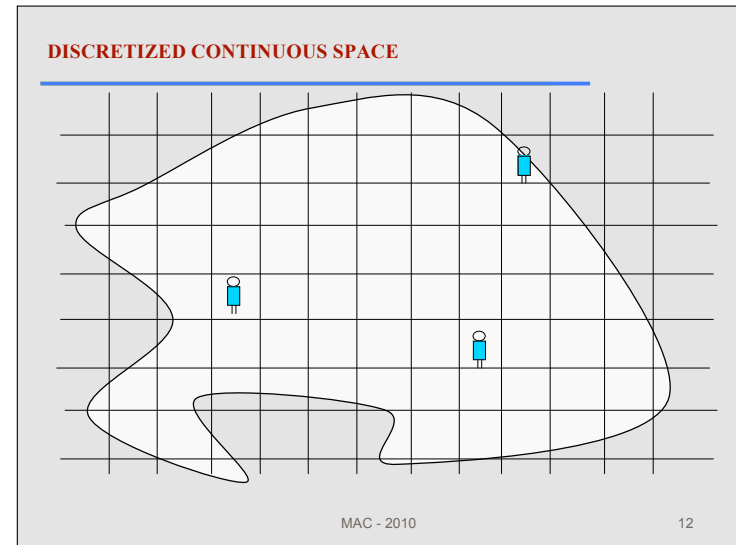
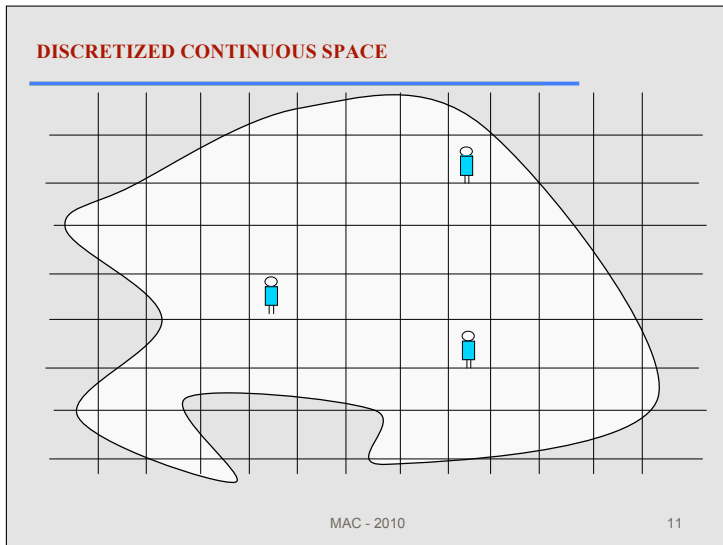
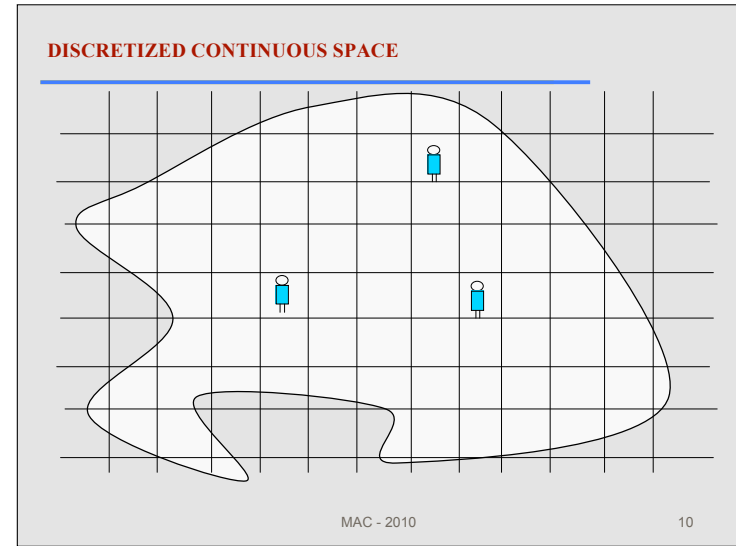
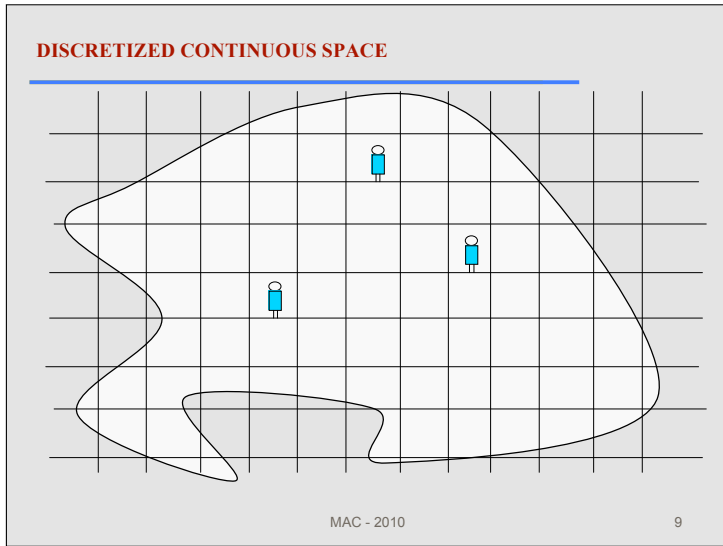
7

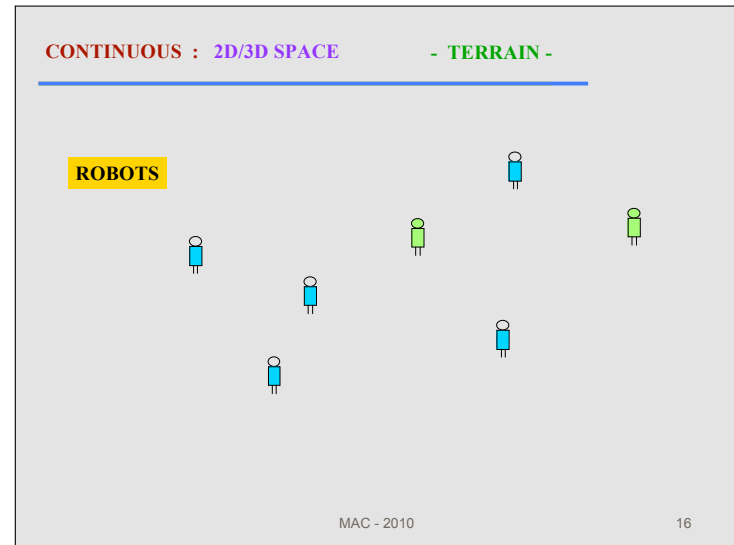
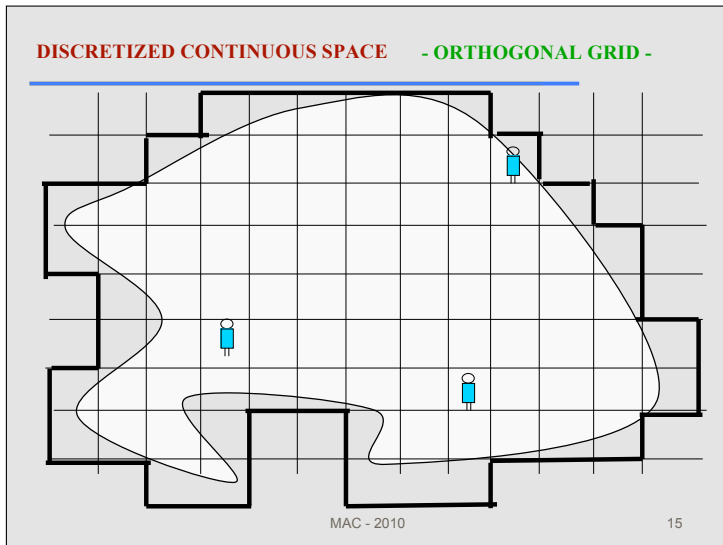
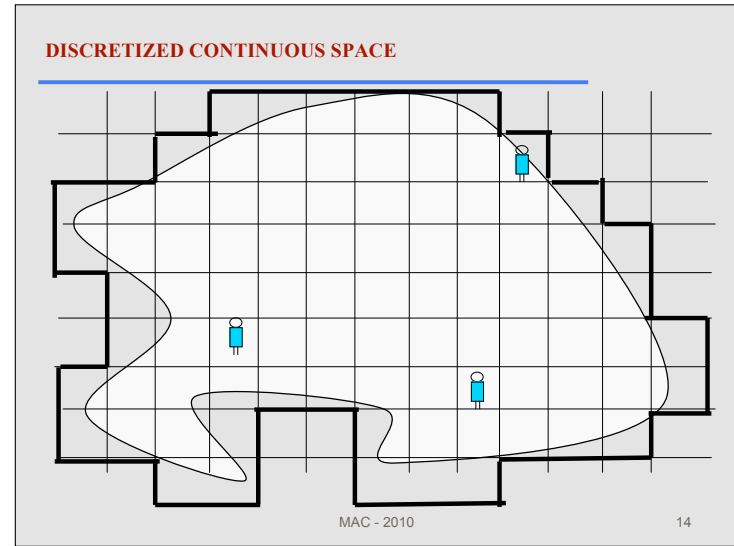
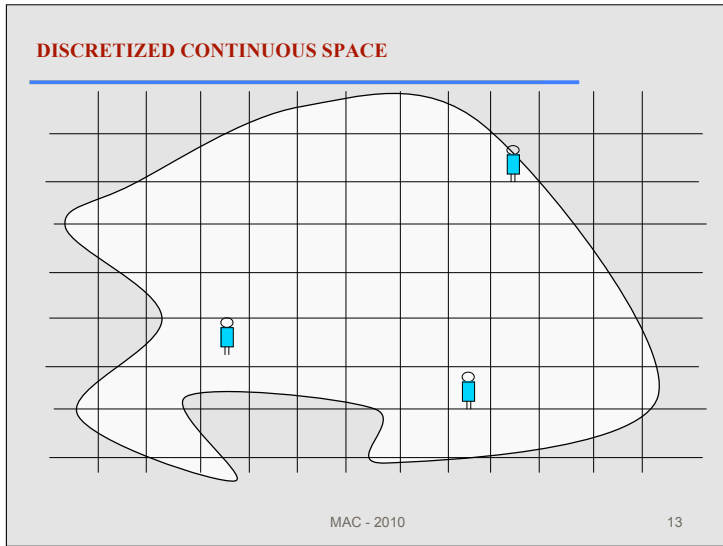
DISCRETIZED CONTINUOUS SPACE



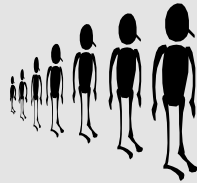
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Distributed Computing by Mobile Robots



Practical Motivations

Very Few
Complex
Specialized
Robots



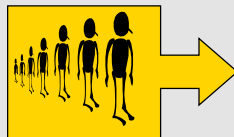
Many
Very Simple
Generic
Micro-Robots

- Easy to manufacture
- Cheap
- Easily replaceable
- Reusable

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Autonomous Mobile Robots



TASK

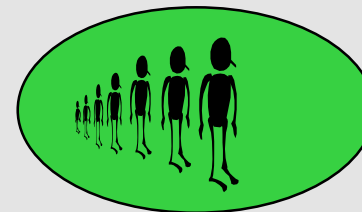
- A robot - alone - is computationally weak
- Cooperation of robots is essential to perform complex tasks

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System

Many Simple
Micro-Robots



System (SWARM)

{ Upgradable
Robust

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System

In addition to **sensing** and **computing** the **micro-robots** are capable of **locomotion**

System

Mobility adds useful system capabilities

- Patrol a wide area
- Re-positioning for better surveillance
- Work in hostile/dynamic environments
 - Risky area surrounding or surveillance
 - Search and rescue missions
 - Space exploration
 - Military operations

Cooperation of autonomous mobile robots

Robotics
Artificial Intelligence
Behaviour of Social Animals (Insects)
Swarms and Ants
Control
Distributed Algorithms

Main Research Question

Minimum level of capabilities
(sensorial, computational, and motorial)
robots must possess to **COLLECTIVELY**
solve a given task **deterministically**

Main Research Question

Minimum level of capabilities
(sensorial, computational, and motorial)
robots must possess to **COLLECTIVELY**
solve a given task **deterministically**

How **weak** can each single robot and still
globally accomplish the given task ?

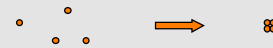
How much **local power** is necessary to perform
global computations ?

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Basic Coordination Tasks

Gathering



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Basic Coordination Tasks

Gathering



Specific Patterns



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Basic Coordination Tasks

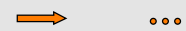
Gathering



Specific Patterns



Alignment

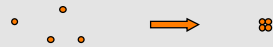


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Basic Coordination Tasks

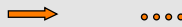
Gathering



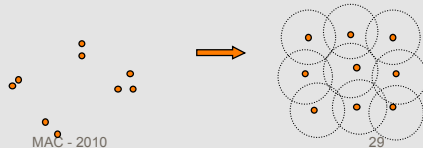
Specific Patterns



Alignment



Scattering



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Complex Tasks

Mine Sweeping

Hazardous Retrieval

Rescue Operation

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COMPUTATIONAL MODEL

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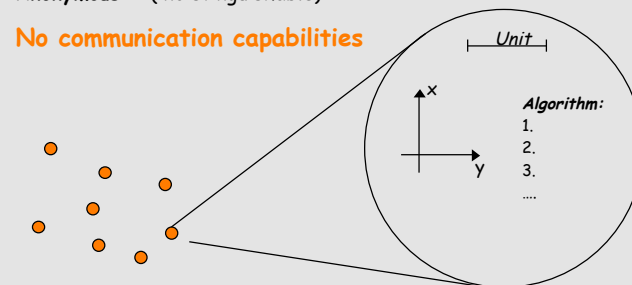
Weak Mobile Robots

Autonomous (no central control)

Homogeneous (run same software)

Anonymous (indistinguishable)

No communication capabilities



The robots are points.

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Robot's Basic Capabilities

•Processing and Storage

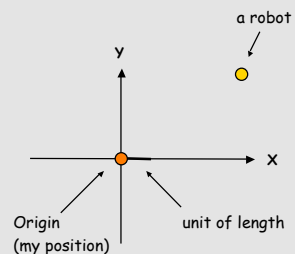
- limited
- execute **same protocol**

•Sensorial

- "see" other robots
- local coordinate system

•Motorial

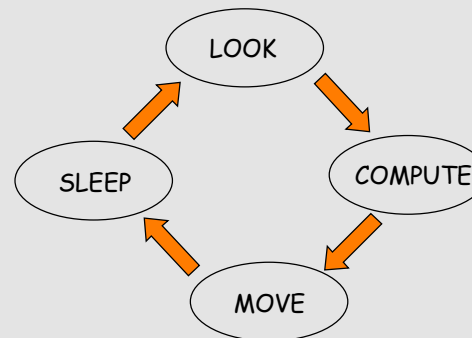
- move towards a destination



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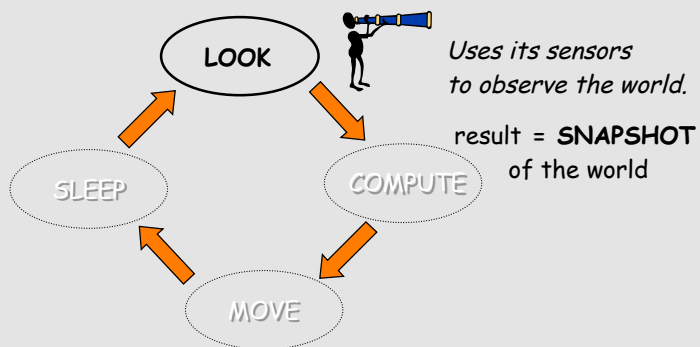
Robot's behaviour : Life Cycle



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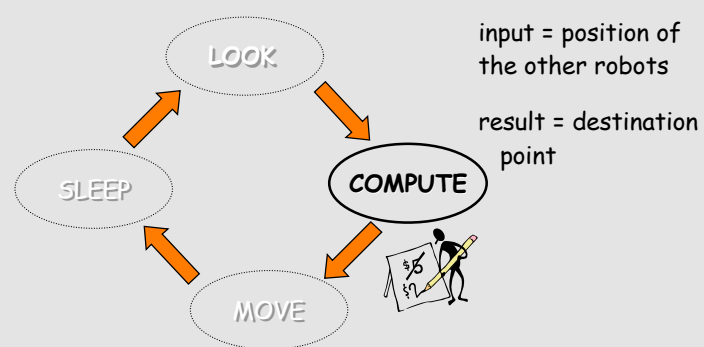
Robot's behaviour : Life Cycle



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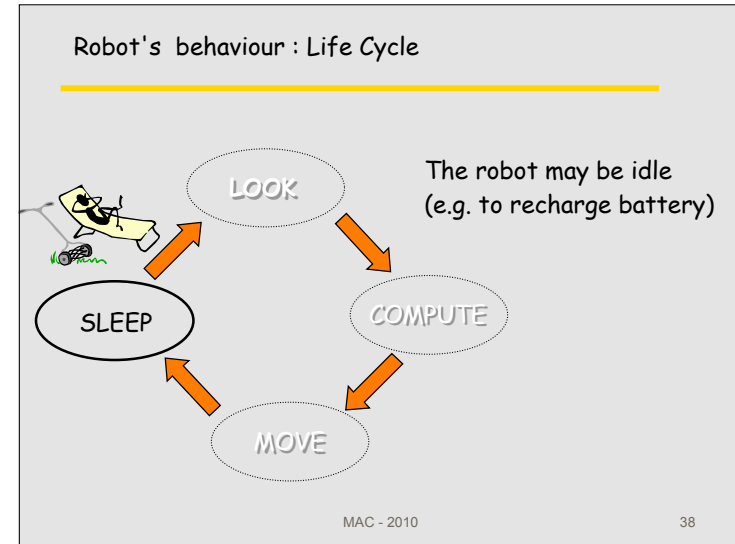
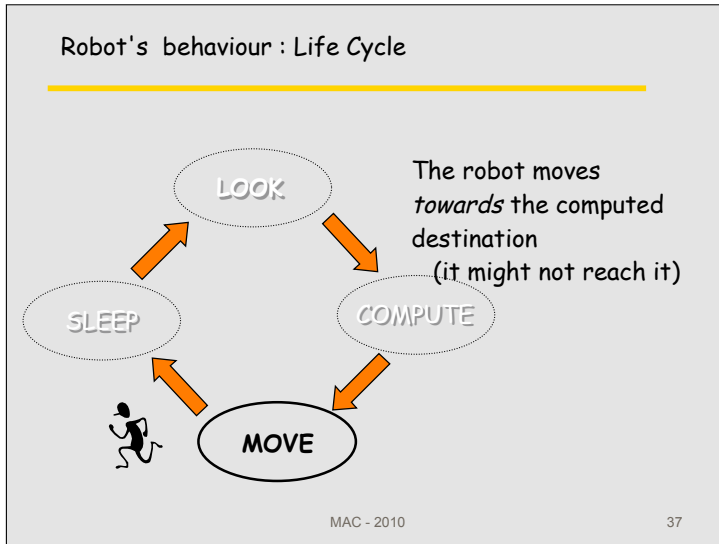
35

Robot's behaviour : Life Cycle



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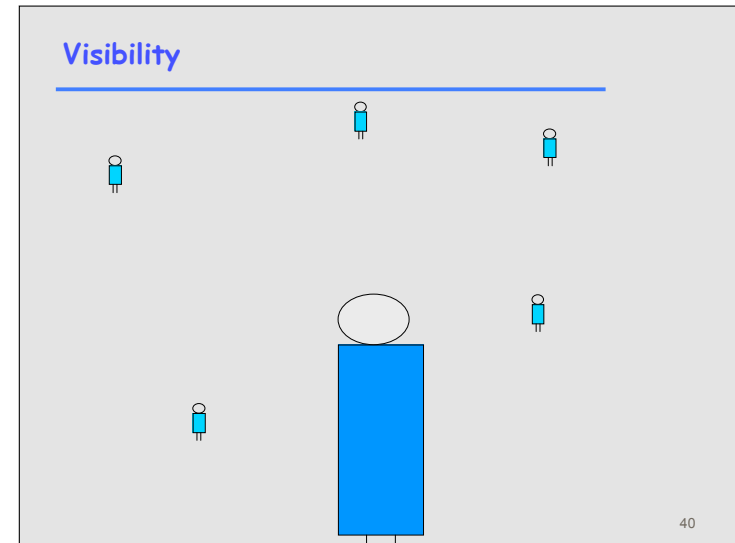
36



Crucial Factors

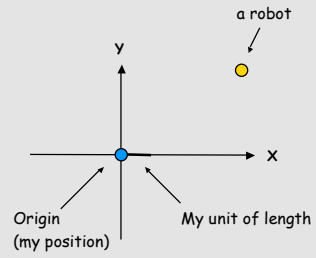
- Visibility

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Visibility

- local coordinate system
- "see" robots as points



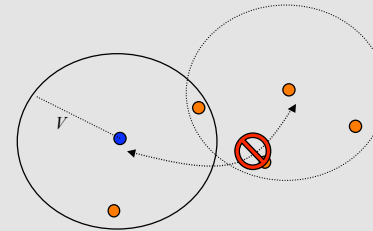
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Visibility

- Radius

Limited

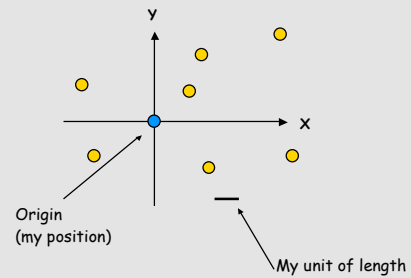


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Visibility : Snapshot

- Snapshot returns the coordinates of the other robots (seen as points) in terms of my local coordinate system

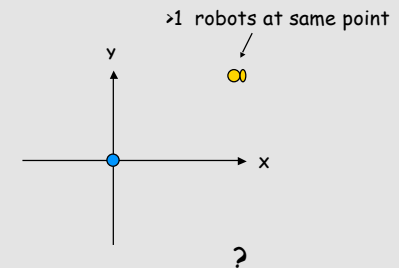


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Visibility : Multiplicity Detection

robots as points



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Visibility : Mobility Detection

in snapshot:

moving robot

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Crucial Factors

- **Visibility** Limited → Unlimited
- **Memory**

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Memory

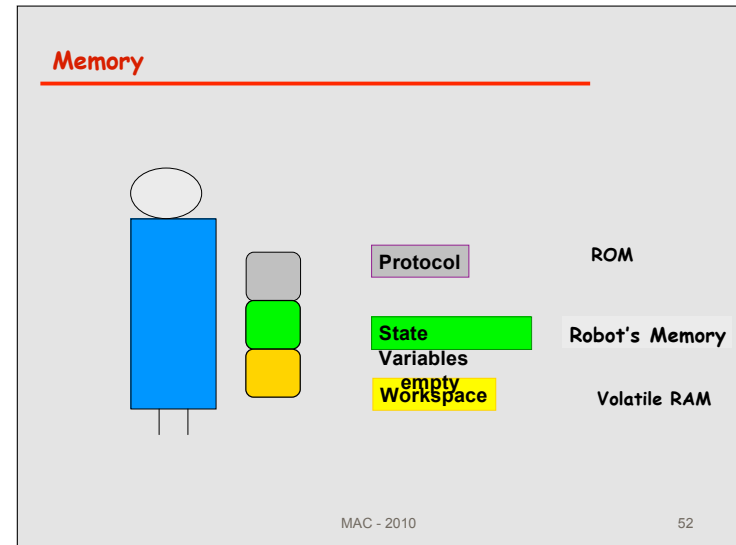
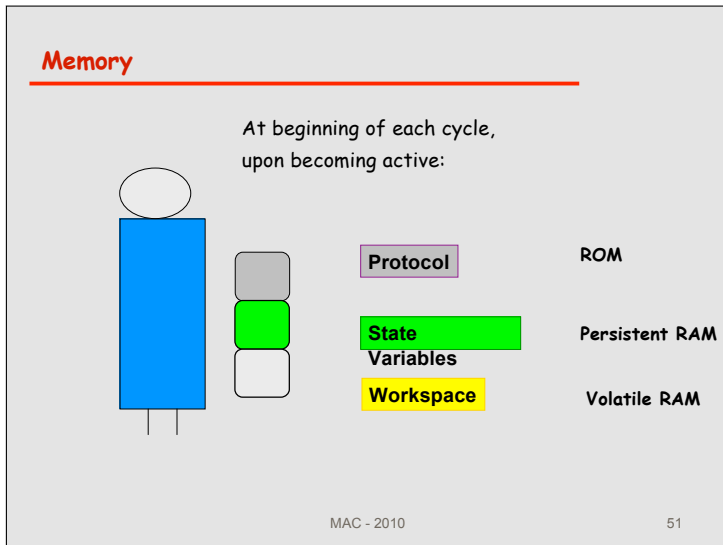
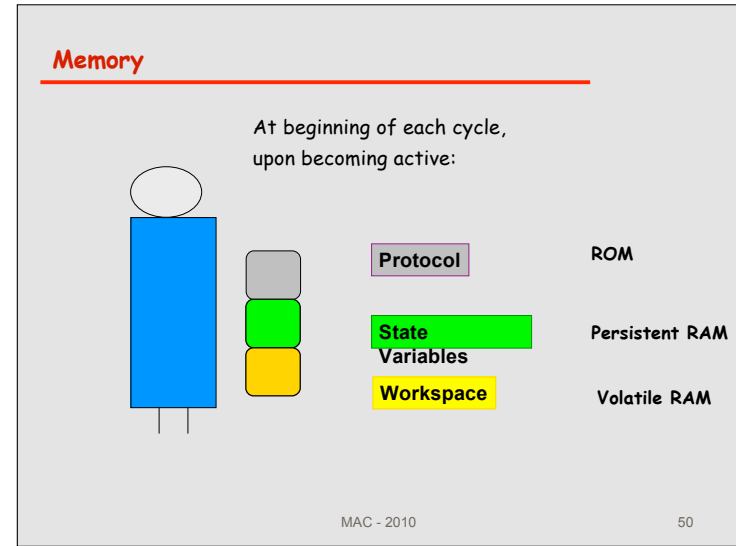
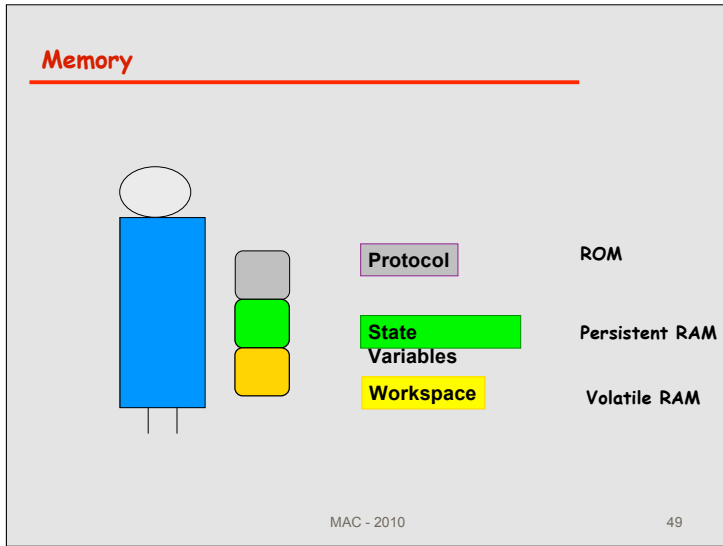
memory

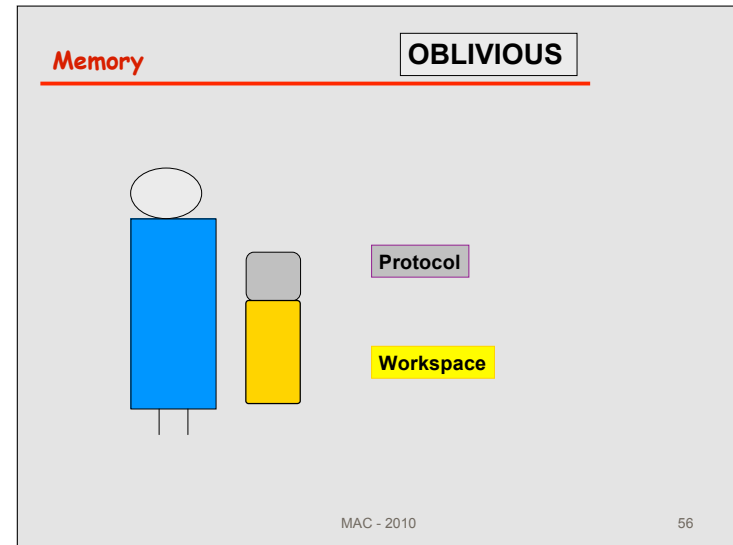
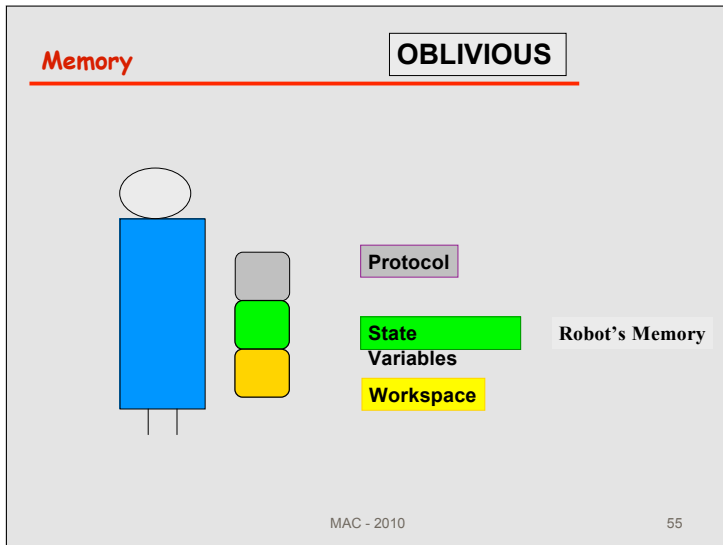
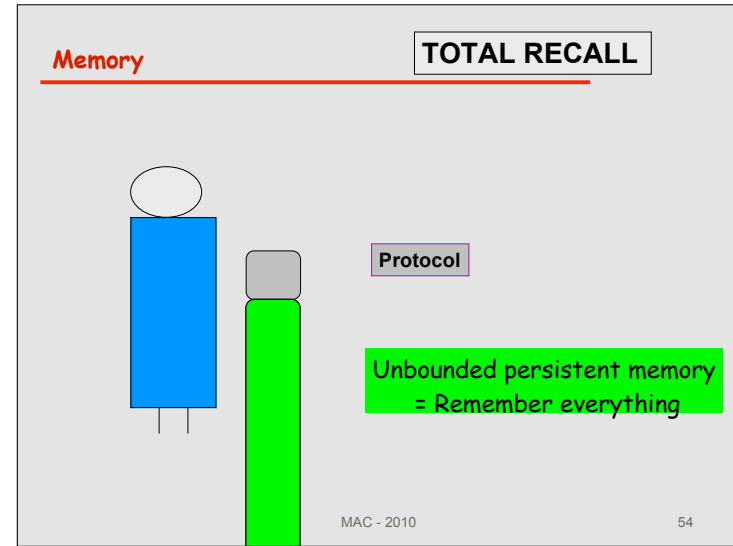
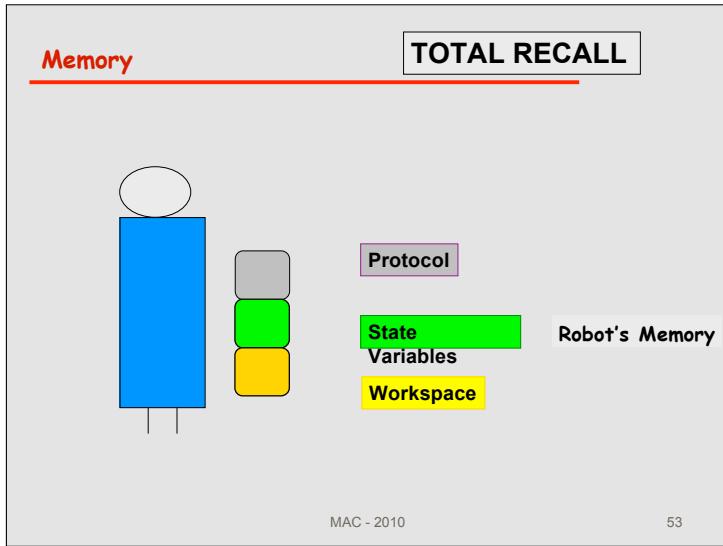
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Memory

memory

MAC - 2010 48





Memory

OBLIVIOUS

At beginning of each cycle,
upon becoming active:

Protocol

empty

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Memory

OBLIVIOUS

No memory of the past:
everytime is like the first time

Protocol

empty

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Crucial Factors

- **Visibility** Unlimited → **Limited**
- **Memory** Persistent → Oblivious

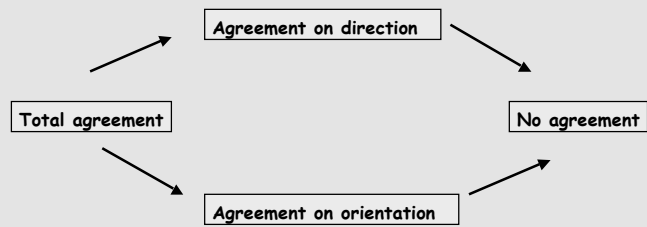
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Crucial Factors

- **Visibility** Unlimited → **Limited**
- **Memory** Persistent → Oblivious
- **Agreement** on local coordinate systems

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Levels of Agreement on Coordinate System

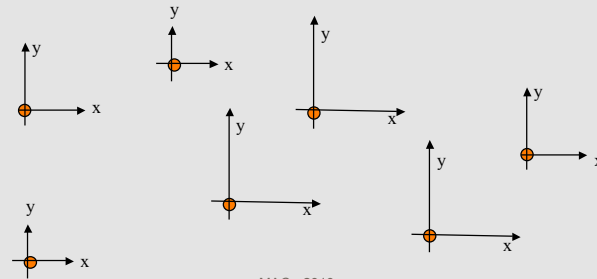


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Agreement 1: Direction and Orientation

The robots agree on a common **direction** and **orientation** of both axes

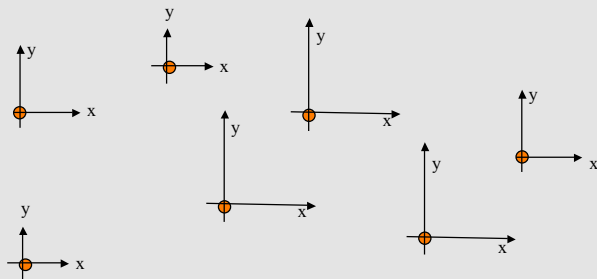


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Total Agreement: Direction and Orientation

The robots agree on a common **direction** and **orientation** of both axes

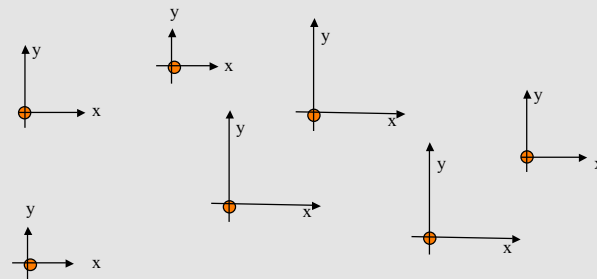


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Total Agreement: Direction and Orientation

Might disagree on **origin** and **unit** of length

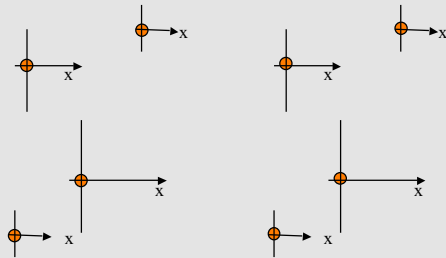


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Agreement 2: without Chirality

The robots agree on a common **direction** and **orientation of one axis**

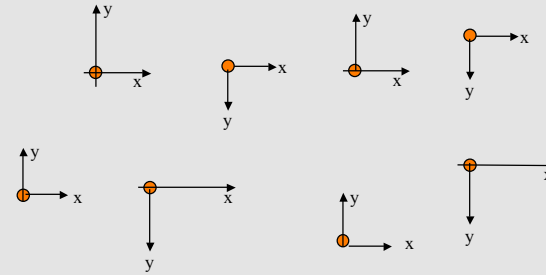


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Agreement 2: without Chirality

The robots agree on a common **direction** and **orientation of one axis**

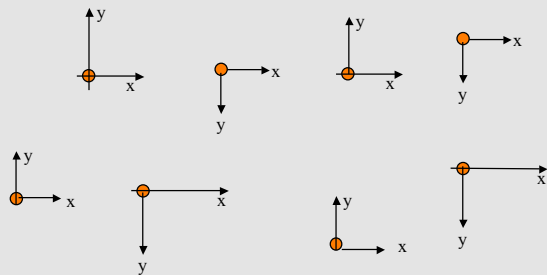


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Agreement 2: without Chirality

Might disagree on **origin** and **unit of length**
might NOT share the same **handedness**

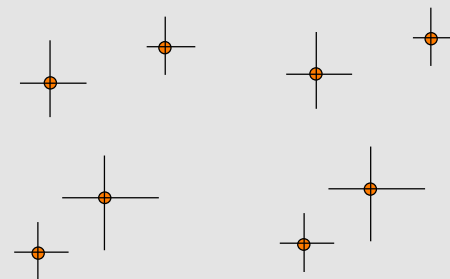


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Agreement 3: Direction

The robots agree on **direction** of both axes

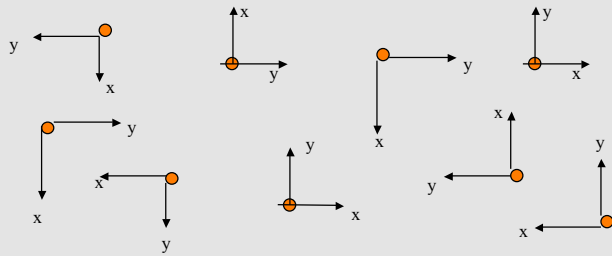


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Agreement 3: Direction

The robots agree on **direction** of both axes

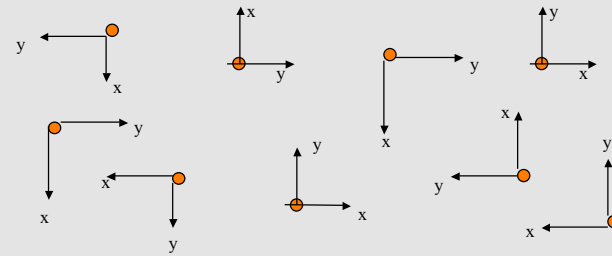


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Agreement 3: Direction

Might disagree on **origin** and **unit** of length
might disagree on **orientation** of axes

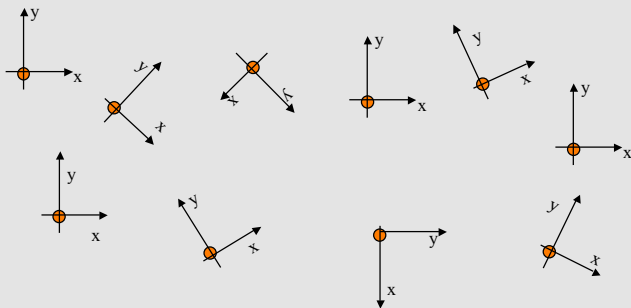


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Agreement 4: Orientation

The robots agree on **orientation** of both axes

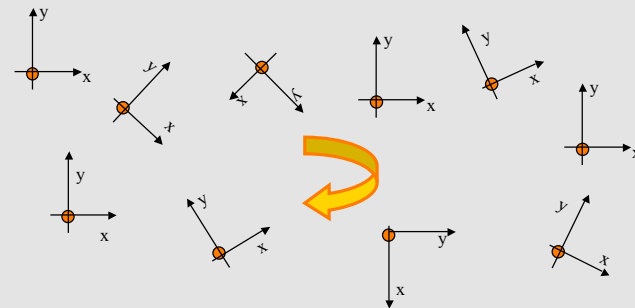


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Agreement 4: Orientation

The robots agree on **circular orientation** of plane

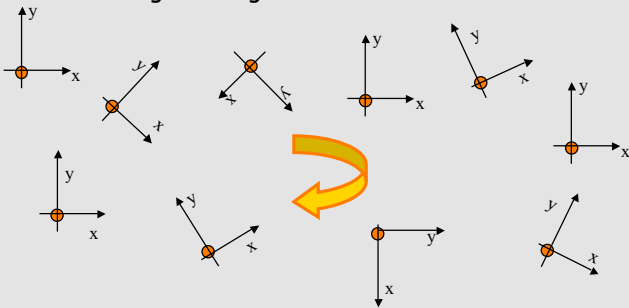


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Agreement 4: Orientation

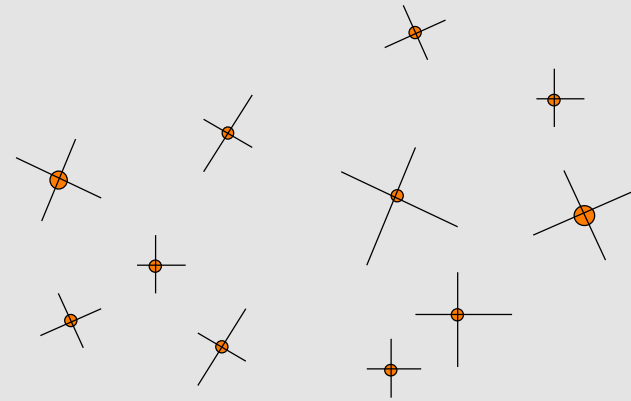
Might disagree on **origin** and **unit** of length
might disagree on **direction** of axes



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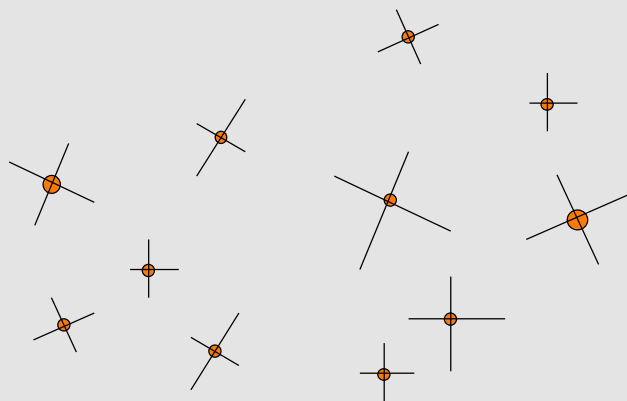
Agreement 5:



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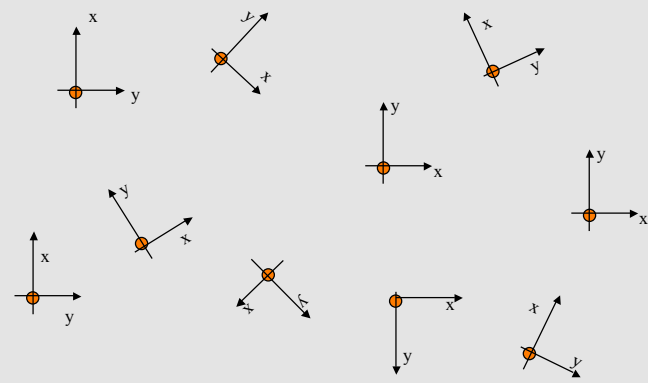
NO Agreement



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NO Agreement



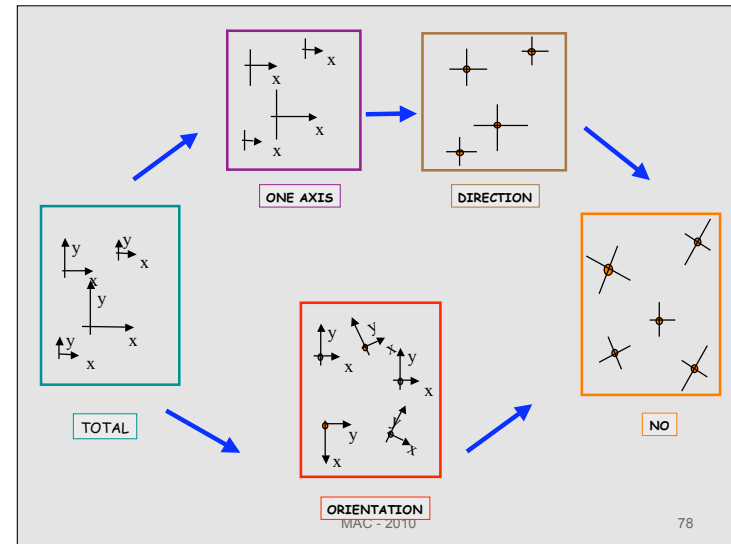
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Levels of Agreement on Coordinate System

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Crucial Factors

- **Visibility** Unlimited → Limited
- **Memory** Persistent → Oblivious
- **Agreement** on local coordinate systems

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Crucial Factors

- **Visibility** Unlimited → Limited
- **Memory** Persistent → Oblivious
- **Agreement** on local coordinate systems
- **Time/Synchronization**

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Time / Synchronization

There are three basic models

- Fully synchronous (FSYNC)
- Semi synchronous (SSYNC)
- Asynchronous (ASYNC)

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Time: Fully Synchronous (FSYNC)

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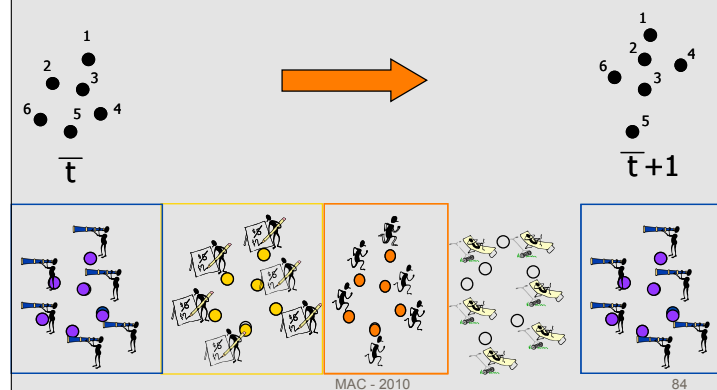
Time: Fully Synchronous (FSYNC)

- There is a **global clock tick** reaching all robots simultaneously
- At each clock tick every robot become **active** and perform its cycle atomically

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Time: Fully Synchronous (FSYNC)



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Time: Semi Synchronous (SSYNC)

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Time: Semi Synchronous (SSYNC)

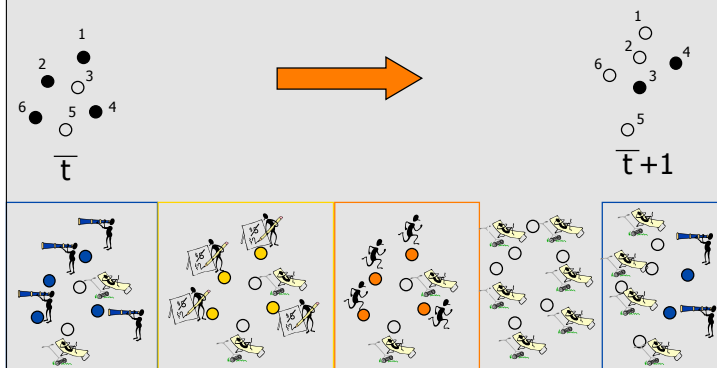
- There is a **global clock tick** reaching all robots simultaneously
- At each clock tick every robot is either **active** or **inactive**, and only active robots perform their cycle atomically.

[Suzuki Yamashita 96]

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Time: Semi Synchronous (SSYNC)



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Time: Semi Synchronous (SSYNC)

Fair Scheduler:
every robots becomes **active**
infinitely often

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Time: Asynchronous (ASYNC)

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Time: Asynchronous (ASYNC)

- There is **NO** global clock and robots do not have a common notion of time
- Each robot becomes active at **unpredictable** time instants
- Each computation and movement takes a finite but **unpredictable** amount of time

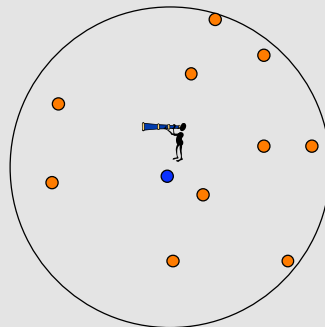
[Flocchini Prencipe Santoro Widmayer 99]

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Time: Asynchronous (ASYNC)

LOOK

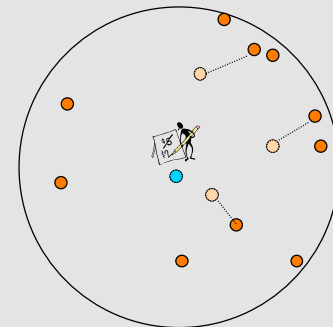


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Time: Asynchronous (ASYNC)

COMPUTE



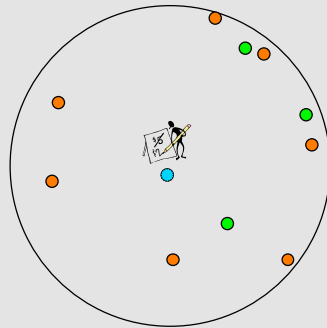
MAC - 2010

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Time: Asynchronous (ASYNC)

Computation based
on
obsolete information

COMPUTE



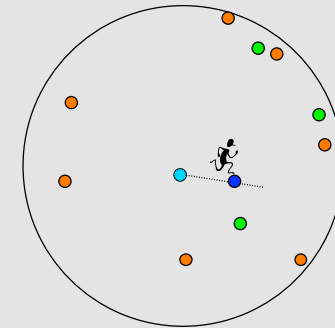
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Time: Asynchronous (ASYNC)

Movement based on
obsolete information

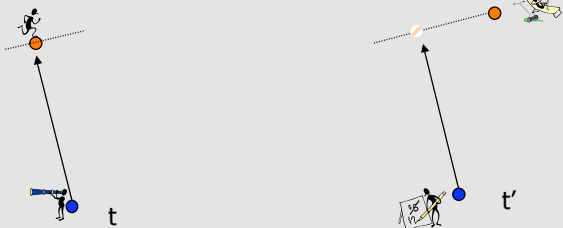
MOVE



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Time: Asynchronous (ASYNC)



A robot could see other robots while they move !

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Crucial Factors

- **Visibility** Unimited \rightarrow **Limited**
- **Memory** Persistent \rightarrow **Oblivious**
- **Agreement** Total \rightarrow **No**
- **Time/Synchronization** FSYNCH \rightarrow **ASYNCH**

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Crucial Factors

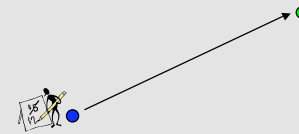
- **Visibility** Unimited → **Limited**
- **Memory** Persistent → **Oblivious**
- **Agreement** Total → **No**
- **Time/Synchronization** FSYNCH → **ASYNCH**
- **Mobility**

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Mobility :

Executing the protocol with the state variables and the current view as input, gives as output a **destination point** ● .

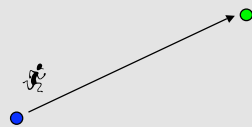


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Mobility 1:

In the MOVE ...

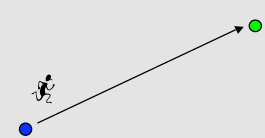


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Mobility 1: UNLIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ●

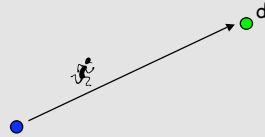


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Mobility 1: UNLIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ●

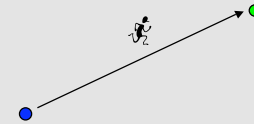


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Mobility 1: UNLIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ●

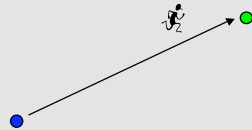


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Mobility 1: UNLIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ●

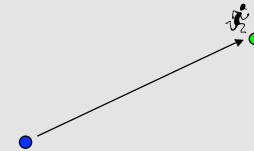


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Mobility 1: UNLIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ●

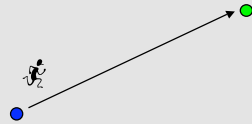


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Mobility 2:

In the MOVE ...

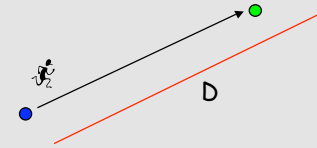


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Mobility 2: LIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ● or has traveled a distance D.

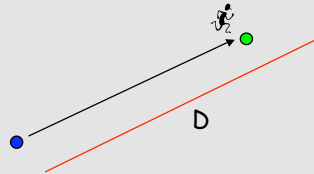


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106

Mobility 2: LIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ● or has traveled a distance D.

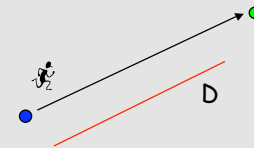


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Mobility 2: LIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ● or has traveled a distance D.

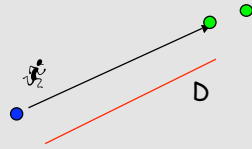


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Mobility 2: LIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ● or has traveled a distance D.

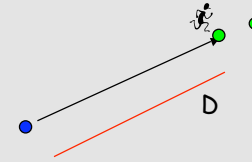


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Mobility 2: LIMITED MOTORIAL ENERGY

In the MOVE the robot moves continuously until it reaches ● or has traveled a distance D.



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Mobility 3:

In the MOVE ...

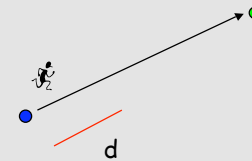


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Mobility 3: UNPREDICTABLE DISTANCE

In the MOVE the robot can stop anytime.
If it does not reach ●, it moves at least d.

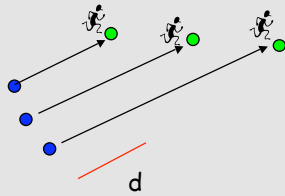


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Mobility 3: UNPREDICTABLE DISTANCE

In the MOVE the robot can stop anytime.
If it does not reach ●, it moves at least d .

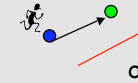


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Mobility 3: UNPREDICTABLE DISTANCE

In the MOVE the robot can stop anytime.
If it does not reach ●, it moves at least d .

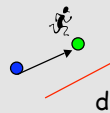


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Mobility 3: UNPREDICTABLE DISTANCE

In the MOVE the robot can stop anytime.
If it does not reach ●, it moves at least d .



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HOMOGENEOUS

every robot executes the same **protocol**

all robots are **identical** and **indistinguishable**

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TERMINATION

NON-OBLIVIOUS

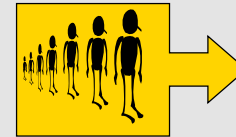
within finite time, every robot enters a **terminal** state

OBLIVIOUS

within finite time, every robot performs only **null** actions

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IMPOSSIBILITY

No **deterministic** protocol exists which **always** correctly solves the problem **terminating** in finite time

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Arbitrary
Pattern Formation

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Pattern Formation

Unlimited Visibility

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Pattern Formation

Pattern P is a set of (distinct) coordinates in the plane.

Number of robots is equal to the number of coordinates.

At the beginning,

the robots are in distinct arbitrary positions.

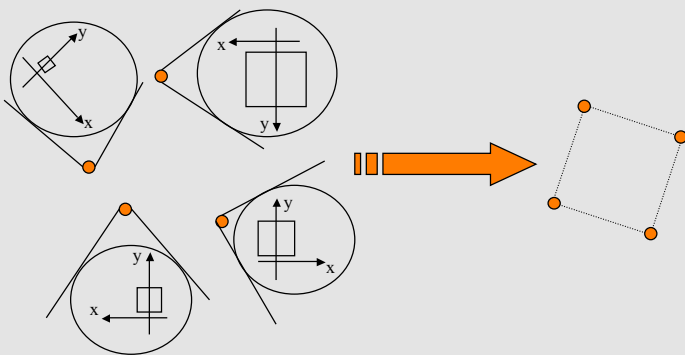
At the end,

the positions of the robots must "coincide"
with P (*translation, scaling or rotation*).

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Pattern Formation



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Computation/Protocol

FORMATION

the pattern is **formed within finite time**
(terminating protocol)

CONVERGENCE

the pattern is **never formed** but the robots'
position **converges** to the pattern in the limit
(non-terminating protocol)
(terminating approximate solution)

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Arbitrary Pattern Formation

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Arbitrary Pattern Formation

The n robots must form
any input pattern P of n points
starting from *any* initial position

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Arbitrary Pattern Formation

... is not always possible

Theorem

For $n > 2$, if *Arbitrary Pattern Formation*
is solvable with level of agreement A
then
the *Leader Election* problem is solvable
with level of agreement A

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Leader Election Problem

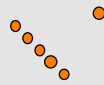
After a *finite* number of cycles, all the
robot *deterministically agree* on (choose)
the same robot L , called the *leader*

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Leader Election Problem

After a *finite* number of cycles, all the robot **deterministically agree** on (choose) the same robot **L**, called the **leader**



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129

Leader Election Problem

After a *finite* number of cycles, all the robot **deterministically agree** on (choose) the same robot **L**, called the **leader**



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APF : Agreement and Computability



Agreement on Coordinate System

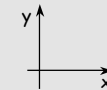
(e.g., compass, GPS)

even if the robots are **asynchronous and oblivious**

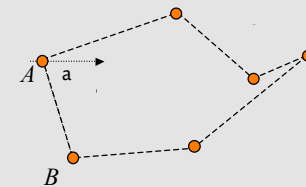
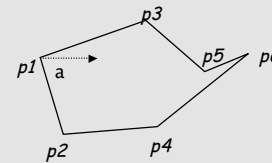
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APF:




: The Algorithm



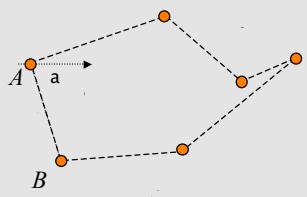
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APF:  : The Algorithm

If I am the leftmost ...

If I am the second leftmost ...



If I am not the leftmost and the angle between the leftmost and the second leftmost is $\neq \alpha$: I DON'T MOVE

If I am not the leftmost and the angle between the leftmost and the second leftmost is $= \alpha$: ...


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APF: Total Agreement

Theorem.
 With a common coordinate system, a set of *asynchronous anonymous oblivious* robots can form **any** input pattern in **finite** time.


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APF : Agreement and Computability



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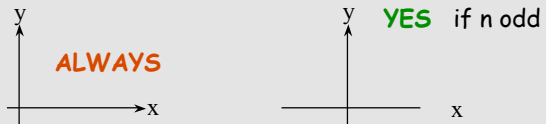
APF : Agreement and Computability



No Chirality

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APF : Agreement and Computability



even if asynchronous and oblivious

[Flocchini Prencipe Santoro Widmayer 02]

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APF : Agreement and Computability

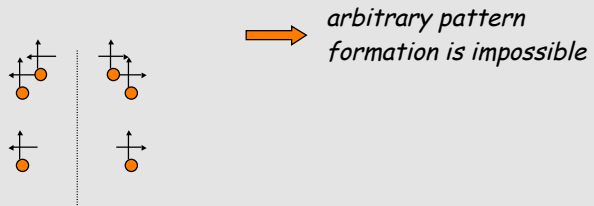


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Theorem. If the number of robots is even, the **Leader Election** problem is unsolvable without chirality.

Even if they have unlimited persistent memory
And are fully synchronous. [FPSW 99]

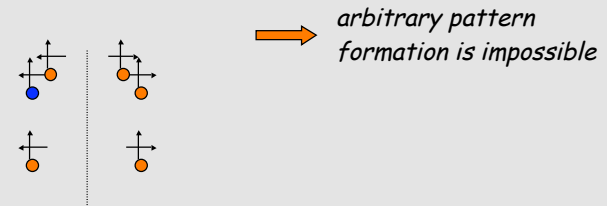


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139

Theorem. If the number of robots is even, the **Leader Election** problem is unsolvable without chirality.

Even if they have unlimited persistent memory
And are fully synchronous. [FPSW 99]

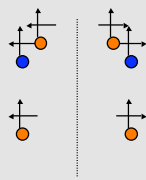


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Theorem. If the number of robots is even, the **Leader Election** problem is unsolvable without chirality.

Even if they have unlimited persistent memory
And are fully synchronous. [FPSW 99]

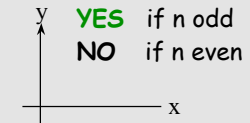


arbitrary pattern formation is impossible

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APF : Agreement and Computability



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Pattern Formation : n even and no chirality

Q: Which patterns can be formed when n is even and no chirality?

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Pattern Formation : n even and no chirality

Q: Which patterns can be formed when n is even and no chirality?

A: **NO** asymmetric patterns

even if

- unbounded persistent memory
- fully synchronous system

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NO ASYMMETRIC PATTERNS

By contradiction: \mathcal{A} lets the robots form **asymmetric patterns**, starting from an *arbitrary* initial configuration

Let us execute \mathcal{A} to form \mathcal{P} (asymmetric)

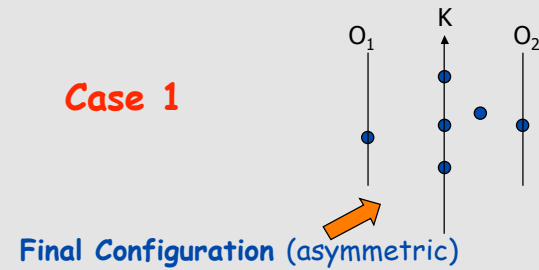
June 13, 2007

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NO ASYMMETRIC PATTERNS

Case 1



June 13, 2007

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NO ASYMMETRIC PATTERNS

Case 1

Contradiction!!



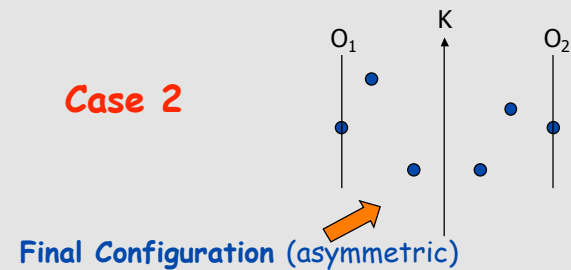
June 13, 2007

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NO ASYMMETRIC PATTERNS

Case 2

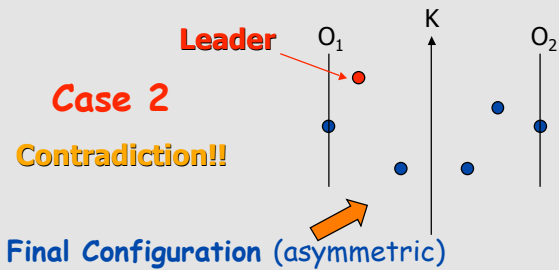


June 13, 2007

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NO ASYMMETRIC PATTERNS



June 13, 2007

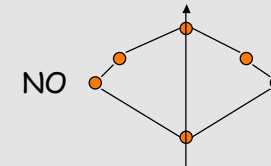
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Pattern Formation : n even and no chirality

Q: Which patterns can be formed when n is even and no chirality?

A: symmetric patterns



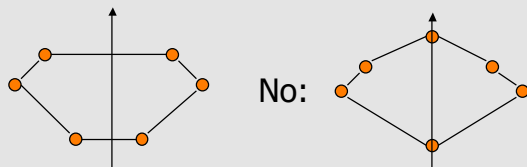
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150

Pattern Formation : n even and no chirality

Q: Which patterns can be formed when n is even and no chirality?

A: symmetric patterns



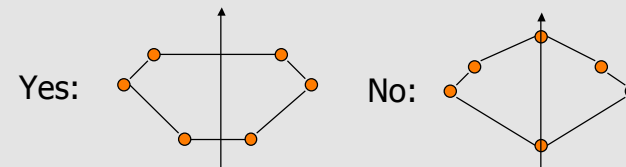
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Pattern Formation : n even and no chirality

Q: Which patterns can be formed when n is even and no chirality?

A: symmetric patterns with >0 empty axis of symmetry.



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Pattern Formation : n even and no chirality

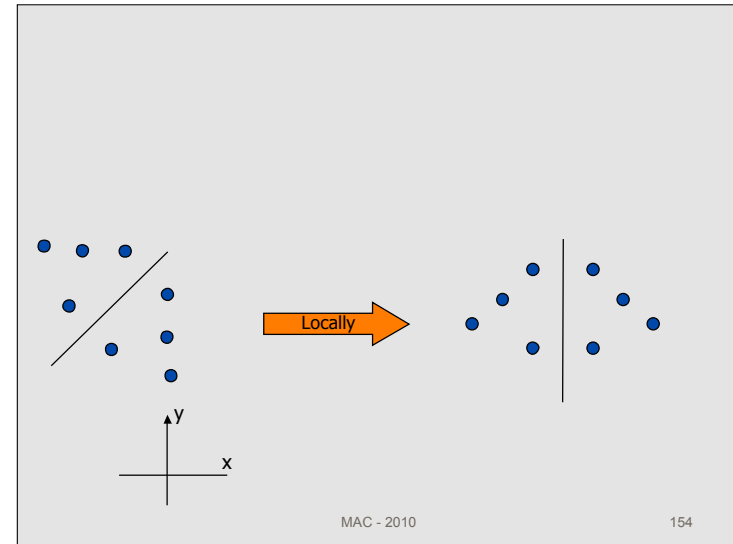
A: symmetric patterns with >0 empty axis of symmetry.

Local operations:

- Compute an **empty axis** of symmetry **S** (the same for all robots) of input pattern **P**
- Rotate **P** so that **S** is parallel to the agreed orientation of **Y**

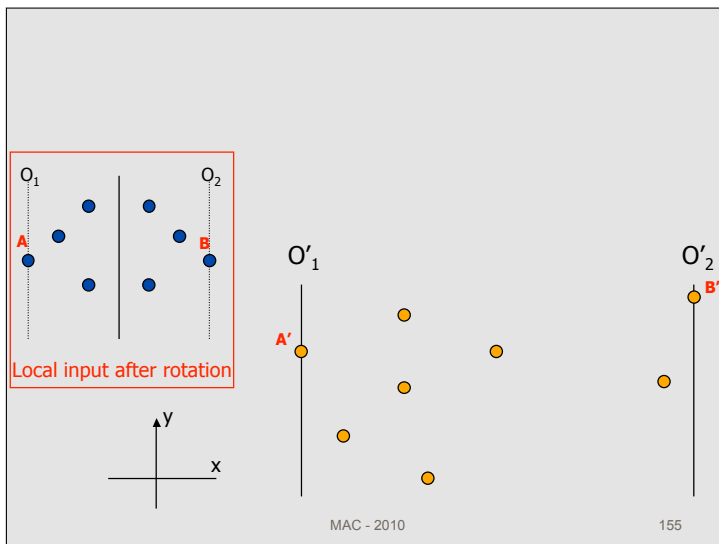
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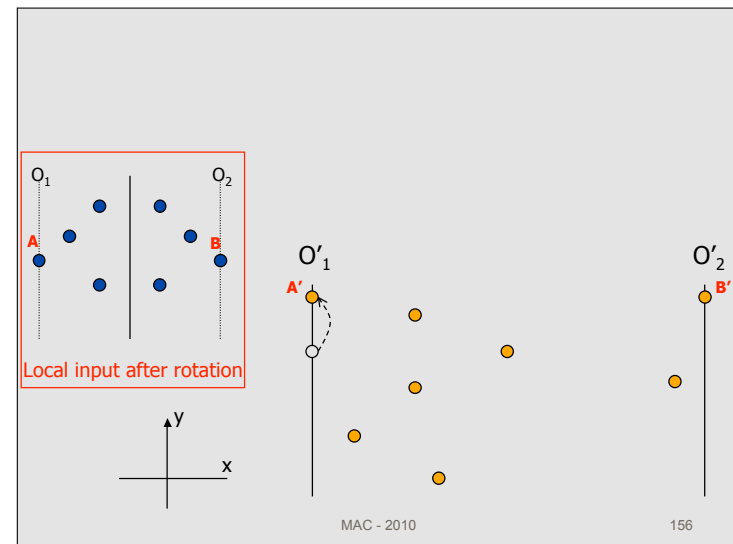
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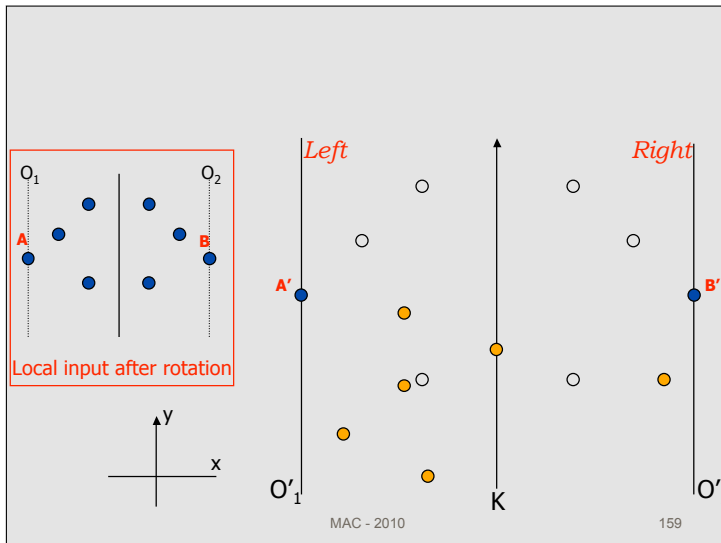
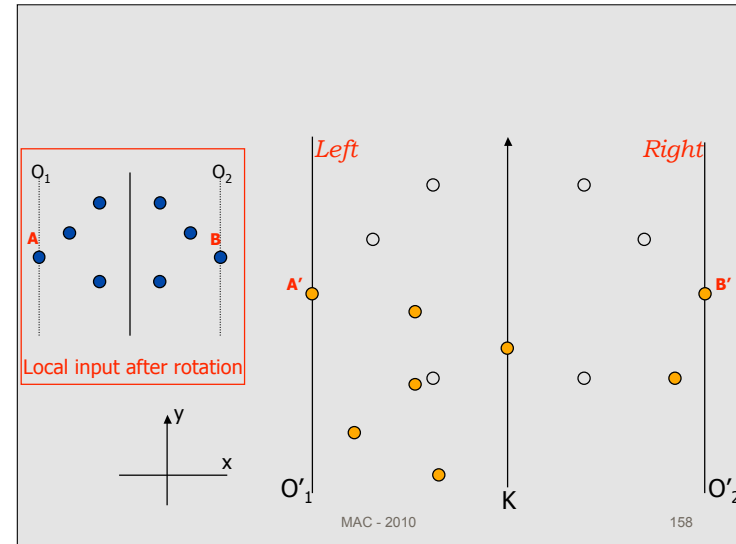
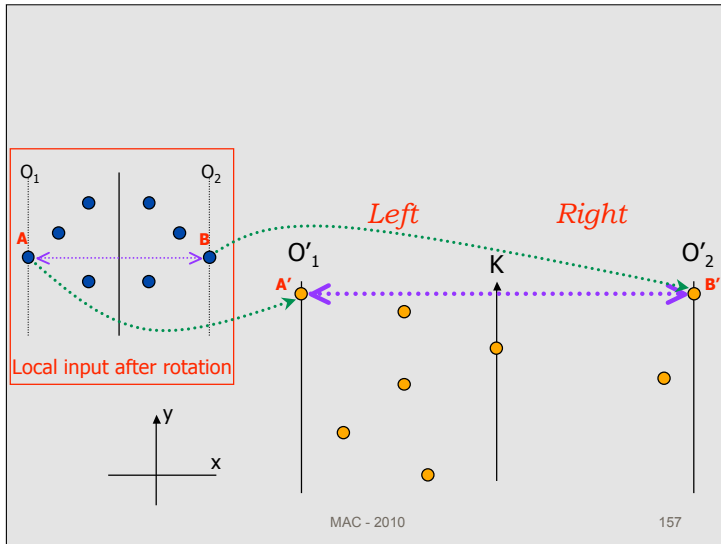
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- Robots on K:**
- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
 - If no *free* robot is inside Left/Right, the **topmost** on K moves towards one of the Final Positions still available
- Robots inside Left/Right:**
- The *free* robot in my side **closest** to a Final Position p still available in my side, moves towards p
 - If there are no Final Positions available in my side, go on K
- MAC - 2010 160

Robots on K:

- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
- If no *free* robot is inside Left/Right, the **topmost on K** moves towards one of empty Final Positions

Robots inside Left/Right:

- The *free* robot in my side closest to a Final Position p still available in my side, moves towards p
- If there are no Final Positions available in my side, go on K

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Robots on K:

- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
- If no *free* robot is inside Left/Right, the **topmost on K** moves towards one of empty Final Positions

Robots inside Left/Right:

- The *free* robot in my side closest to a Final Position p still available in my side, moves towards p
- If there are no Final Positions available in my side, go on K

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Robots on K:

- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
- If no *free* robot is inside Left/Right, the **topmost on K** moves towards one of empty Final Positions

Robots inside Left/Right:

- The *free* robot in my side closest to an empty Final Position in my side, moves towards it
- If there are no Final Positions available in my side, go on K

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Robots on K:

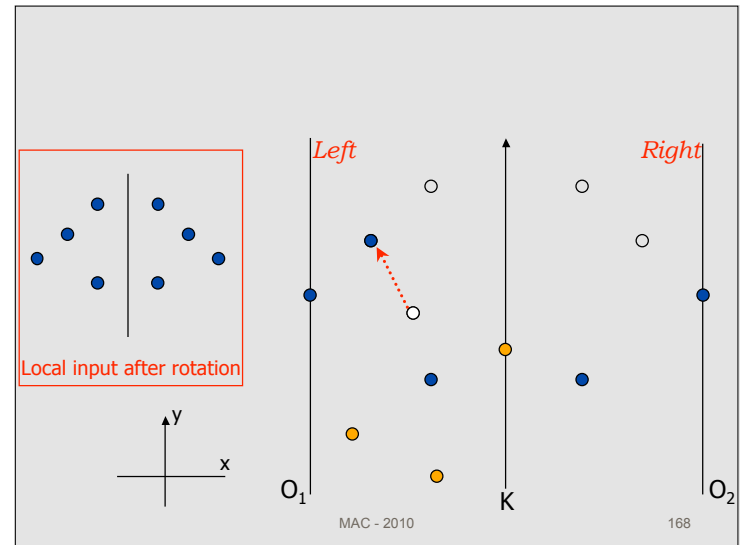
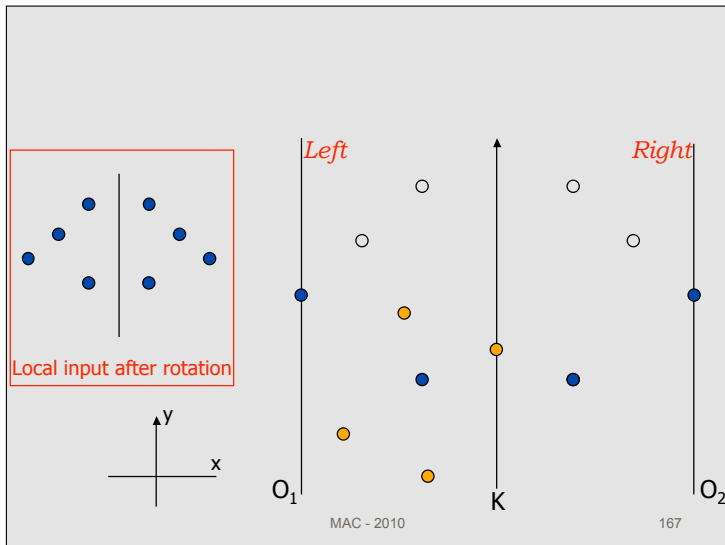
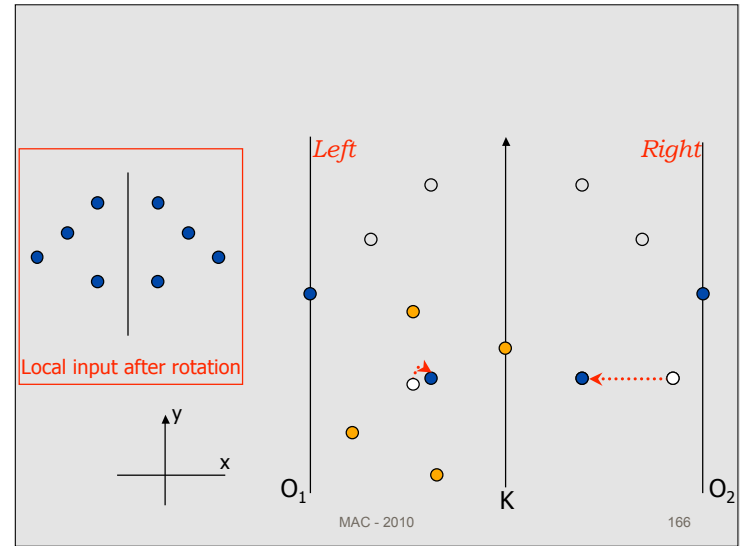
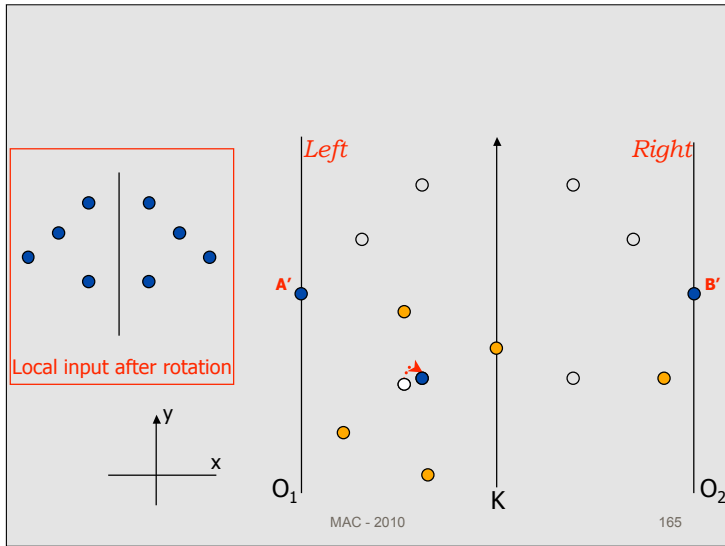
- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
- If no *free* robot is inside Left/Right, the **topmost on K** moves towards one of empty Final Positions

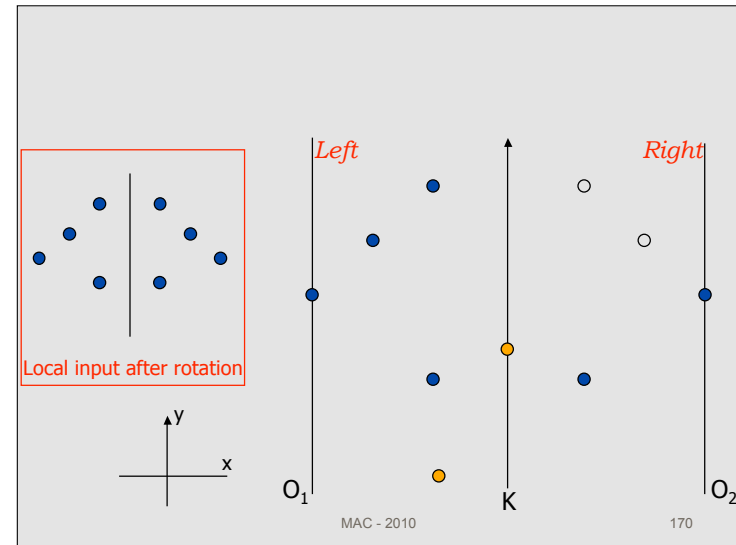
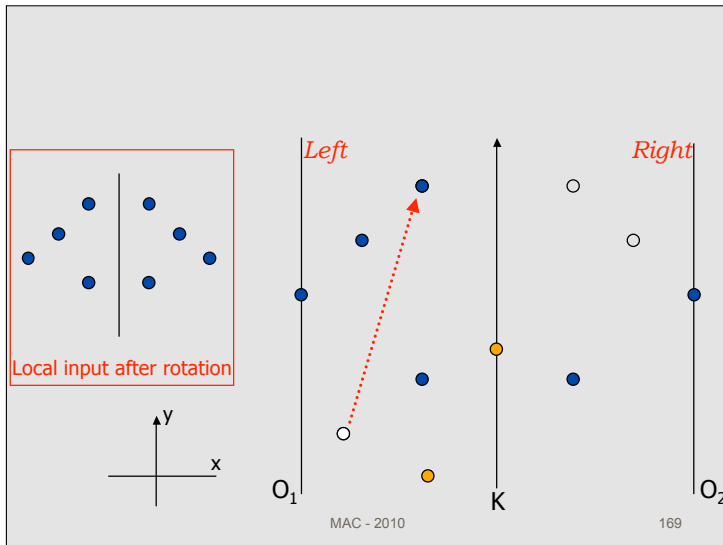
Robots inside Left/Right:

- The *free* robot in my side closest to an empty Final Position in my side, moves towards it
- If there are no Final Positions available in my side, go on K

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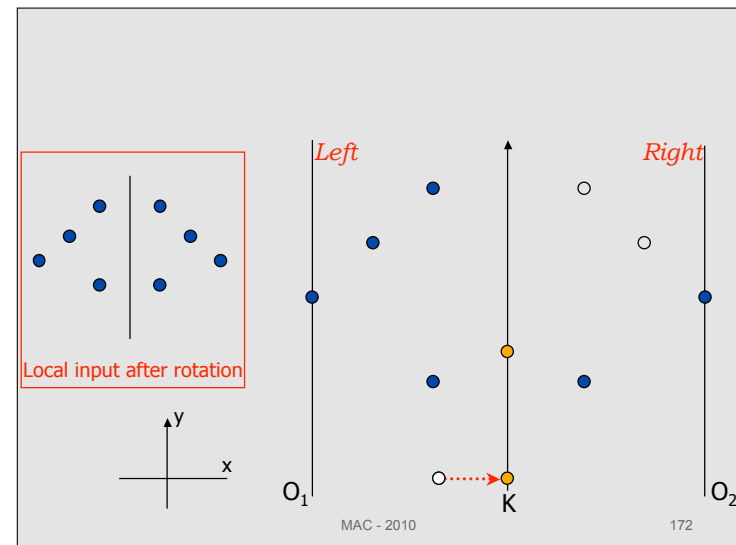
Robots on K:

- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
- If no *free* robot is inside Left/Right, the *topmost* on K moves towards one of empty Final Positions

Robots inside Left/Right:

- The *free* robot in my side closest to an empty Final Position in my side, moves towards it
- If there are no Final Positions available in my side, go on K

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Robots on K:

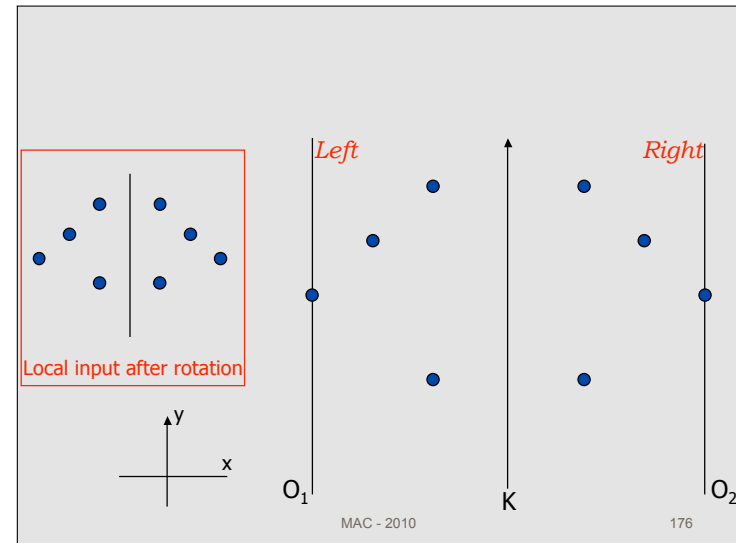
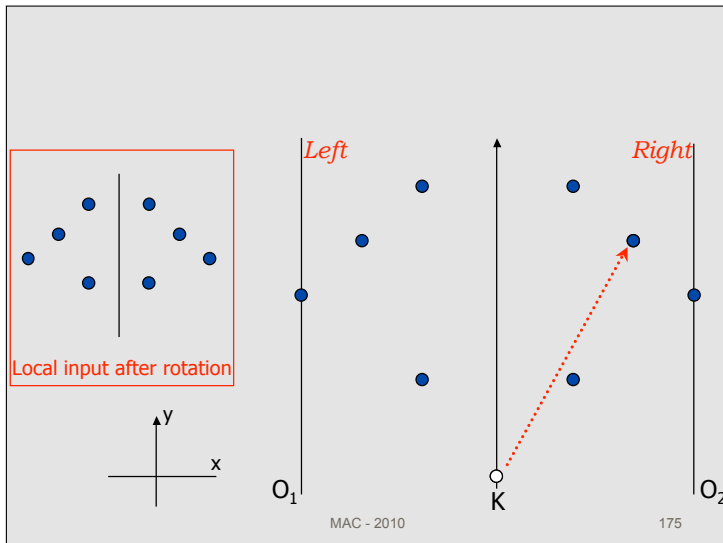
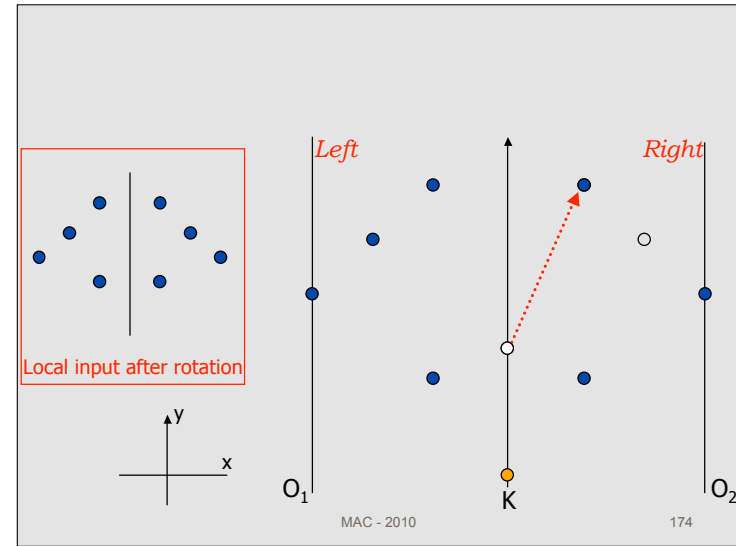
- Do not move as long as there are robots inside Left/Right not on Final Positions (*free* robots)
- If no *free* robot is inside Left/Right, the **topmost** on K moves towards one of empty Final Positions

Robots inside Left/Right:

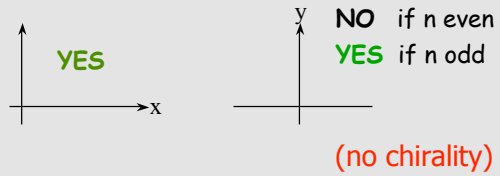
- The *free* robot in my side closest to an empty Final Position in my side, moves towards it
- If there are no Final Positions available in my side, go on K

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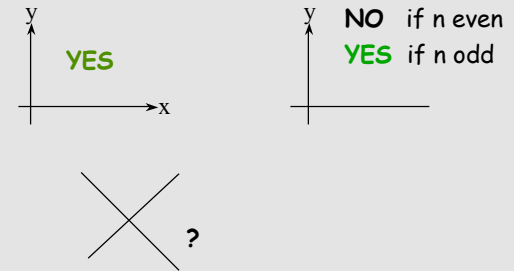
APF : Agreement and Computability



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APF : Agreement and Computability



No agreement

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APF : Agreement and Computability

even if

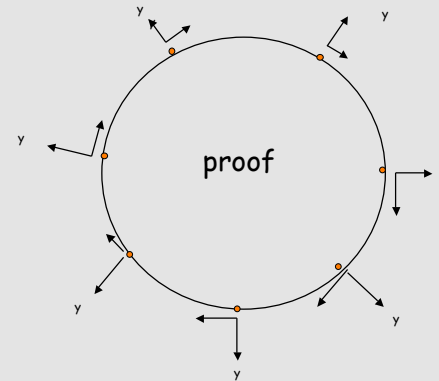
- unbounded persistent memory
- fully synchronous system



No agreement

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n-GON

A diagram showing a regular n-gon inscribed in a circle. The vertices of the n-gon are marked with small red dots. From each vertex, an arrow labeled 'y' points outwards, perpendicular to the side of the n-gon. The word 'proof' is written in the center of the circle.

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any totally synchronous execution will transform the n-GON into another n-GON

A diagram identical to the one on slide 181, showing a regular n-gon with vertices marked by red dots and arrows labeled 'y' pointing outwards from each vertex. The word 'proof' is in the center.

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... thus the only patterns that *can* be formed are n-GONS

A diagram identical to the one on slide 181, showing a regular n-gon with vertices marked by red dots and arrows labeled 'y' pointing outwards from each vertex. The word 'proof' is in the center.

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APF : NO AGREEMENT

even if

- unbounded persistent memory
- fully synchronous system

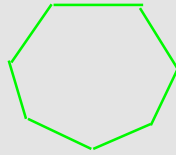
~~IMPOSSIBLE~~

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APF : NO AGREEMENT

corollary

The only patterns that *could* be formed without agreement are n-GONS



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Pattern Formation : no agreement

Q: can n-GONS be formed when there is no agreement ?

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Pattern Formation : no agreement

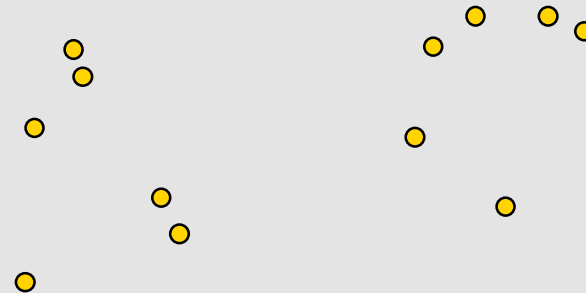
Q: can n-GONS be formed when there is no agreement ?

Forming n-gon = uniform circle formation

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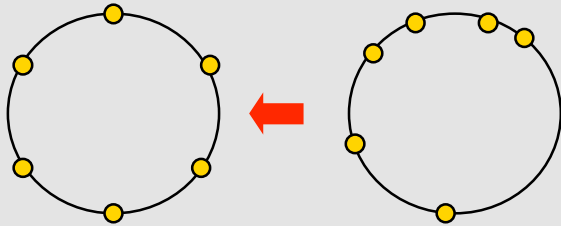
Circle Formation



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Uniform Circle Formation



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Pattern Formation : no agreement

Q: can n-GONs be formed when there is no agreement ?

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Pattern Formation : no agreement

Q: can **uniform circles** be formed when there is no agreement ?

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Pattern Formation : no agreement

Q: can **uniform circles** be formed when there is no agreement ?

YES with **unlimited persistent memory** in **SSYNCH**

[Yamashita and Suzuki, 1996]

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Pattern Formation : no agreement

Q: can **uniform circles** be formed when there is no agreement ?

YES with **unlimited persistent memory** in **SSYNCH**

[Yamashita and Suzuki, 1996]

CONVERGENCE oblivious [Defago Konagaya, 2002]

[Chatzigiannakis et al., 2004] [Defago Souissi, 2008]

SPECIAL SIZES oblivious [Dieudonne and Petit, 2007]

Pattern Formation : no agreement

Q: can **uniform circles** be formed when there is no agreement ?

YES with **unlimited persistent memory** in **SSYNCH**

[Yamashita and Suzuki, 1996]

YES **obliviously** in **SSYNCH**

[Yamashita and Suzuki, 2008]

[Petit and Dieudonne, 2009]

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Pattern Formation : no agreement

YES with **unlimited persistent memory** in **SSYNCH**

[Yamashita and Suzuki, 1996]

YES **obliviously** in **SSYNCH**

[Yamashita and Suzuki, 2008]

[Petit and Dieudonne, 2009]

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Pattern Formation : no agreement

$$\text{PAT}_{\text{SSYNCH}}[\text{OBLIVIOUS}] = \text{PAT}_{\text{SSYNCH}}[\text{NON-OBLIVIOUS}]$$

YES with **unlimited persistent memory** in **SSYNCH**

[Yamashita and Suzuki, 1996]

YES **obliviously** in **SSYNCH**

[Yamashita and Suzuki, 2008]

[Petit and Dieudonne, 2009]

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Pattern Formation : no agreement

PAT [OBLIVIOUS]_{SSYNCH} = PAT [NON-OBLIVIOUS]_{SSYNCH}

= PAT [OBLIVIOUS]_{FSYNCH} = PAT [NON-OBLIVIOUS]_{FSYNCH}



OBLIVIOUSNESS
IS NOT A HANDICAP
IN FSYNCH AND SSYNCH

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Pattern Formation : no agreement

WHAT ABOUT IN ASYNCH ?



OBLIVIOUSNESS
IS NOT A HANDICAP
IN FSYNCH AND SSYNCH

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Pattern Formation : no agreement

WHAT ABOUT IN ASYNCH ?

PAT [OBLIVIOUS]_{ASYNCH} ? = ? PAT [NON-OBLIVIOUS]_{ASYNCH}

CAN OBLIVIOUS ROBOTS FORM A
UNIFORM CIRCLE ASYNCHRONOUSLY ?

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Pattern Formation : no agreement

Some claims ...

[Katreniak 2005] [Petit Dieudonne, 2009]

CAN OBLIVIOUS ROBOTS FORM A
UNIFORM CIRCLE ASYNCHRONOUSLY ?

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