

The Computational Power of Oblivious Mobile Robots

Forming a series of geometric patterns

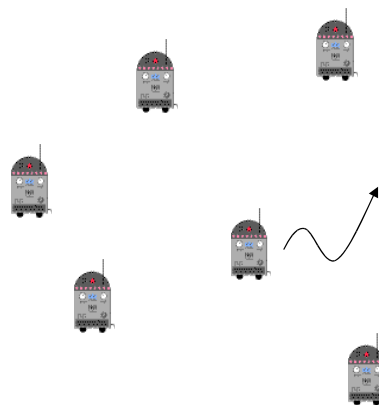
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Joint work with **Paola Flocchini, Nicola Santoro, and Masafumi Yamashita**

Team of Stupid Robots

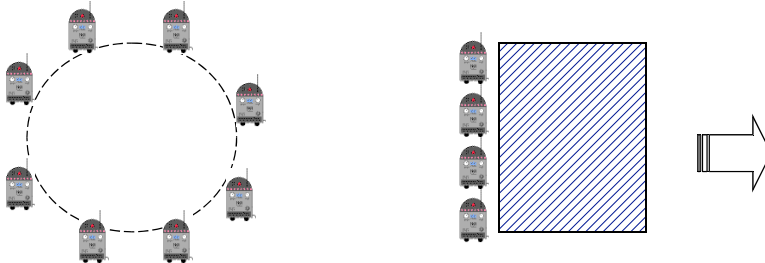
- **Autonomous** Robots
- Moving on a plane
- Can not communicate
- Have Vision (unlimited)
- Have no memory
- Semi-synchronous



Which tasks are possible for a team of stupid robots?

Pattern Formation

Can the robots arrange themselves on a *circle* ? on a *line* ?
or other geometric patterns?



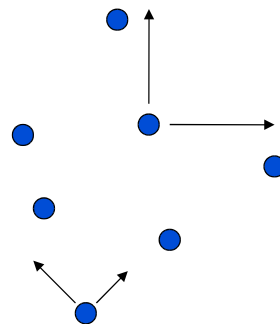
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Robot Model

- Points on the Plane [Dimensionless]
- Repeat
 - LOOK : Positions of robots
 - COMPUTE : new location
 - MOVE : to the computed location
- No agreement on Orientation
- No agreement on unit distance
- Identical algorithms
- Robots may be active or passive in ant time step



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Models of Synchrony

- **Synchronous Model**

Robots act in synchronized time steps; Start at same time;

- **Semi-Synchronous Model [SYm]**

In each time step, a subset of robots are active;

Active robots complete exactly one LOOK-COMPUTE-MOVE step.

- **Asynchronous Model [CORDA]**

Robots complete each step in an arbitrary amount of time.

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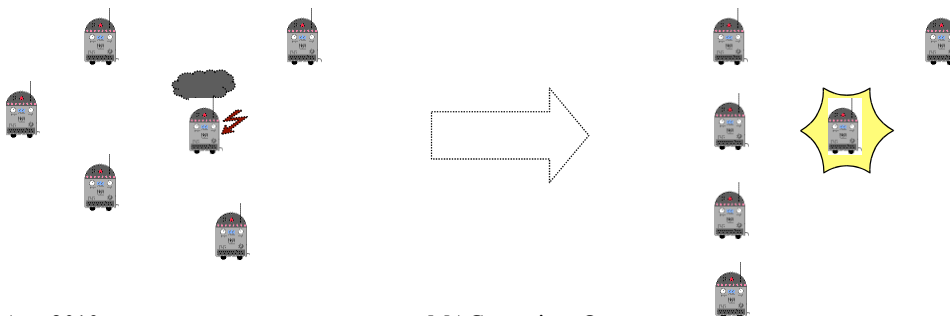
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Motivation

Why Oblivious ?

- Robots can **crash** (and recover at a later time).
- Robots may join at any time, in any state.
- Simple to design and analyze algorithms!



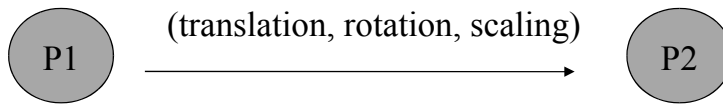
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Arbitrary Pattern Formation (APF)

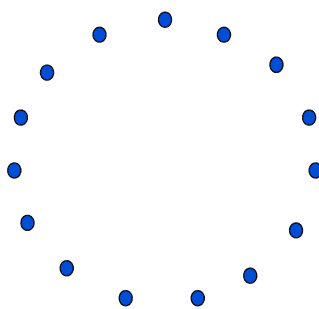
- **Pattern** : A set of points $\in \mathbb{R}^2$
- **Isomorphic Patterns** : $P1 \sim P2$



- A team of robots *forms a pattern* P if,
There is some coordinate system Z , such that :
Locations of robots = points in P (according to Z)
[Multiple robots may occupy the same point.]

Known Results

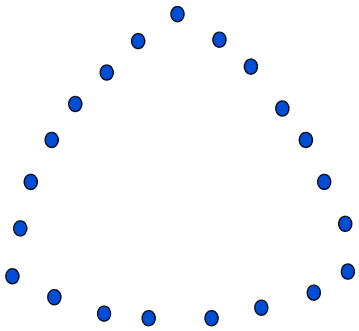
- **Circle Formation:** [Sugihara & Suzuki 1990]



a := farthest neighbor
 b := closest neighbor
 c := midpoint of the line segment (a,b)
If (distance to c) $>$ $|bc|$
 move one step towards c
If (distance to c) $<$ $|bc|$
 move one step away from c
Else, move one step away from b

Known Results

- **Circle Formation:** [Sugihara & Suzuki 1990]

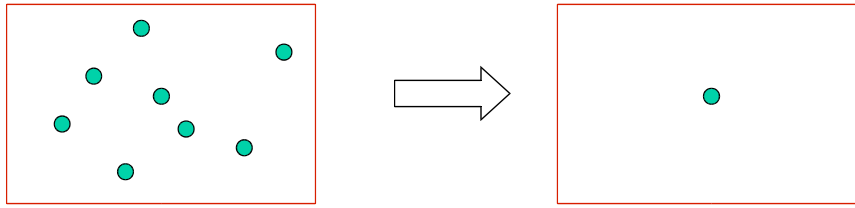


```
a := farthest neighbor
b := closest neighbor
c := midpoint of the line segment (a,b)
If (distance to c) > |bc|
    move one step towards c
If (distance to c) < |bc|
    move one step away from c
Else, move one step away from b
```

Circle Formation

- [Sugihara & Suzuki 1990] : Approximate Circle
- [Katreniak 2005] : Non-uniform Circle
- [Dieudonné & Petit 2007] : Uniform Circle ($k \neq 4$)
- [Défago & Souissi 2008] : Convergence to Uniform Circle
- [Suzuki & Yamashita 1999, 2010] : Uniform Circle

POINT formation (gathering)



Theorem [SY 1999]

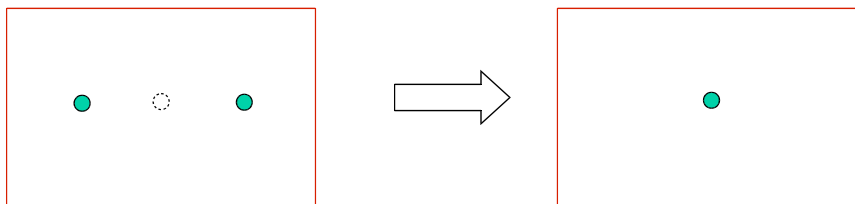
For $k = 2$ oblivious robots, it is **impossible** to form POINT.

However, the two robots can converge to a point!

Theorem [SY 1999]

For $k \geq 2$ non-oblivious robots, POINT formation is always possible.

POINT formation (gathering)



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For $k = 2$ oblivious robots, it is **impossible** to form POINT.

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For $k \geq 2$ non-oblivious robots, POINT formation is always possible.

Oblivious vs. Non-Oblivious

Arbitrary Pattern Formation

Theorem [SY 2010]

For $k \geq 3$ and any arbitrary pattern P ,

k *non-oblivious* robots can form P \Leftrightarrow k *oblivious* robots can form P

Oblivious robots are almost as powerful as non-oblivious robots!

Series of Patterns

- Can oblivious robots form a **series** of Patterns

$P_1, P_2, P_3, P_4, \dots$

How to remember if we do not have memory?

A robot system forms the series $\langle P_1, P_2, P_3, \dots, P_m \rangle$

if it goes through the series of configurations

$C_0 C_1 C_2, \dots, C_i, \dots, C_j, \dots, C_x, \dots, C_y, \dots$
 P_1, P_2, P_3, P_m

Which series are formable?

Theorem: No terminating algorithm can form a **finite** series of patterns $\langle P_1, P_2, P_3, \dots, P_m \rangle$, where $m > 1$

Want to form the infinite series: $(P_1, P_2, P_3, \dots, P_m)^\infty$

A series $\langle P_1, P_2, P_3, \dots, P_m \rangle$ is *formable* \Leftrightarrow

There exists an algorithm for forming the infinite series

$$(P_1, P_2, P_3, \dots, P_m)^\infty$$

Our Model (Adversary)

- **Semi-Synchronous:**

At each time step, the *adversary* decides which robots are active;

[Fairness: Each agent is activated infinitely often]

- **Lack of Orientation**

Adversary decides the local coordinate system of each robot at each step;

- **Arbitrary Initial Configuration**

Adversary decides the initial location of the robots.

Our Model (assumptions)

- **Agreement on Clockwise Direction**

Robots agree on clockwise / counter-clockwise



- **Multiplicity Detection**

Robots can **count** the number of robots at each location.

Three Scenarios

- **Visibly Distinct Robots**

- Each robot is distinct (has visible ID)!

- **Indistinguishable Robots with ID's**

- Robots look identical but are given invisible ID's!

- **Anonymous Robots**

- All robots look identical and execute the same algorithm!

Visibly Distinct Robots

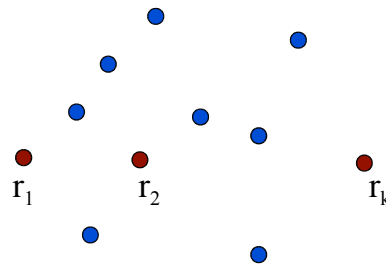
- k distinct robots

$$r_1, r_2, r_3, \dots, r_k$$

- To form pattern P_i

Encode information as ratio of distances:

$$\frac{\text{Dis}(r_1, r_k)}{\text{Dis}(r_1, r_2)} = F(P_i)$$

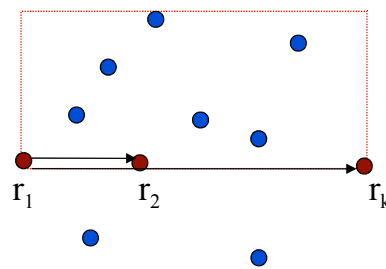


Visibly Distinct Robots

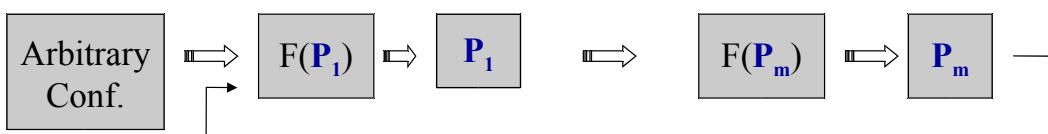
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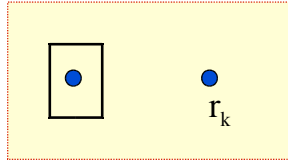


Robots can agree on a coordinate system!

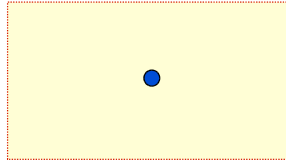


Visibly Distinct Robots

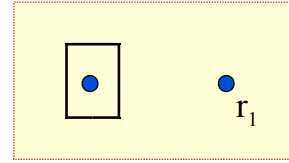
- What about forming the **POINT** pattern?



Before POINT



POINT



After POINT

Visibly Distinct Robots

Theorem: The series $\langle P_1, P_2, P_3, \dots, P_m \rangle$ is formable iff

- $\text{size}(P_i) \leq k$, for all i ,
- P_i is non-isomorphic to P_j , $i \neq j$

Second Scenario

- **Visibly Distinct Robots**

- Each robot is distinct (has visible ID)!

- **Indistinguishable Robots with ID's**

- Robots look identical but are given invisible ID's!

- **Anonymous Robots**

- All robots look identical and execute the same algorithm!

Indistinguishable Robots with ID's

- k distinct (but indistinguishable) robots

$$r_1, r_2, r_3, \dots, r_k$$

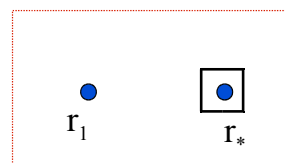
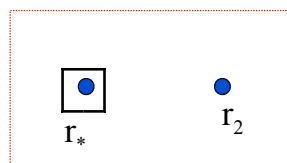
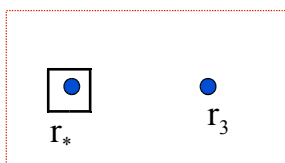
Theorem: (impossibility)

With $k = 3$ robots, the following series is not formable

$$(P_1, P_2, \dots, P_1) \text{ [POINT]} (P_{1+1}, \dots, P_m) \text{ [TWO-POINTS]}$$

where $\text{size}(P_i) = 3$, for all i . ($1 < m$)

No robot can distinguish between



Indistinguishable Robots with ID's

Possibility:

Theorem: With $k = 3$ robots, the series S^∞ is formable if S is a cyclic rotation of the following:

- $(P_1, P_2, P_3, \dots, P_m)$ [POINT] [TWO-POINTS]
- $(P_1, P_2, P_3, \dots, P_m)$ [TWO-POINTS][POINT]
- $(P_1, P_2, P_3, \dots, P_m)$ [POINT]
- $(P_1, P_2, P_3, \dots, P_m)$ [TWO-POINTS]
- $(P_1, P_2, P_3, \dots, P_m)$

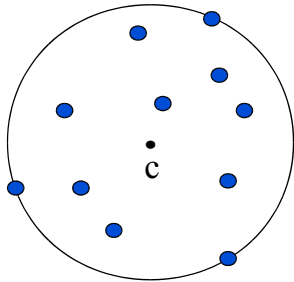
where $\text{size}(P_i) = 3$, for all i .

Indistinguishable Robots with ID's

Theorem: With $k > 3$ robots,
 $\langle P_1, P_2, P_3, \dots, P_m \rangle$ is formable iff

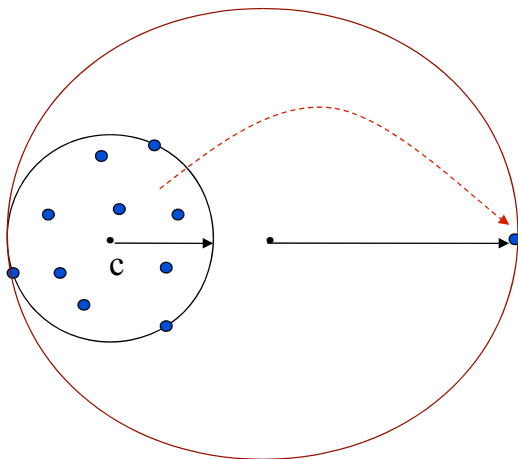
- $\text{size}(P_i) \leq k$, for all i ,
- P_i is non-isomorphic to P_j , $i \neq j$

Algorithm for $k > 3$ robots



Smallest Enclosing Circle (SEC): [UNIQUE]

Algorithm for $k > 3$ robots



Bi-Circular Configuration (BCC):

Encodes info about a pattern using ratio of two radii

Third Scenario

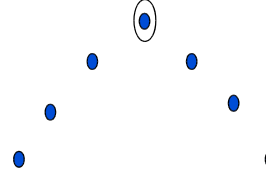
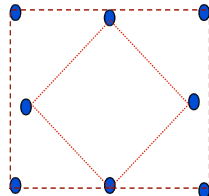
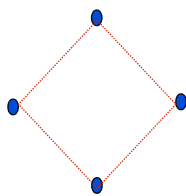
- **Visibly Distinct Robots**
 - Each robot is distinct (has visible ID)!
- **Indistinguishable Robots with ID's**
 - Robots look identical but are given invisible ID's!

- **Anonymous Robots**
 - All robots look identical and execute the same algorithm!

Anonymous Robots

Assumption: The robots start from distinct locations.

Robots in symmetric locations have similar view!



Symmetry $\rho = 4$

$\rho = 4$

$\rho = 1$

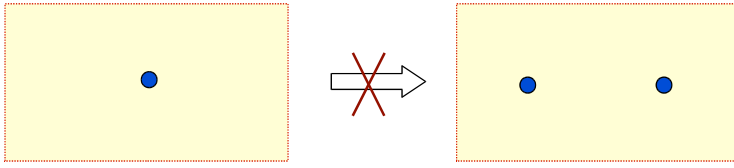
Size $k = 4$

$k = 8$

$k = 7$

Anonymous Robots

Two robots in the same location may never be separated!



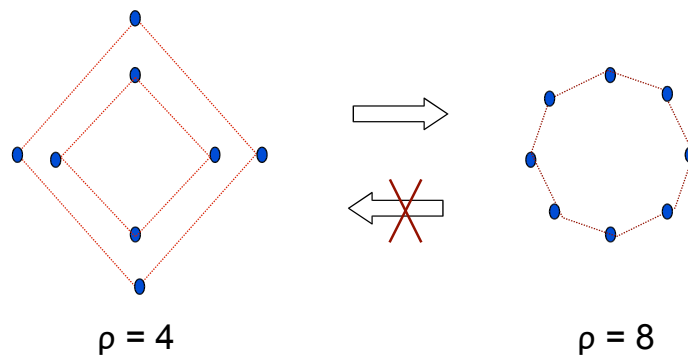
Necessary Condition:

In a series of patterns $(P_1, P_2, P_3, \dots, P_m)^\infty$,
Sizes must be preserved, i.e.
 $\text{size}(P_i) = \text{size}(P_j)$, for all i, j

Anonymous Robots

Necessary Condition (2):

In a series of patterns $(P_1, P_2, P_3, \dots, P_m)^\infty$,
The symmetricity must be preserved, i.e.
 $\rho(P_i) = \rho(P_j)$, for all i, j



Anonymous Robots

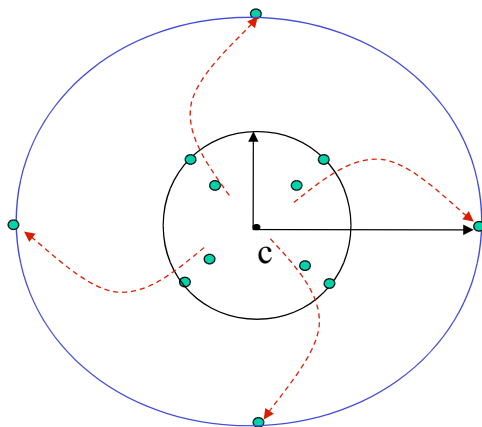
Sufficient Condition:

Theorem: The series $\langle P_1, P_2, P_3, \dots, P_m \rangle$ is formable, if

(1) $\text{Size}(P_i) = \text{Size}(P_j) = k$, for all i, j

(2) $\rho(P_a) = \rho(P_b) = q * \rho(C_0)$, for all a, b , ($q \geq 1$)

Algorithm for Anonymous Robots



Q-Symmetric Circular Configuration SCC(q):

Encodes info about a pattern using ratio of two radii

Algorithm for Anonymous Robots

- The symmetry in a configuration is maintained only if the adversary activates all robots with same view, at the same time.
- In a symmetric configuration, form $SCC(F(P_i))$ to signal the formation of pattern P_i
- In an asymmetric configuration, form $BCC(F(P_i))$ to signal the formation of the next pattern P_i
- Ensure that every possible intermediate configuration contains information about the pattern to be formed.

To Summarize...

Characterization of formable series:

The series $\langle P_1, P_2, P_3, \dots, P_m \rangle$ is formable,

- ♦ With k visibly distinct robots if
$$\text{Size}(P_i) \leq k, \text{ for all } i,$$
- ♦ With k indistinguishable ranked robots if
$$k > 3 \text{ and } \text{Size}(P_i) \leq k$$
- ♦ With k anonymous robots if
 - (1) $\text{Size}(P_i) = \text{Size}(P_j) = k$ and
 - (2) $\rho(P_a) = \rho(P_b) = q * \rho(C_0)$, for all a, b , ($q \geq 1$)

Conclusions

- We can overcome obliviousness of mobile robots not just for a single pattern but for a series of patterns.

Implement Global Memory without Local Memory!

- The technique of encoding information using ratio of distances can be used in other scenarios.
- These techniques can tolerate *measurement errors* up to certain extent.

Open Questions

- What happens when some robots are stalled permanently?
 - ☒ *Should the pattern include the failed robots?*
 - ☒ *How many failures can be tolerated?*
- What happens with robots have limited mobility?
 - ☒ *e.g. each robot moves unit distance in each step.*
- What happens in the case of limited visibility?

Thank you!

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