How simple robots benefit from looking back

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motivation



 Robot inside an unknown polygon



- Robot inside an unknown polygon
- Tasks:



- Robot inside an unknown polygon
- Tasks:
 - meet identical robots



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 - draw a map



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- Q: How simple can we make the robot?



- Robot inside an unknown polygon
- Tasks:
 - meet identical robots
 - draw a map
- Q: How simple can we make the robot?
 - \Rightarrow find simplistic design



visibilities



• Vertices are mut. visible: segment is inside polygon

visibilities



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- Vertices are mut. visible: segment is inside polygon
- Visibility graph: edge for every pair of visible verts



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 ⇒topological map



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 ⇒robots form a clique
- Mapping problem:
 ⇒reconstruct vis. graph



robot model



• We assume n is given



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 - while at a vertex:



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 ⇒ origin of last move

Vertices as Viewpoints

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view

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Capture information a robot can collect about v
view

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 ⇒view from v

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 - \Rightarrow collection of all paths

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• level-k-view:





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 - \Rightarrow collection of all paths
- level-1-view: $v^{1} = (L_{1}, L_{2}, ..., L_{d})$
- level-k-view:

view



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- level-1-view: $v^{1} = (L_{1}, L_{2}, ..., L_{d})$
- level-k-view: $v^{k} = (N_{1}^{k-1}, N_{2}^{k-1}, ..., N_{d}^{k-1})$

Vertices as Viewpoints classes

classes

 group all vertices with same v[∞] into classes C_i

classes



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 - \Rightarrow periodic on boundary

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$$|C_i| = I$$
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classes



- group all vertices with same v[∞] into classes C_i
 - ⇒periodic on boundary ⇒ $|C_i| = |C_j| \forall i,j$
- $|C_i| = I$: distinguishable \checkmark
- Norris95: vⁿ⁻¹ is enough!
 (same resulting classes)

The Class C* definition



definition



 C* is the lexicographically smallest class that forms a clique

definition



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- Will show:
 Every polygon has a class that forms a clique (!)

definition



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- Will show: Every polygon has a class that forms a clique (!)

 \Rightarrow C^{*} is well defined, unique



ears



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 - ⇒vertices in the same class as an ear are ears



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 - ⇒vertices in the same class as an ear are ears
- Every polygon has an ear


existence of a clique

• Cut ears repeatedly...



existence of a clique

Cut ears repeatedly...
 ⇒cut the entire class



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existence of a clique



- Cut ears repeatedly...
 ⇒cut the entire class
 ⇒no class will split!
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 ⇒must be a clique!
 - ⇒contains all vertices of some original class

⇒Every polygon has a class that is a clique!

problem re-definition



problem re-definition



 Views of level n-1 are sufficient to infer classes

problem re-definition



• Views of level n-l are sufficient to infer classes \Rightarrow task in terms of classes

problem re-definition



 Views of level n-1 are sufficient to infer classes
 ⇒task in terms of classes

• Given:

Meeting and Mapping problem re-definition



- Views of level n-1 are sufficient to infer classes
 ⇒task in terms of classes
- Given:
 - classes along boundary

problem re-definition



 Views of level n-1 are sufficient to infer classes \Rightarrow task in terms of classes

Given:

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- classes of neighbors

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Tasks:

problem re-definition



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Given:

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Tasks:

meet other robots

problem re-definition



 Views of level n-1 are sufficient to infer classes \Rightarrow task in terms of classes

Given:

- classes along boundary
- classes of neighbors

Tasks:

- meet other robots
- infer visibility graph

Meeting and Mapping meeting







Meeting and Mapping meeting 5 • C^{*} is unique 12 • C* can be inferred 10 9 neighbors class vert $C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3, C_4$ C **C**₂ $C_3, C_4, C_1, C_2, C_4, C_1, C_2, C_4, C_1$ 3 C_{4}, C_{1}, C_{2} 4 $C_4, C_1, C_2, C_4, C_1, C_2, C_3$ 5 $C_{1}, C_{2}, C_{4}, C_{1}, C_{2}, C_{3}, C_{4}$ $C_{2,C}$ 6 C $C_{3}, C_{4}, C_{1}, C_{2}, C_{4}, C_{1}, C_{2}, C_{4}, C_{1}$ 7 Ca C_4, C_1, C_2 C₄ 8 $C_1, C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3$ 9 C₂,C₄,C₁,C₂,C₄,C₁,C₂,C₃,C₄ Ci 10 C_2 C₃,C₄,C₁,C₂,C₄,C₁,C₂,C₄,C₁ C_3 C_{4}, C_{1}, C_{2} C_4 $C_1, C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3$ 12

Meeting and Mapping meeting 5 • C^{*} is unique 12 C^{*} can be inferred 10 9 class neighbors vert C₂,C₄,C₁,C₂,C₄,C₁,C₂,C₃,C₄ C \Rightarrow Meeting is trivial: move C₃,C₄,C₁,C₂,C₄,C₁,C₂,C₄,C **C**₂ 3 C_{4}, C_{1}, C_{2} along boundary until a C4,C1,C2,C4,C1,C2,C3 5 $C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3, C_4$ vertex in C^* 6 C $C_{3}, C_{4}, C_{1}, C_{2}, C_{4}, C_{1}, C_{2}, C_{4}, C_{1}$ 7 Ca C_4, C_1, C_2 8 C₄ $C_1, C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3$ $C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3, C_4$ 9 10 C_2 C₃,C₄,C₁,C₂,C₄,C₁,C₂,C₄,C₁ C_3 C_{4}, C_{1}, C_{2} C_4 $C_1, C_2, C_4, C_1, C_2, C_4, C_1, C_2, C_3$ 12





 need to identify vertices in the list of neighbors



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- If own class is a clique:



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 ⇒segment + class → ID
- \Rightarrow C^{*} vertices can be done



•	Identify edges (v_i, v_{i+k}) of
	increasing distances k

vert	class	neighbors
i	С	$C_{A}, C_{B}, C_{C},, C_{L},, C_{M},, C_{X}, C_{Y}, C_{Z}$



vert	class		neighbors	
i	С	$C_A, C_B, C_C,$,C _L ,,C _M ,	,C _X ,C _Y ,C _Z
		\checkmark	?	\checkmark

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- Is the next unidentified vertex v_j = v_{i+k} or not?



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- y is number of dist. k-l
 backward-edges of v_{i+k}



vert	class		neighbors	S
i	С	$C_A, C_B, C_C, .$,C _L ,,C _M	,,C _X ,C _Y ,C _Z
		·		
		5	2	\checkmark
		•	•	•

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- y is number of dist. k-l
 backward-edges of v_{i+k}
- Move to v_j and look back • $v_j = v_{i+k} \Rightarrow LB = -(y+1)$



vert	class	neighbors
i	С	$C_{A}, C_{B}, C_{C},, C_{L},, C_{M},, C_{X}, C_{Y}, C_{Z}$



•	We show:
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- Assume $v_j \neq v_{i+k}$ but LB = -(y+1)



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	$LB = -(y+I) \Rightarrow v_j = v_{i+k}$

- Assume $v_j \neq v_{i+k}$ but LB = -(y+1)
- v_{i+k} and v_j are same class

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- v_{i+k} and v_j are same class
 ⇒v_j has y back-edges



vert	class	neighbors
i	C	$C_{A}, C_{B}, C_{C},, C_{L},, C_{M},, C_{X}, C_{Y}, C_{Z}$

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- Assume $v_j \neq v_{i+k}$ but LB = -(y+1)
- v_{i+k} and v_j are same class
 ⇒v_j has y back-edges
 ⇒all back-edges identified



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 ⇒all back-edges identified
- Use C^{*} as a frame



vert	class	neighbors
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 ⇒2 cases



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- We show: $LB = -(y+1) \Rightarrow v_j = v_{i+k}$
- Use C^* as frame \Rightarrow 2 cases
- Case I: there are multiple
 C* between v_i,v_j



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- Use C^* as frame $\Rightarrow 2$ cases
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 C* between v_i,v_j
 - ⇒forbidden polygon 4
- Case II: there is one C*



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Case II: there is one C^{*}
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Case II: there is one C^{*}
 ⇒forbidden polygon 4

 \Rightarrow Criterion for deciding $v_j = v_{i+k}$

• The class of a vertex can be determined in finite time by a look-back robot

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- Every polygon has a class C* that forms a clique

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- The class of a vertex can be determined in finite time by a look-back robot
- Every polygon has a class C* that forms a clique
- Because of this, robots can always meet "easily"
- C* can be used as frame to infer the visibility graph

Thank you!